

Your Data is Blue, But Can You Prove it?

Interactive Proofs for Distribution Properties

(Mostly) based on works with Guy Rothblum

A SURVEY ON DISTRIBUTION TESTING

Your Data is Big. But is it Blue?

Clément L. Canonne¹

May 8, 2015

Your Data

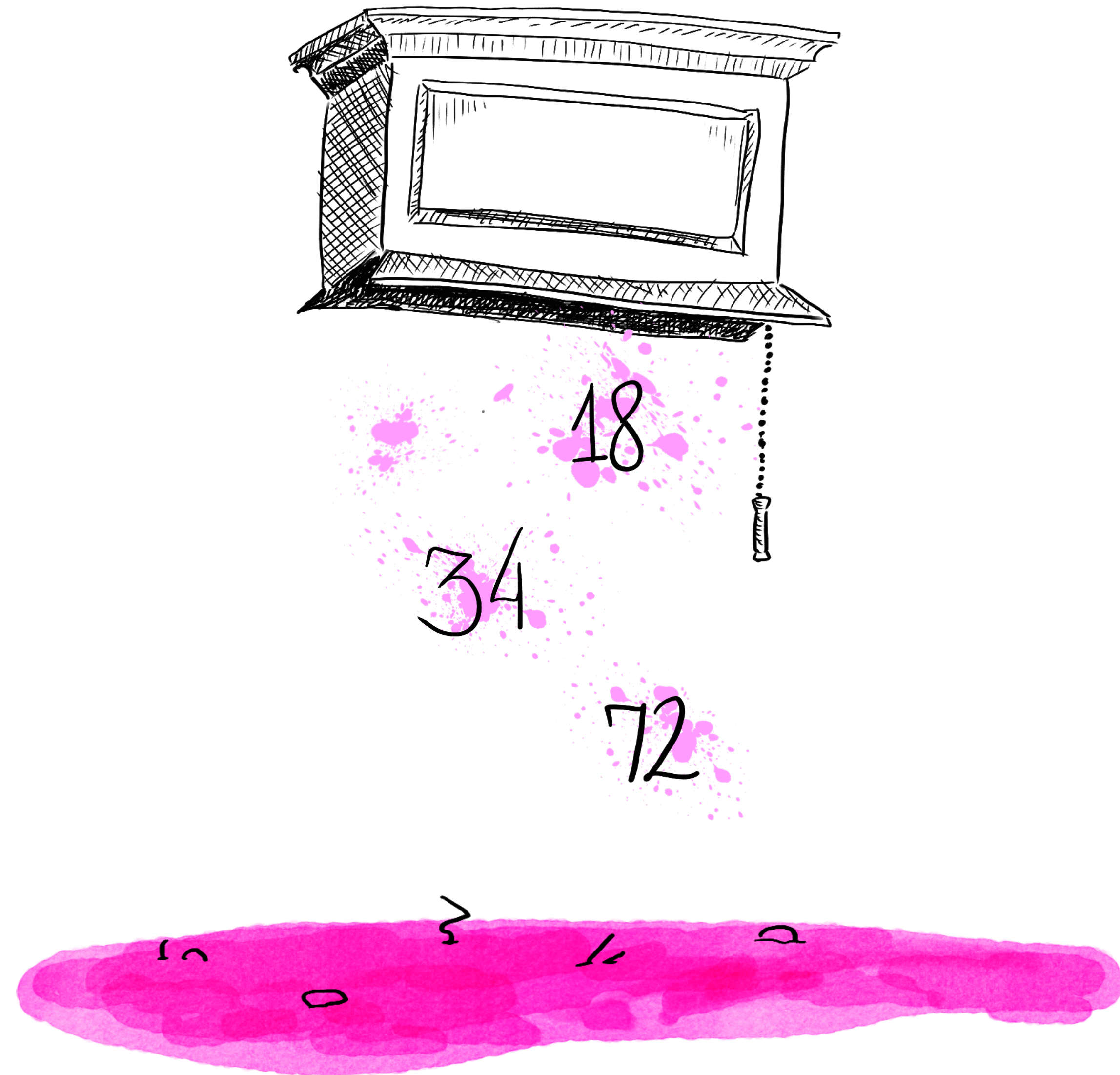
Prove it!

Interactive Proofs for Distribution Properties

(Mostly) based on works with Guy Rothblum

Introduction: Testing Distributions

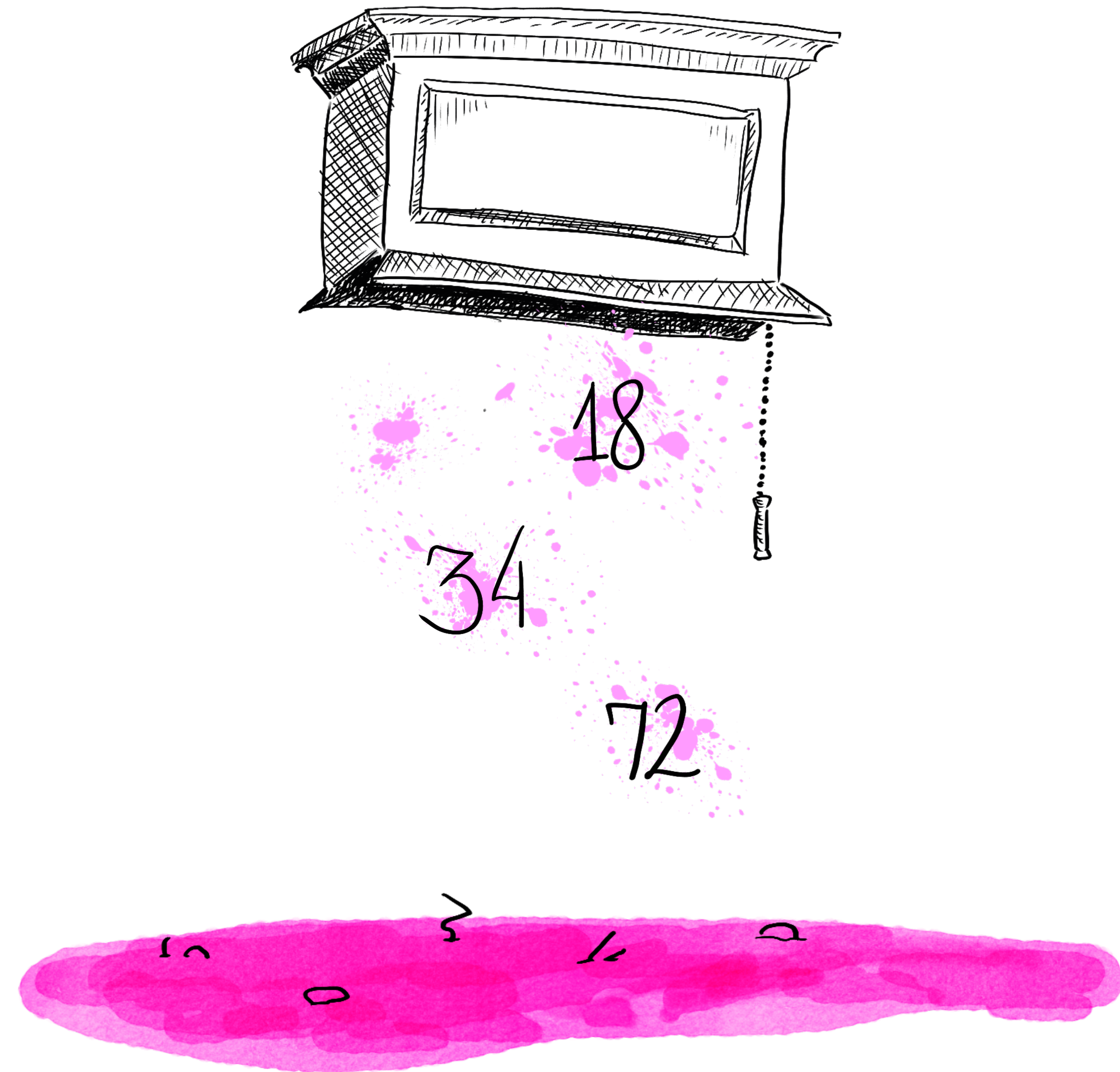
Distribution **D** over **[N]**



Introduction: Testing Distributions

Learn the distribution! $O(N)$

Distribution **D** over **[N]**

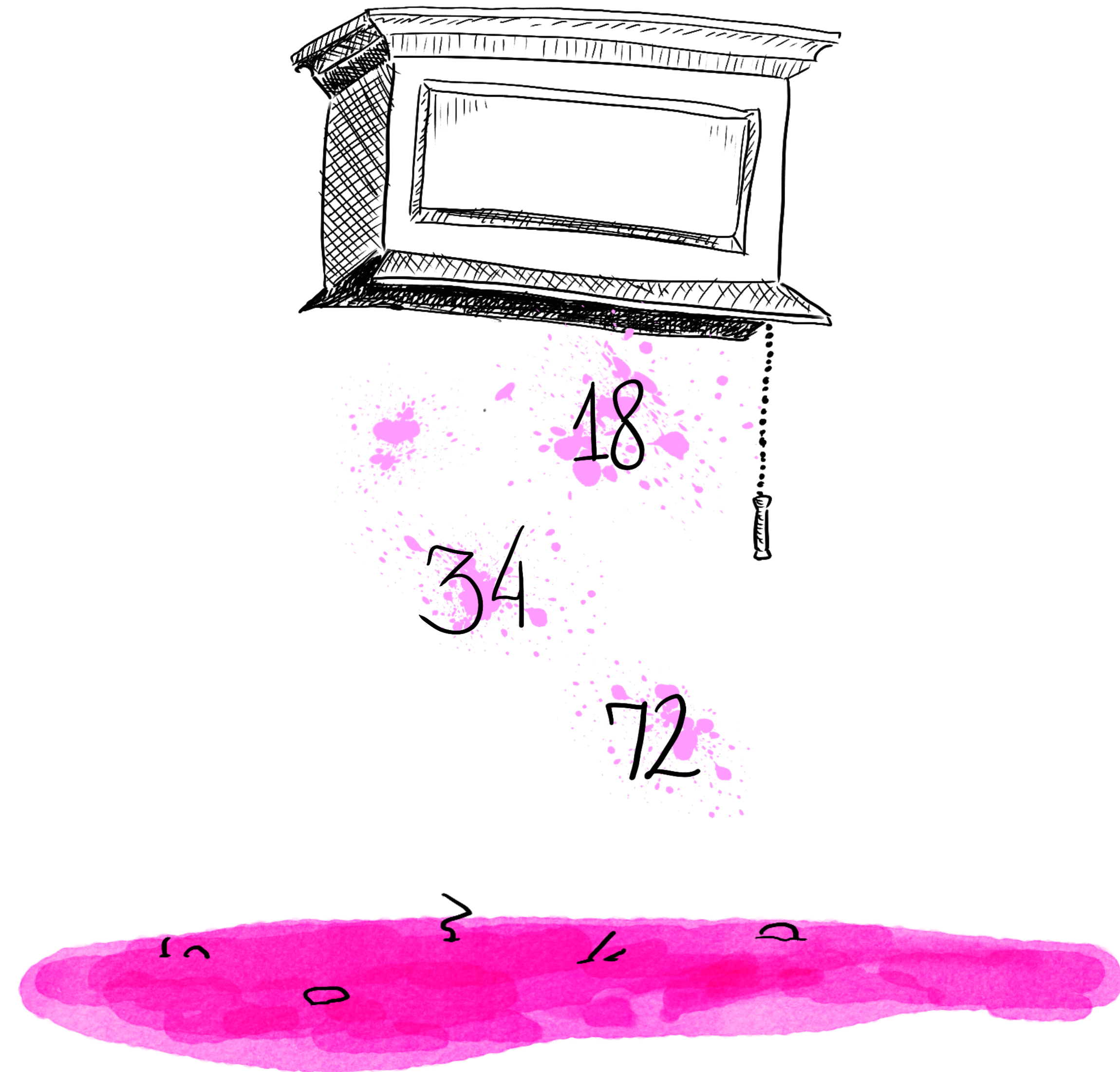


Introduction: Testing Distributions

Learn the distribution! $O(N)$

Easy: approximate mean, variance. $\text{poly}(\epsilon^{-1})$

Distribution **D** over **[N]**



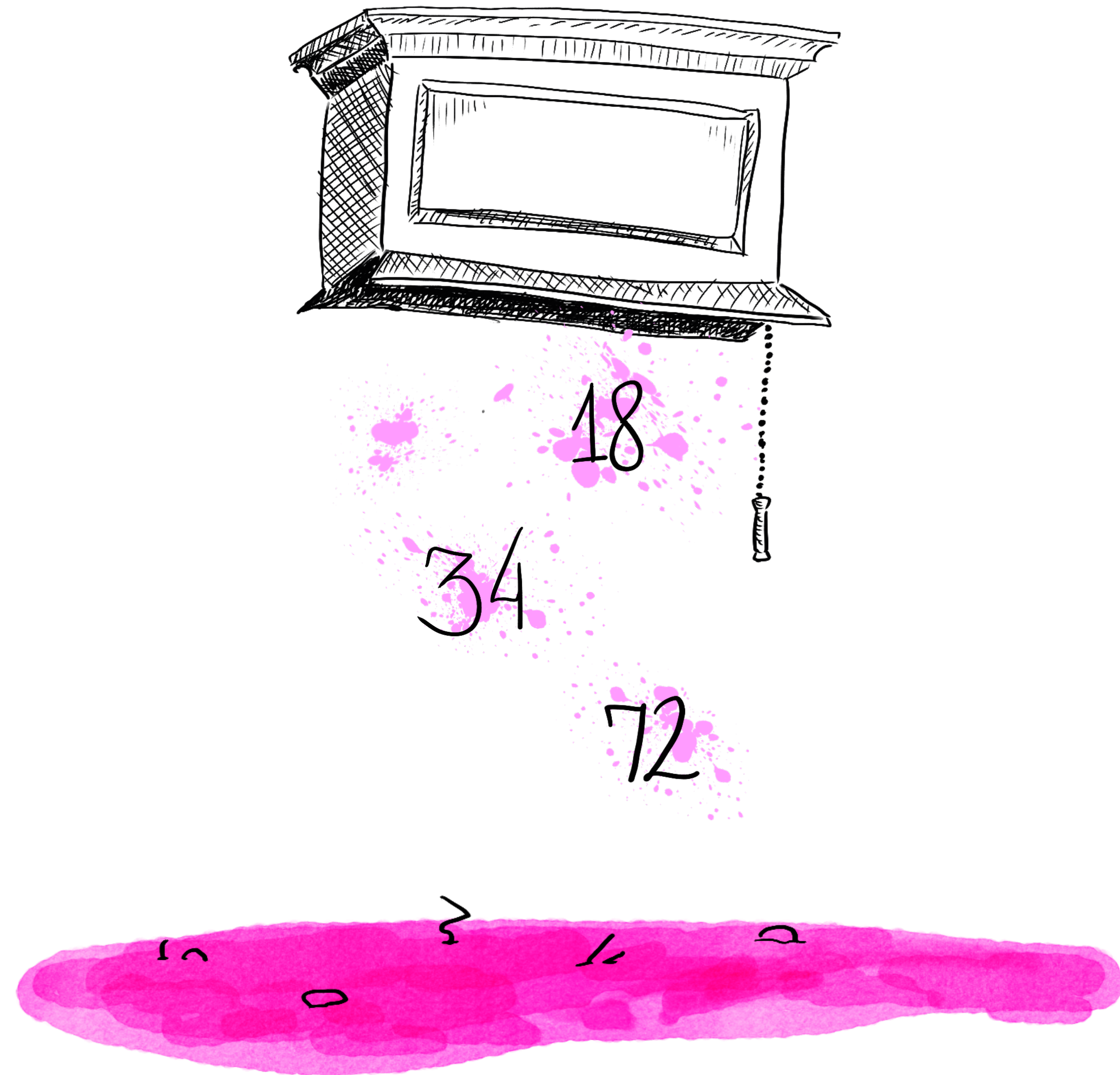
Introduction: Testing Distributions

Learn the distribution! $O(N)$

Easy: approximate mean, variance. $\text{poly}(\epsilon^{-1})$

Hard: is D uniform or *far* from it? $\Theta(\sqrt{N})$
[GR'00,...,Pan'08]

Distribution D over $[N]$



Introduction: Testing Distributions

Learn the distribution! $O(N)$

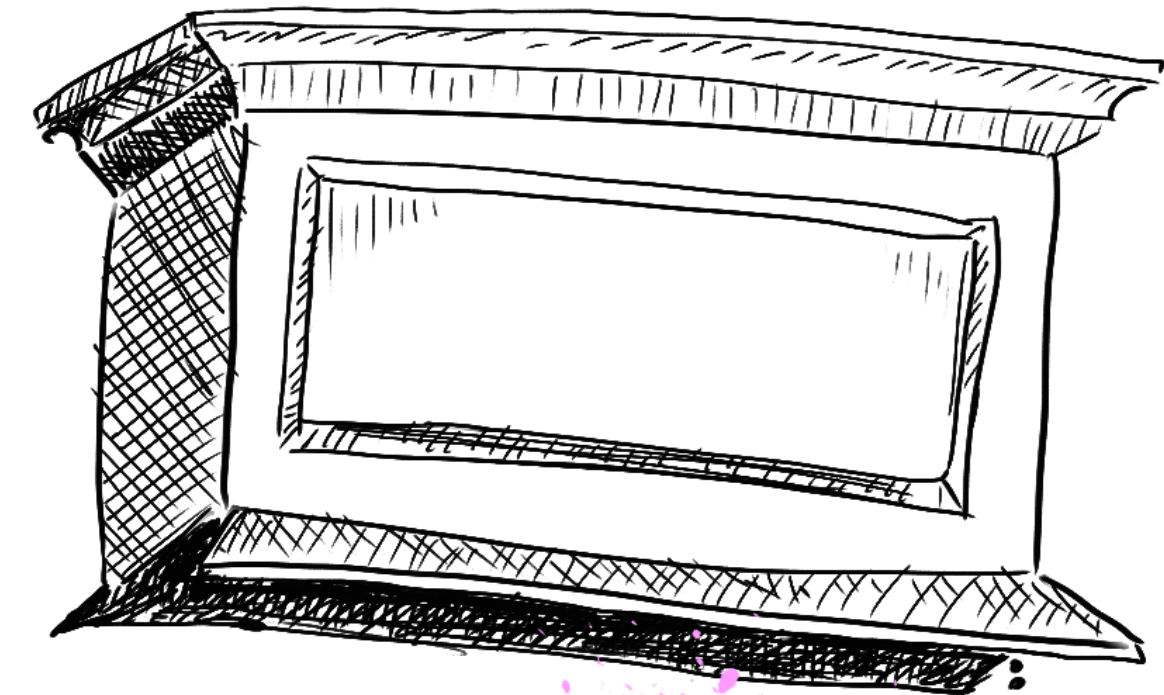
Easy: approximate mean, variance. $\text{poly}(\epsilon^{-1})$

Hard: is D uniform or *far* from it? $\Theta(\sqrt{N})$
[GR'00,...,Pan'08]

Harder: $\Theta(N / \log N)$

- Is the distribution *monotone*? [BFRV'10,AGPRY'19,]
- Approximating *distance from uniform*. [VV'10]
- Approximating *Shannon entropy*. [VV'10]

Distribution D over $[N]$



18
34

72



Introduction: Testing Distributions

Distribution **D** over **[N]**

Learn the distribution! $O(N)$

Easy: approx

Hard: is D

Harder: Θ

Very limited access to the input:

$((1, D(1)), (2, D(2)), \dots, (N, D(N)))$

- Is the distribution *monotone*? [BFRV'10, AGPRY'19,]
- Approximating *distance from uniform*. [VV'10]
- Approximating *Shannon entropy*. [VV'10]

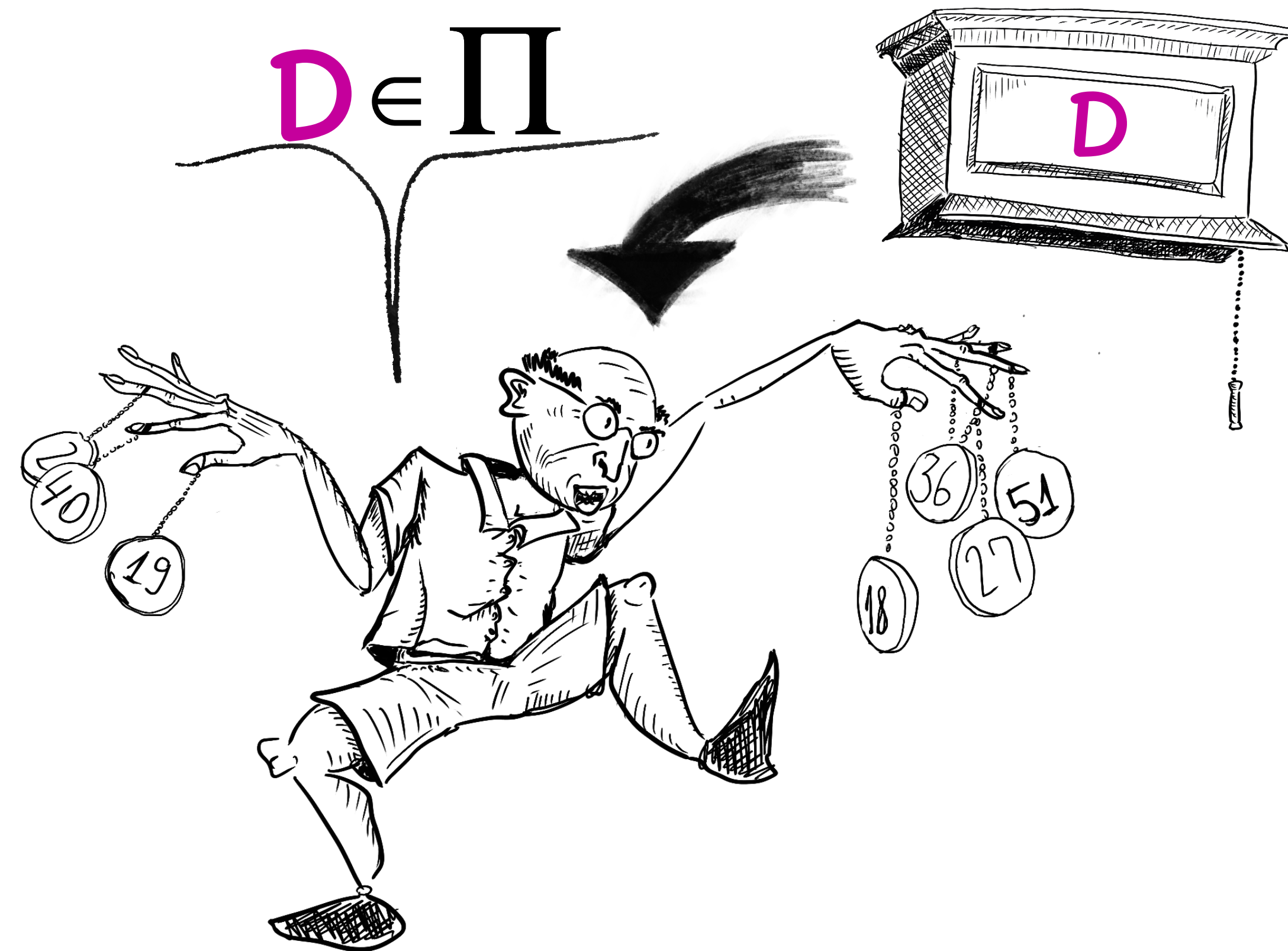


34

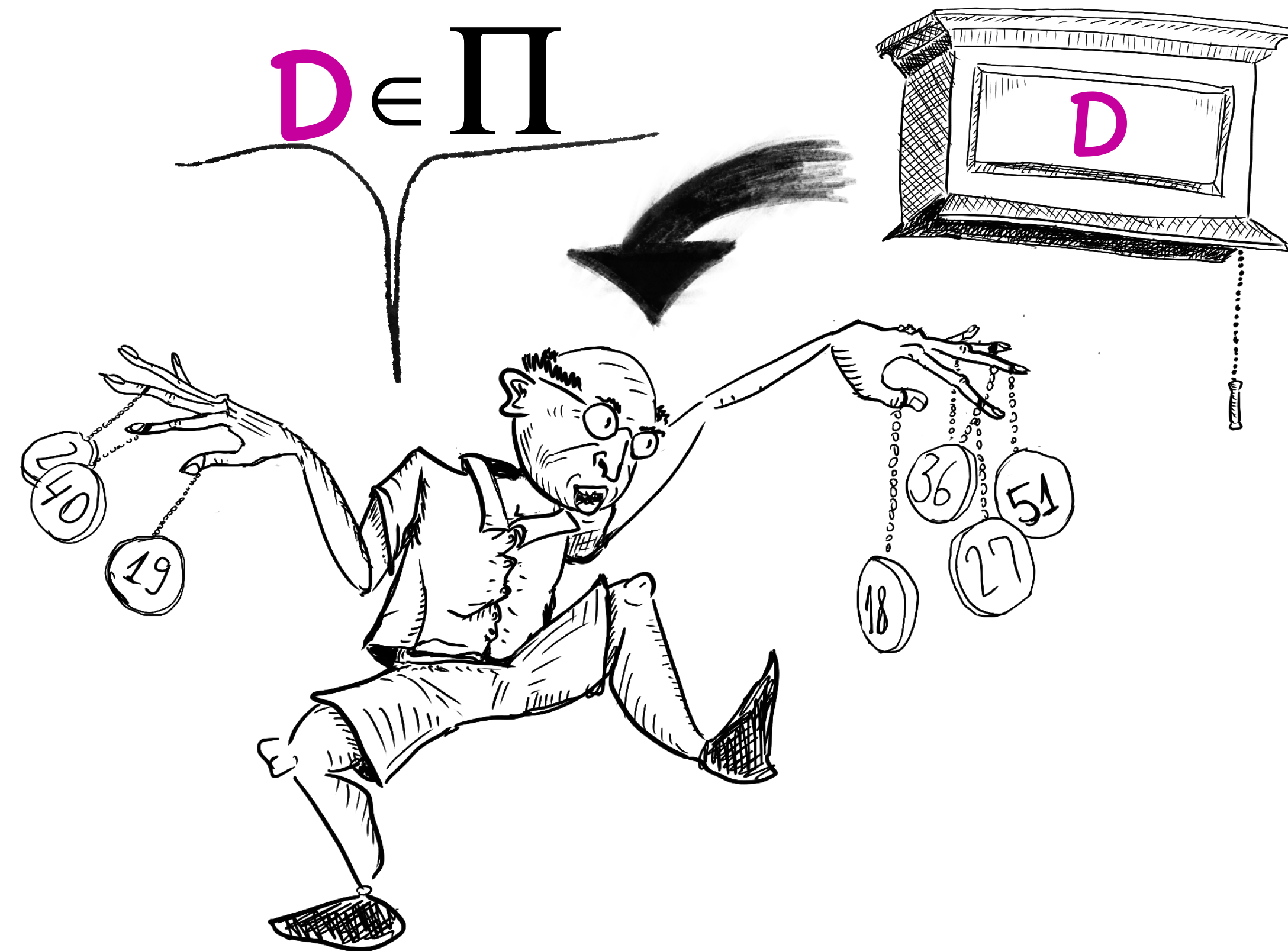
72



Verifying Distribution Properties: Warmup

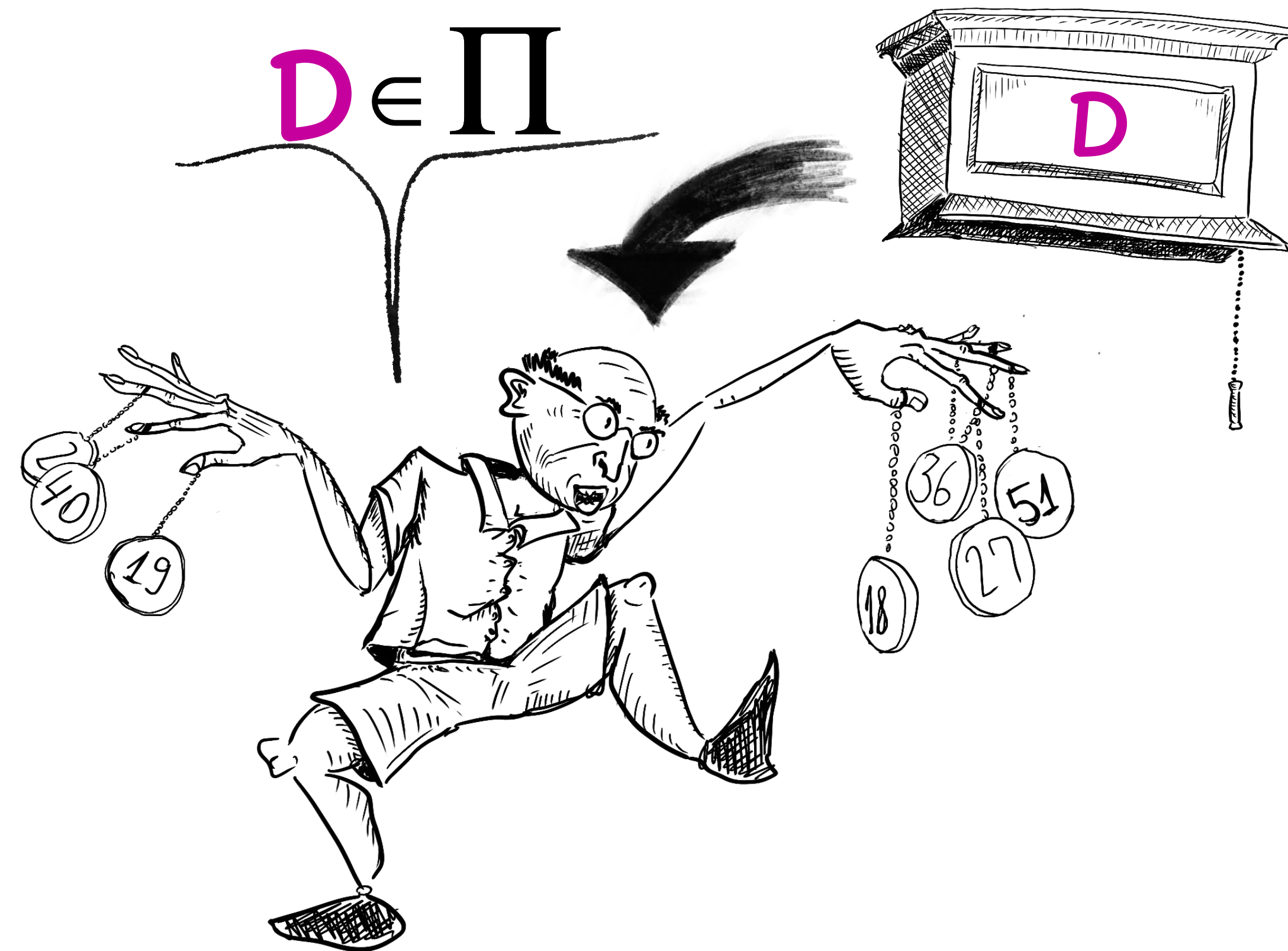


Verifying Distribution Properties: Warmup



Q: Can we *verify* claims with fewer samples?

Verifying Distribution Properties: Warmup



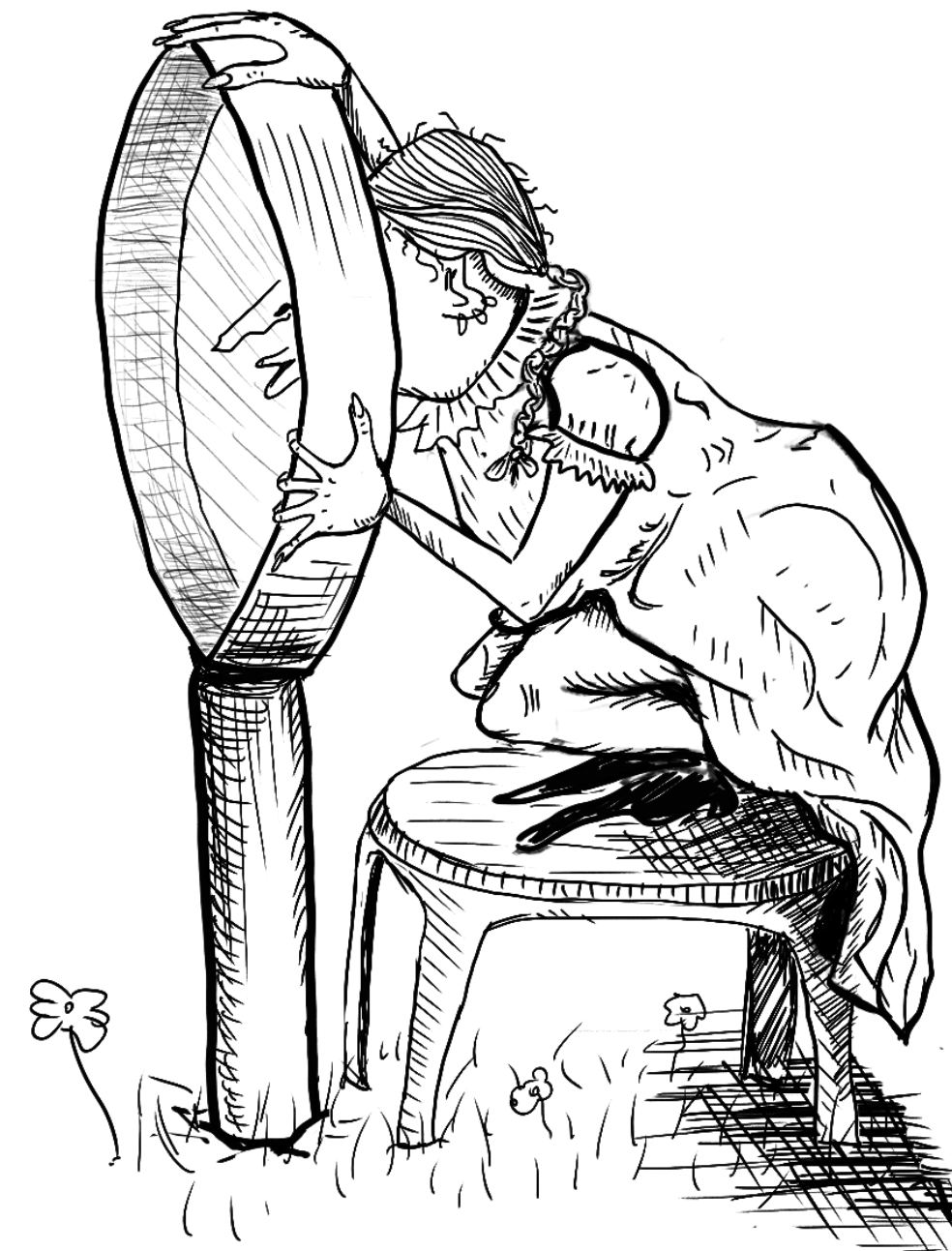
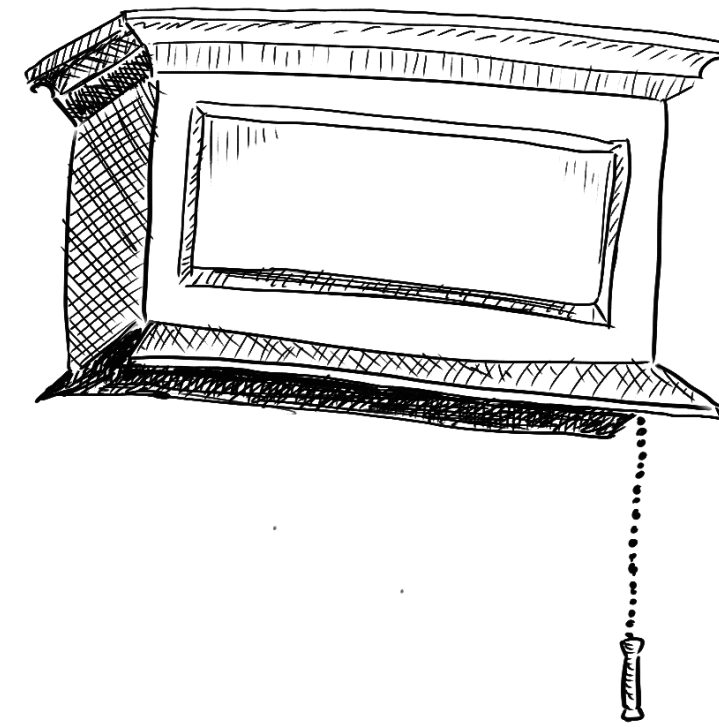
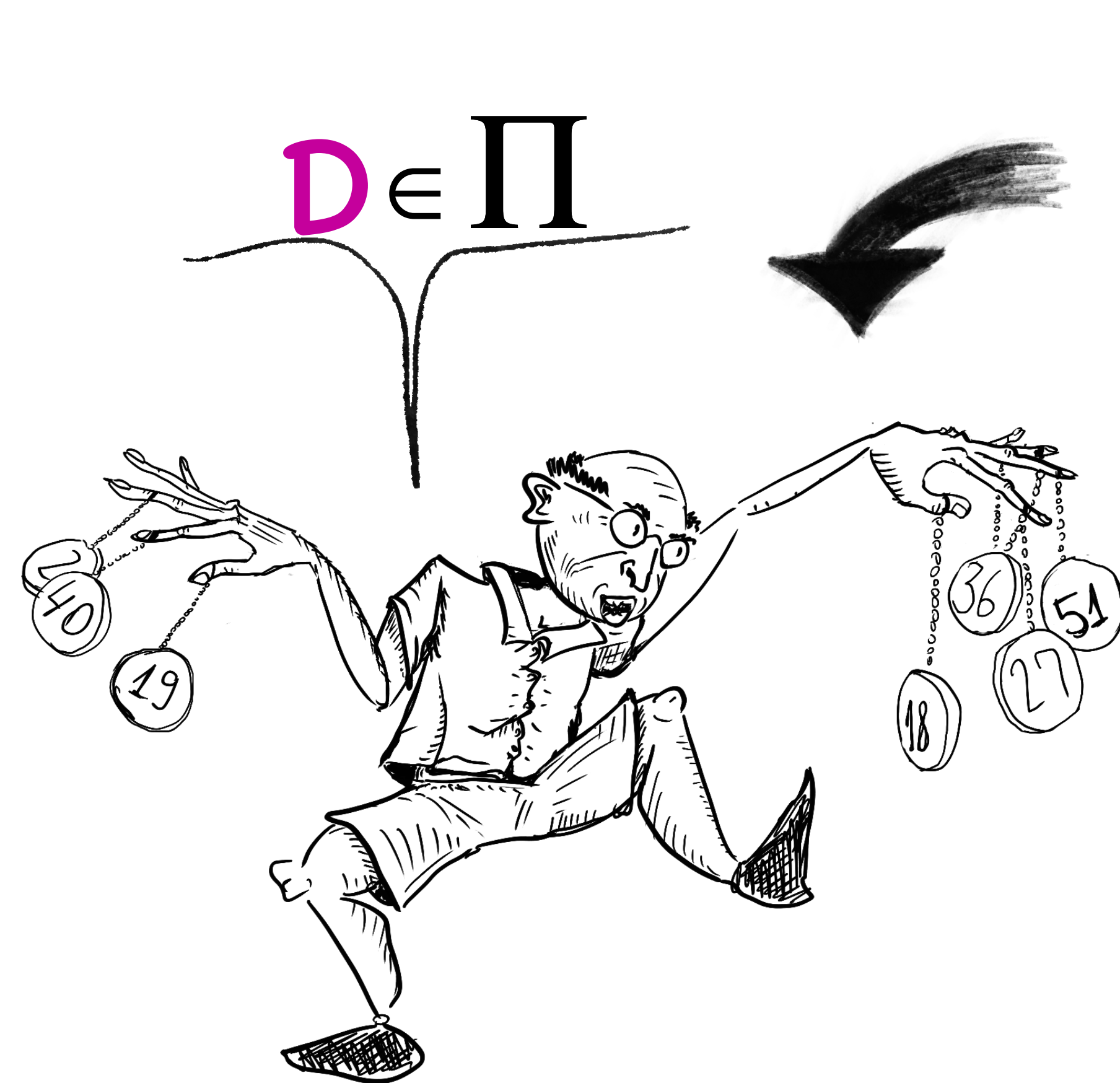
Q: Can we *verify* claims with fewer samples?

- Verifiable *data science*!
- Verifying *AI*?
- Natural theoretical question - new perspectives and insights on old problems.

Verifying Distribution Properties: Warmup

Distribution **D** over **[N]**

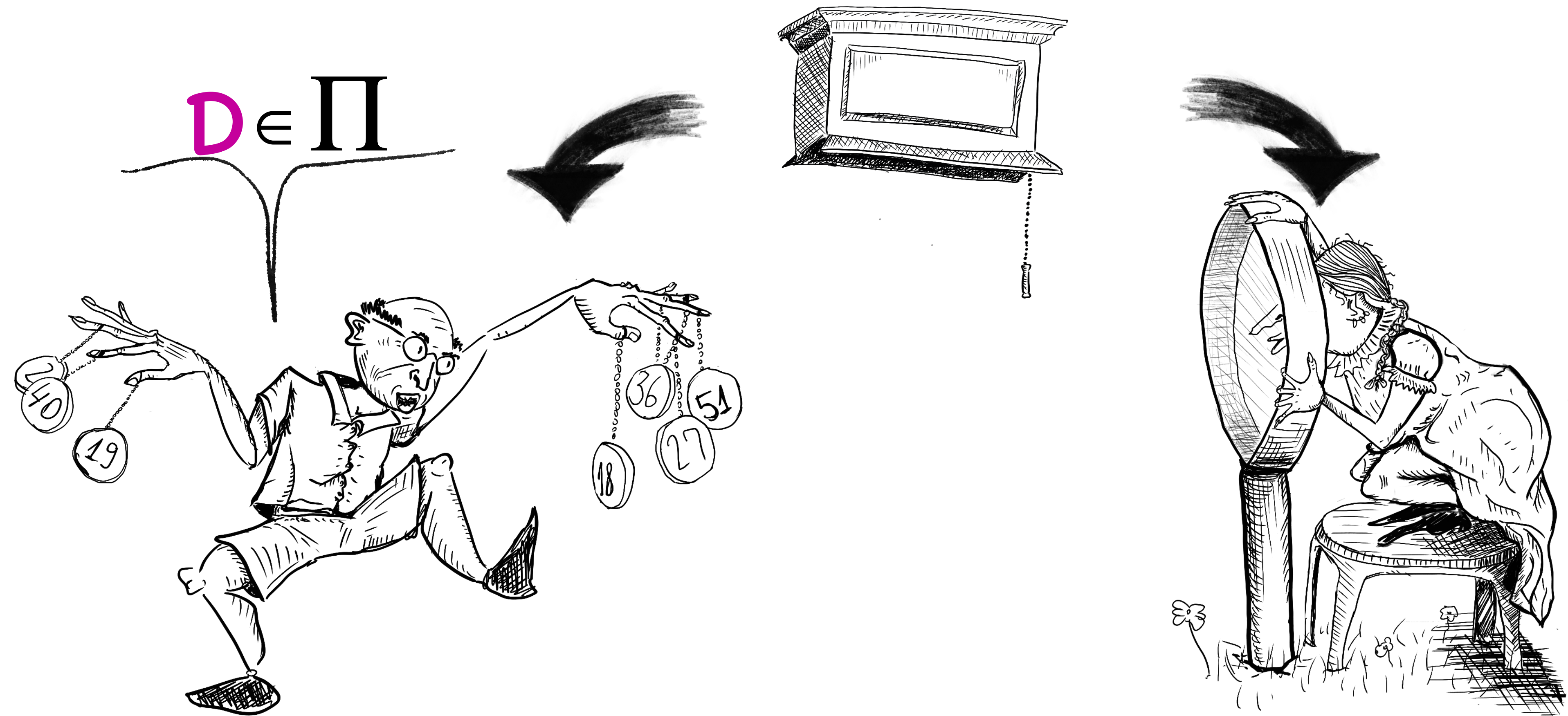
D $\in$ Π



Verifying Distribution Properties: Warmup

Distribution **D** over **[N]**

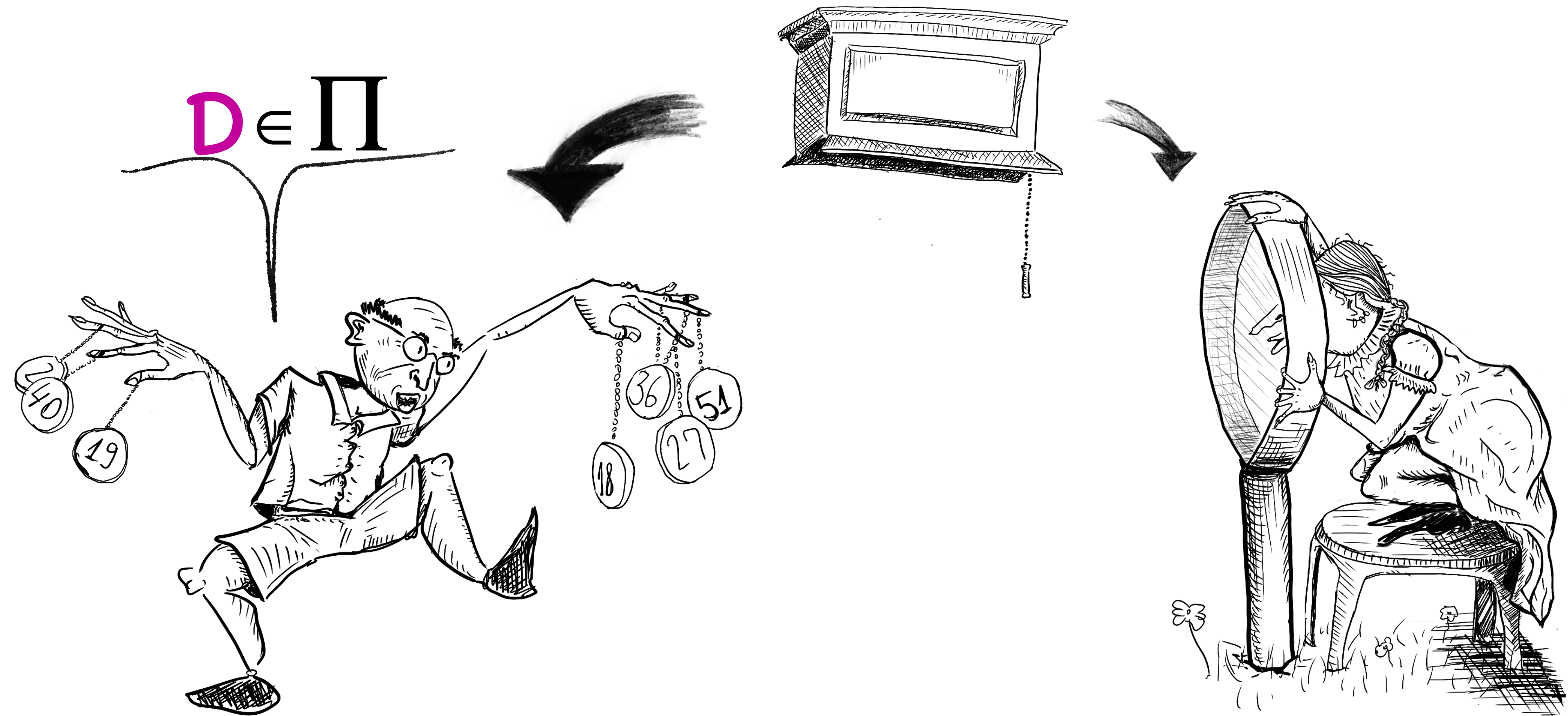
D $\in$ Π



Verifying Distribution Properties: Warmup

Distribution $\mathbf{D}$ over $\mathbf{[N]}$

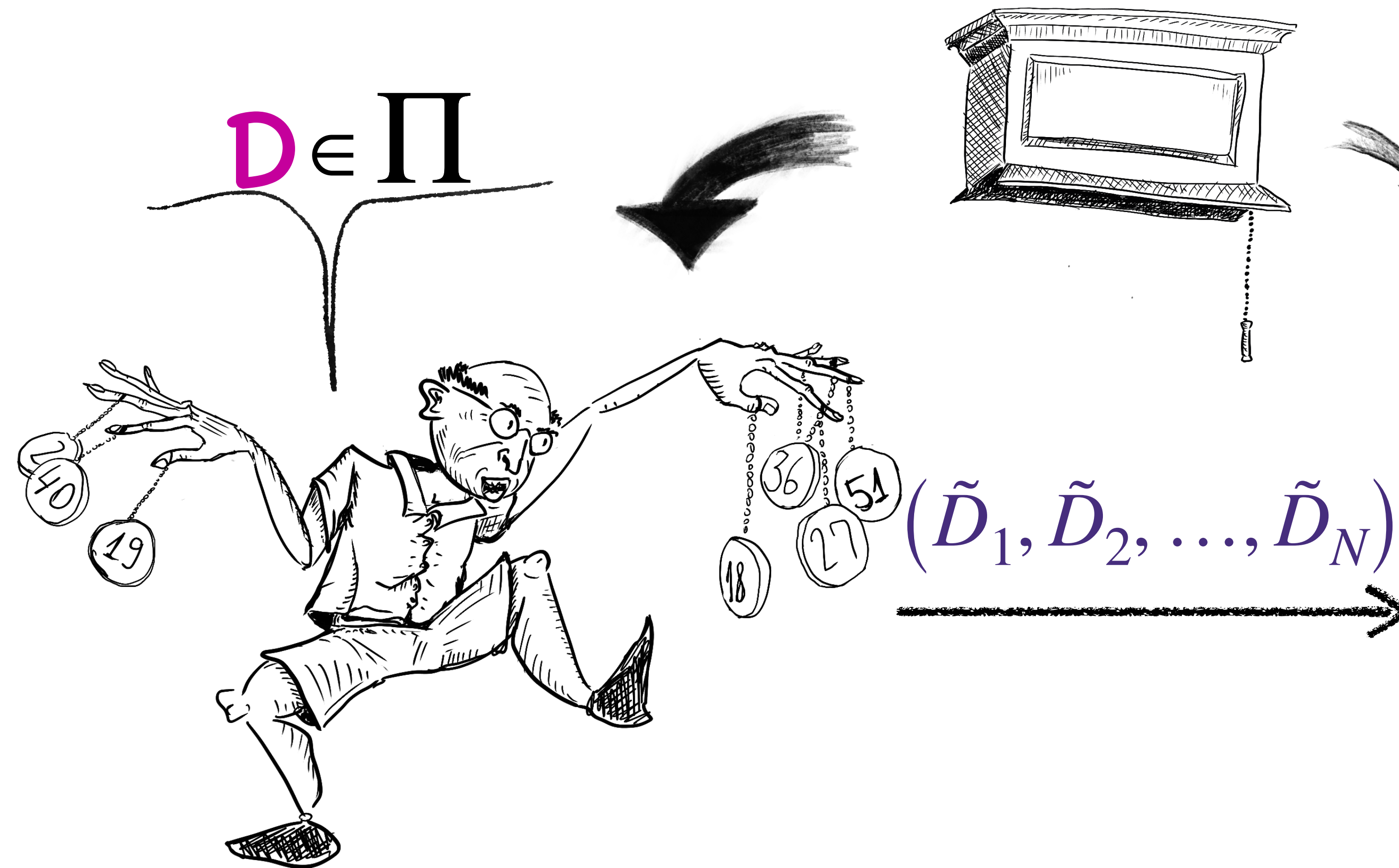
$\mathbf{D} \in \Pi$



Verifying Distribution Properties: Warmup

Distribution $\mathbf{D}$ over $[N]$

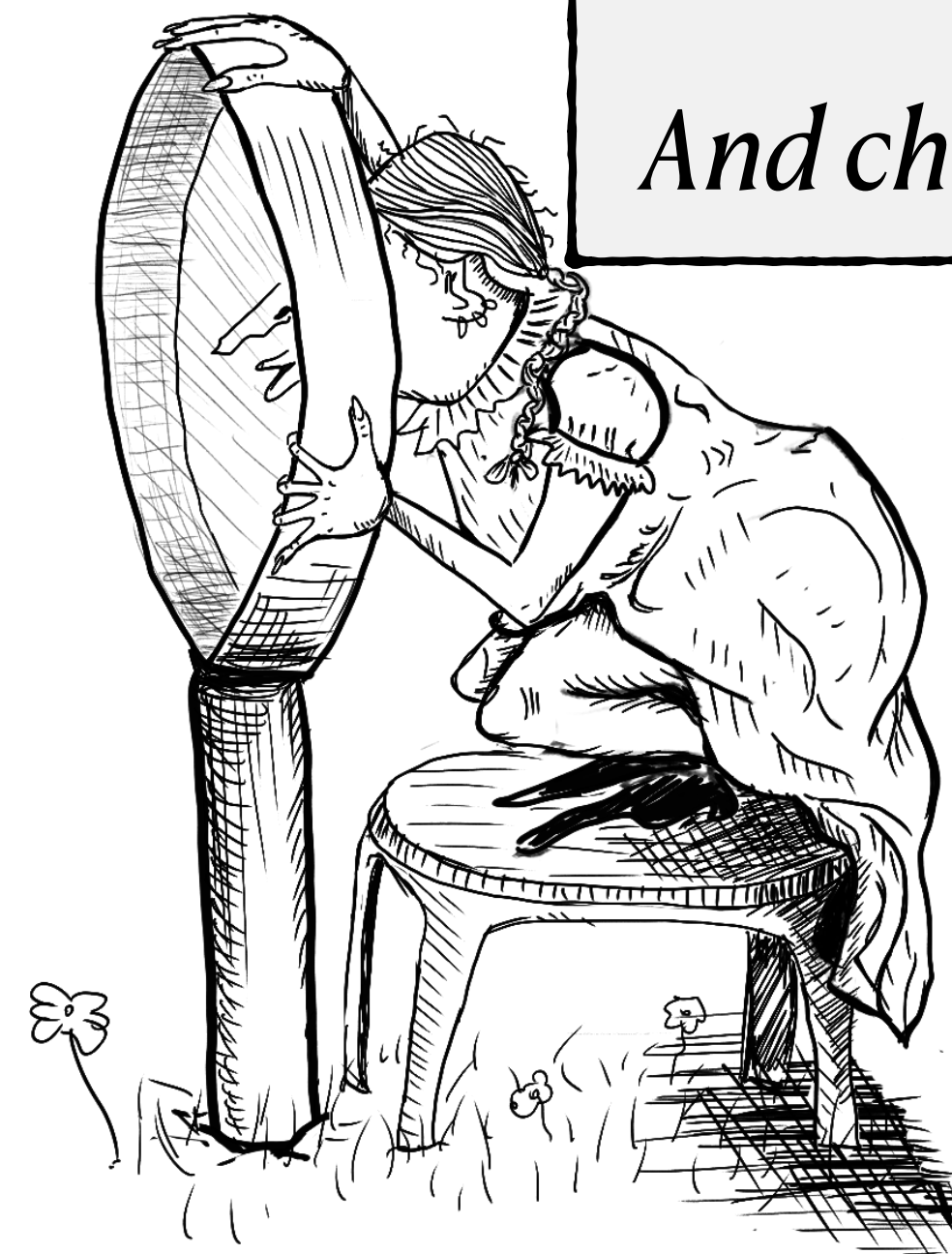
$\mathbf{D} \in \Pi$



$(\tilde{D}_1, \tilde{D}_2, \dots, \tilde{D}_N)$

Samples: $O(N)$

Comm.: $\tilde{O}(N)$



Samples: $O(\sqrt{N})$

Run *Identity Tester* [BFFKR'01, VV'14]
to verify:

$$\delta_{TV}(D, \tilde{D}) \leq 0.01$$

And check $\tilde{D}$ is close to Π

Verifying Distribution Properties: Warmup

Distribution $\mathbf{D}$ over $\mathbf{[N]}$

$\mathbf{D} \in \mathbf{\Pi}$

Run *Identity Tester* [BFFKR'01, VV'14]
to verify:

$$\delta_{TV}(D, \tilde{D}) \leq 0.01$$

And check $\tilde{D}$ is close to $\mathbf{\Pi}$

Can we do better?

$(\tilde{D}_1, \tilde{D}_2, \dots, \tilde{D}_n)$

Samples: $O(N)$

Comm.: $\tilde{O}(N)$

Samples: $O(\sqrt{N})$

Verifying Distribution Properties: Warmup

Distribution $\mathbf{D}$ over $\mathbf{[N]}$

$\mathbf{D} \in \mathbf{\Pi}$

Run *Identity Tester* [BFFKR'01, VV'14]
to verify:

$$\delta_{TV}(D, \tilde{D}) \leq 0.01$$

And check $\tilde{D}$ is close to $\mathbf{\Pi}$

Can we do better?

What properties are efficiently verifiable?

Samples: $O(N)$

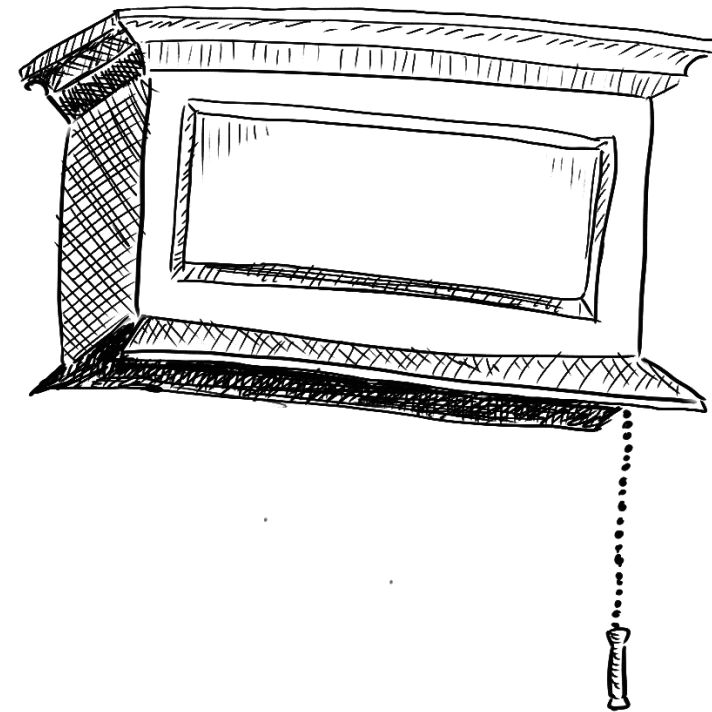
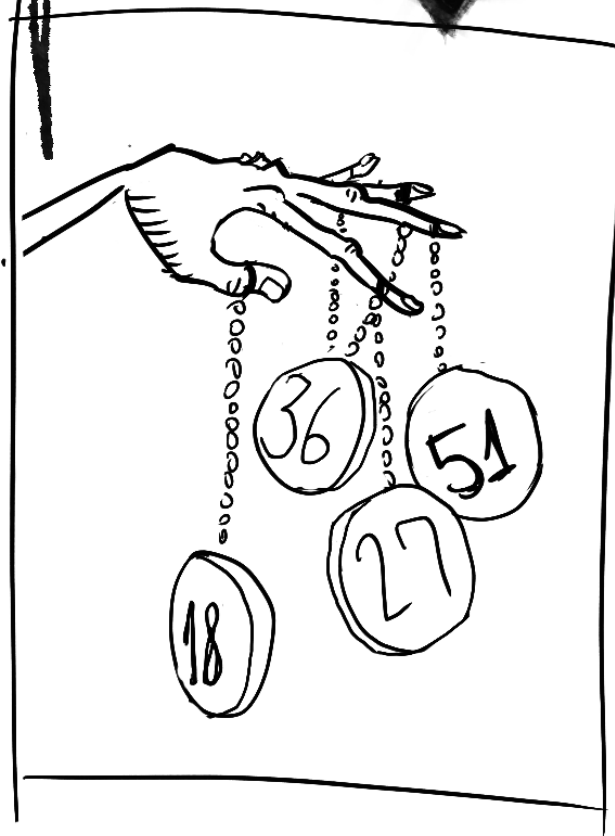
Comm.: $\tilde{O}(N)$

Samples: $O(\sqrt{N})$

Proof Systems for Distributions [GMR'85, CG'17]

$D \in \Pi$

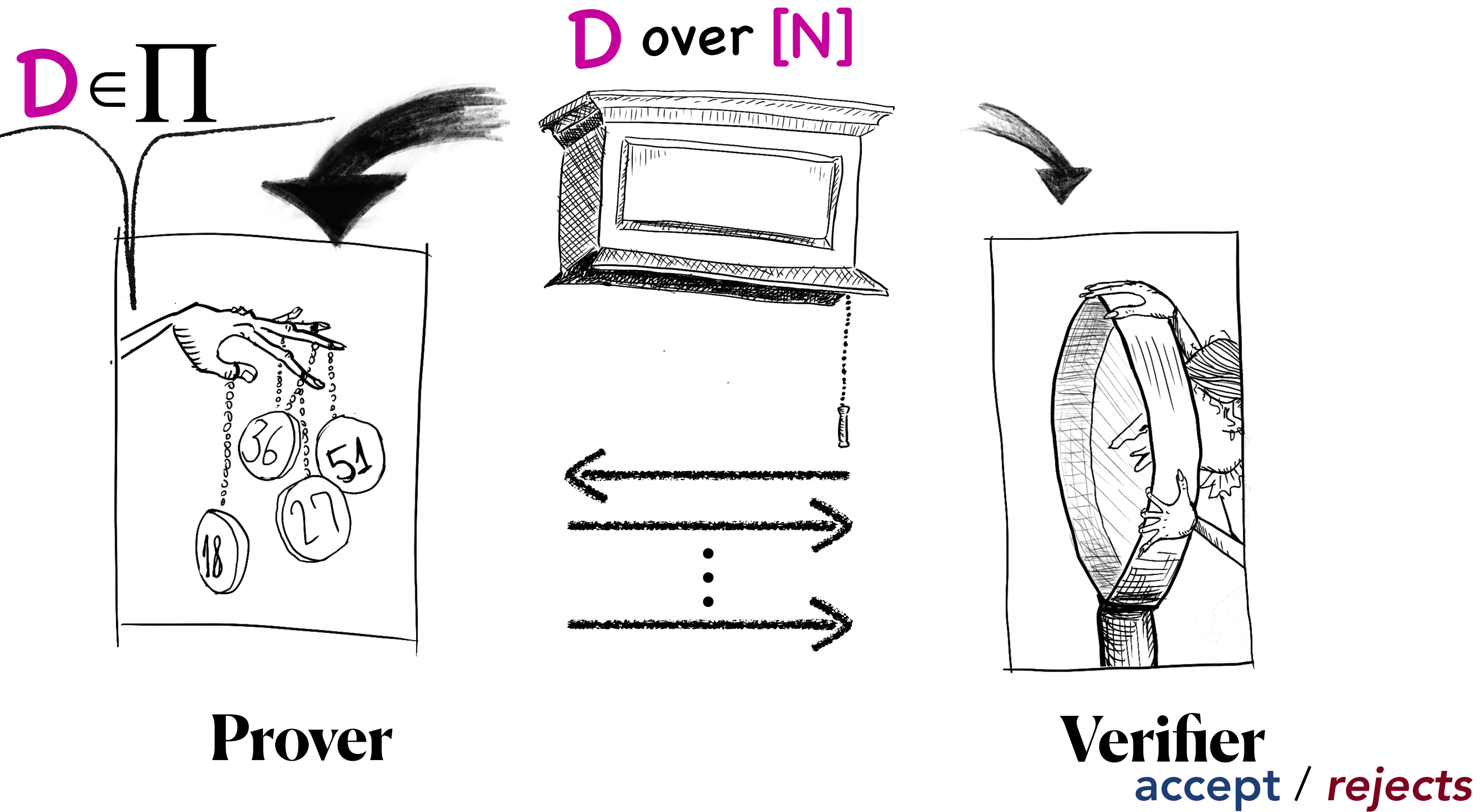
D over $[N]$



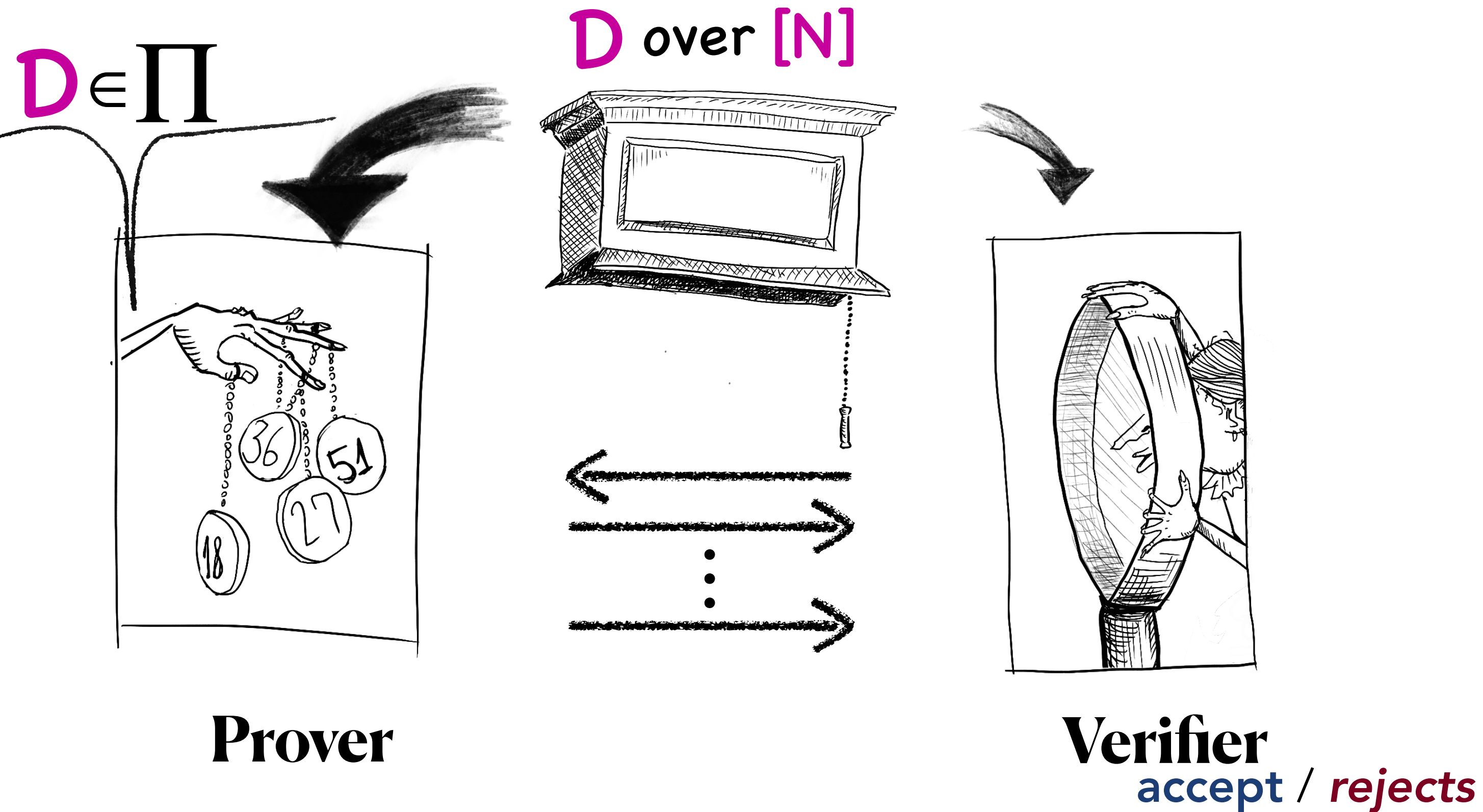
Prover



Proof Systems for Distributions [GMR'85, CG'17]



Proof Systems for Distributions [GMR'85, CG'17]



Completeness

$$\mathbf{D} \in \Pi$$

$\Rightarrow V$ *accepts* whp.

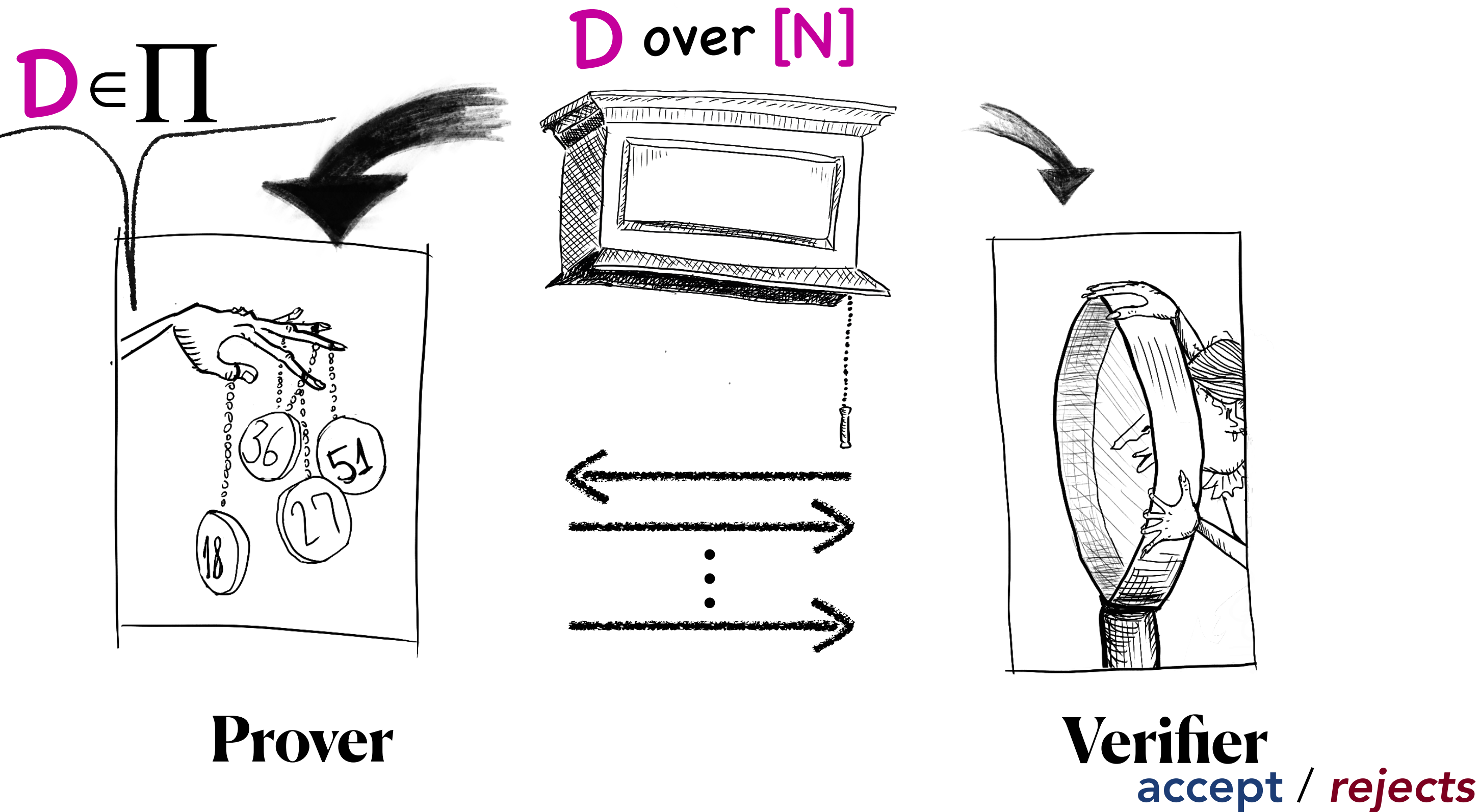
Soundness

$$\delta_{TV}(\mathbf{D}, \Pi) \geq \epsilon$$

$\Rightarrow \forall$ (cheating) strategy P^* ,
 V *rejects* whp.



Proof Systems for Distributions [GMR'85, CG'17]



Completeness

$D \in \Pi$
 $\Rightarrow V$ **accepts** whp.

Soundness

$\delta_{TV}(D, \Pi) \geq \epsilon$
 $\Rightarrow \forall$ (cheating) strategy P^* ,
 V **rejects** whp.

P 's Samples, runtime:
 $\text{poly}(N)$

Communication:
 $o(N)$, small # rounds

V 's Samples, runtime:
 $o(\text{optimal } \Pi \text{ tester})$

**Quantities of
interest**



(Tolerant) Proof Systems for Distributions [GMR'85, CG'17]

$$\delta_{TV}(\mathbf{D}, \Pi) \leq \epsilon_{close}$$

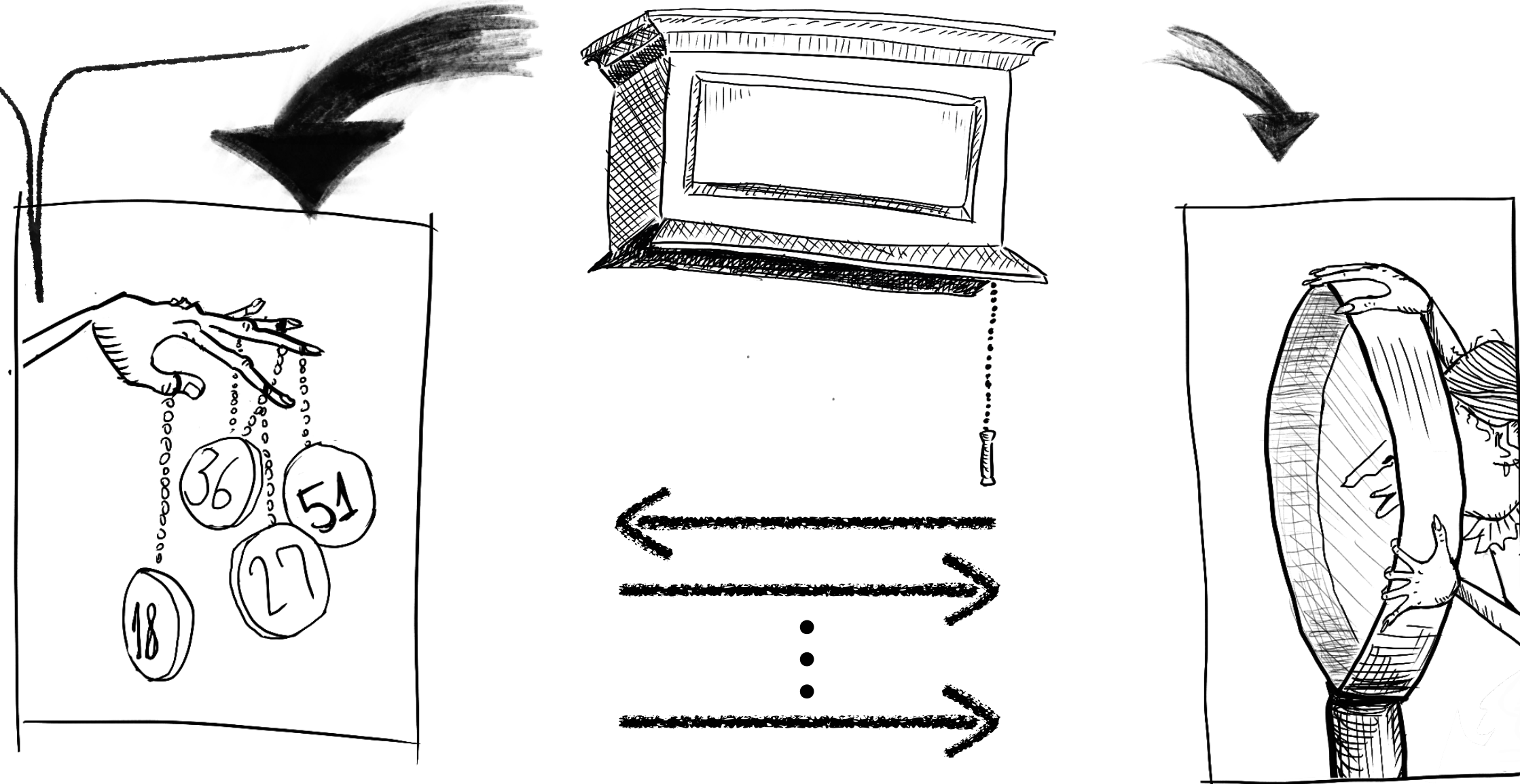
$\mathbf{D}$ over $[N]$

Completeness

$$\delta_{TV}(\mathbf{D}, \Pi) \leq \epsilon_{close} \\ \Rightarrow V \text{ *accepts* whp.}$$

Soundness

$$\delta_{TV}(\mathbf{D}, \Pi) \geq \epsilon_{far} \\ \Rightarrow \forall (\text{cheating}) \text{ strategy } P^*, \\ V \text{ *rejects* whp.}$$



Prover

Verifier
accept / *rejects*

P's Samples, runtime:
 $\text{poly}(N)$

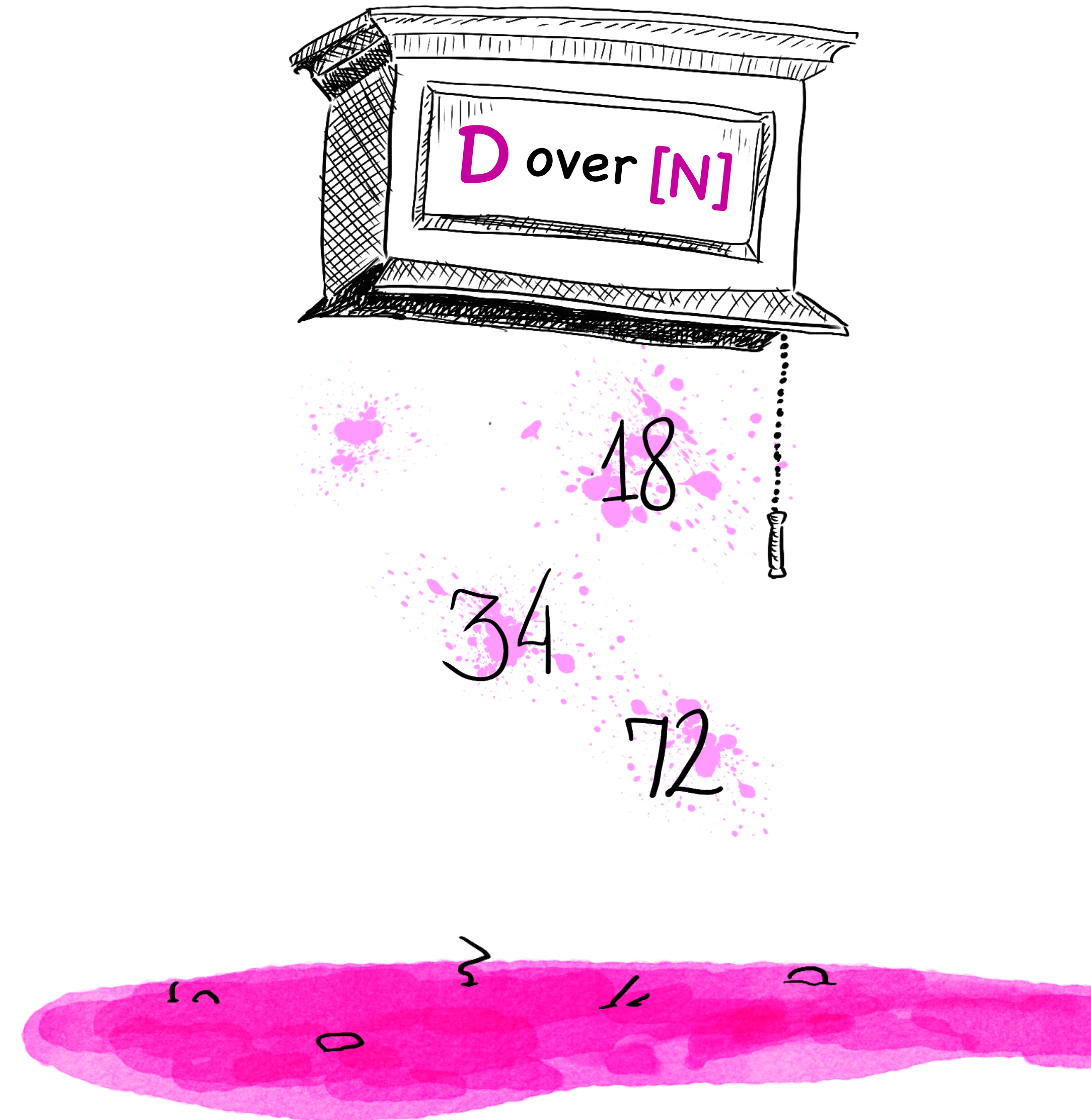
Communication:
 $o(N)$, small # rounds

V's Samples, runtime:
 $o(\text{optimal } \Pi \text{ tester})$

*Quantities of
interest*

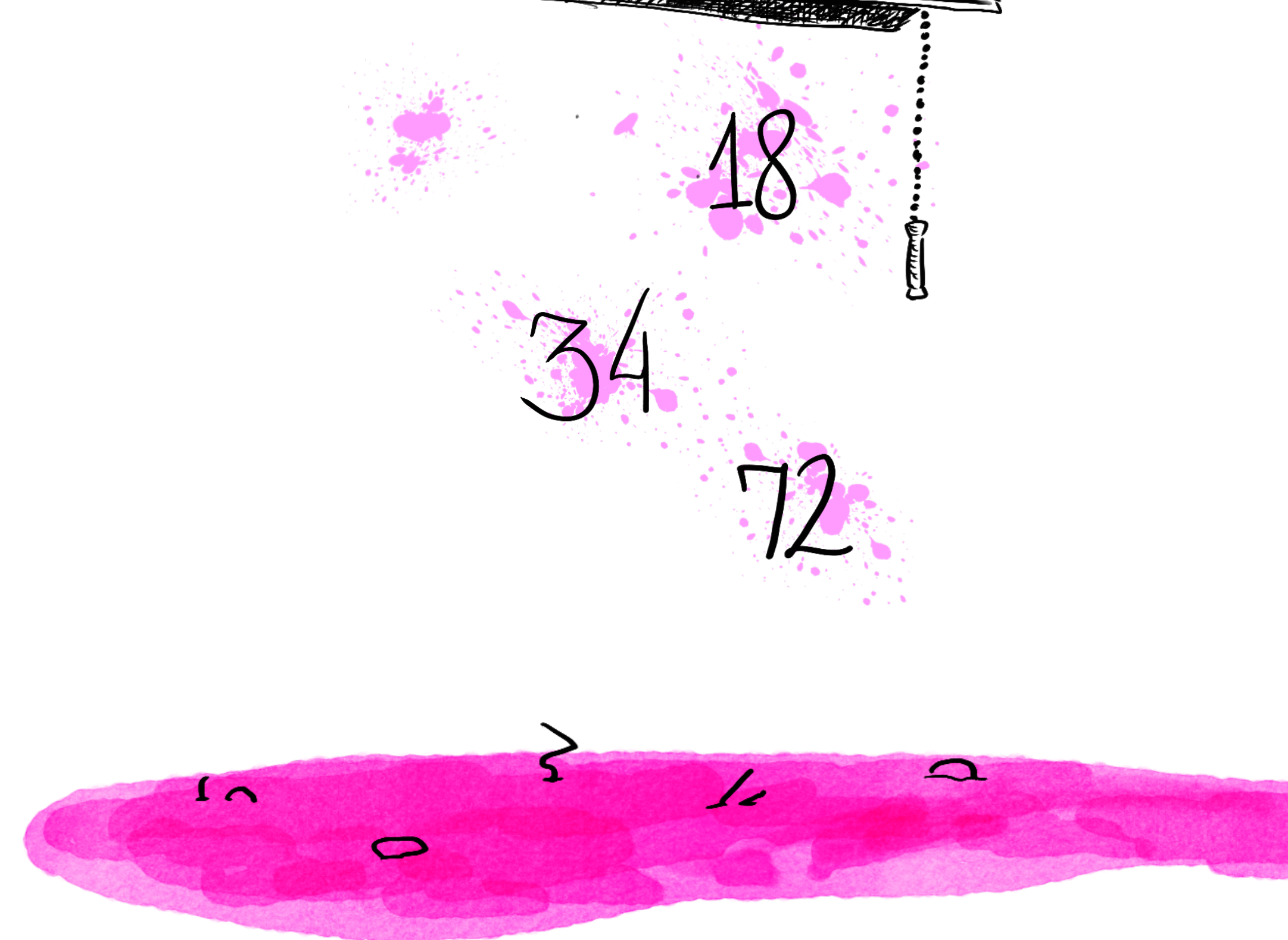
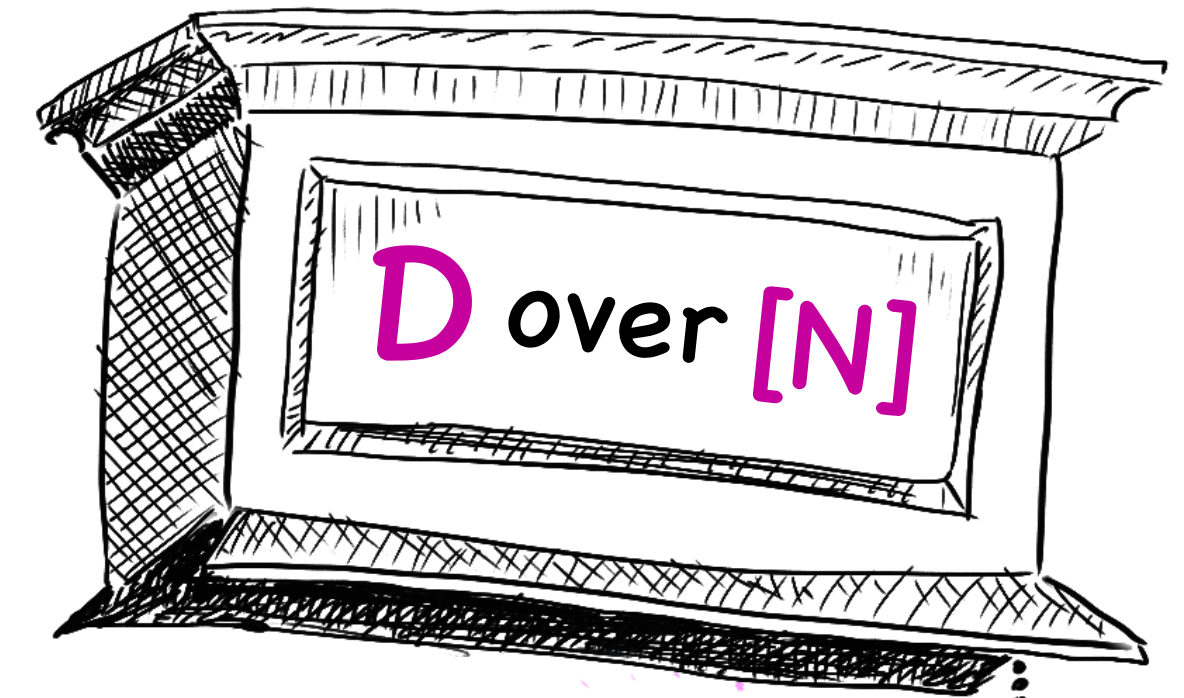


Example: Verifying Uniformity Over Half Domain



Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.

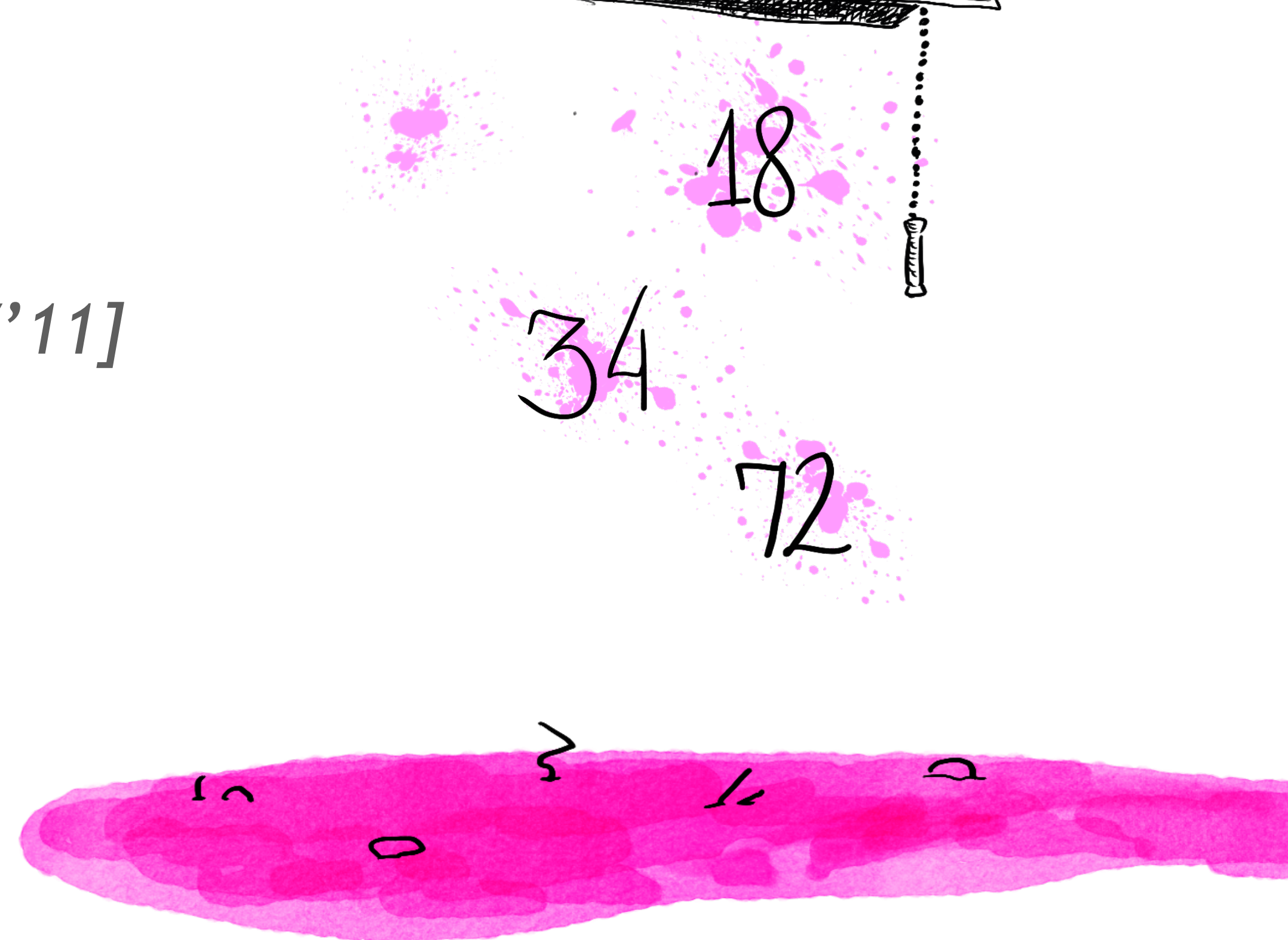
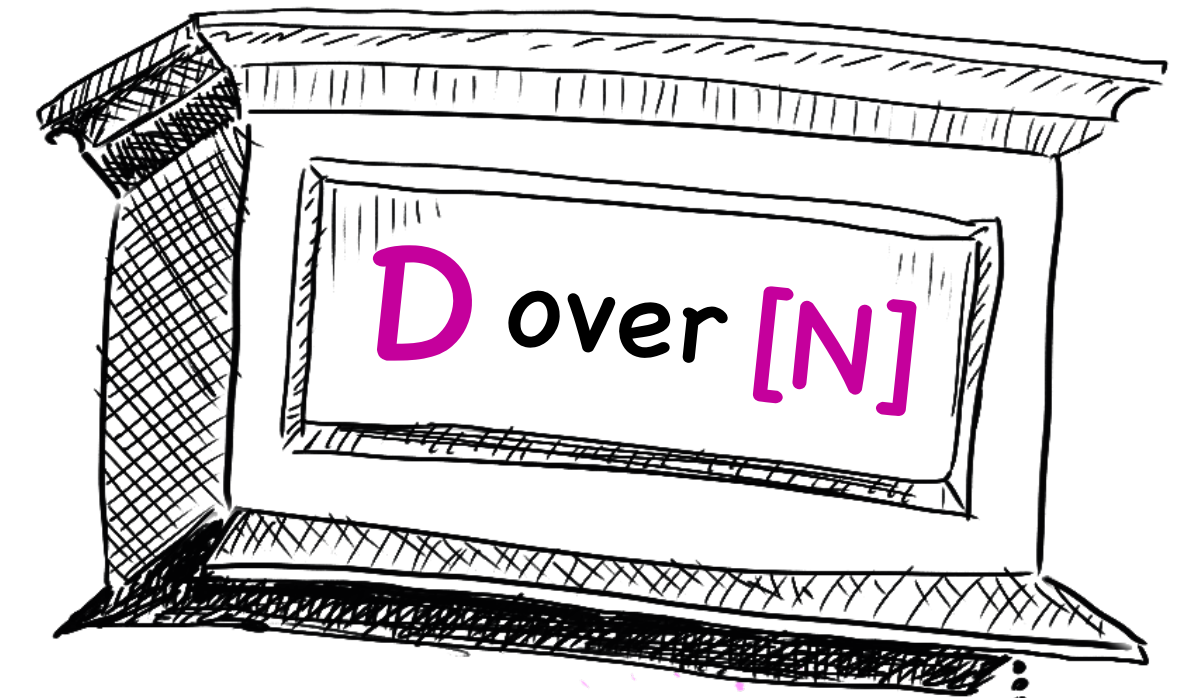


Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over *some half* of their domain.

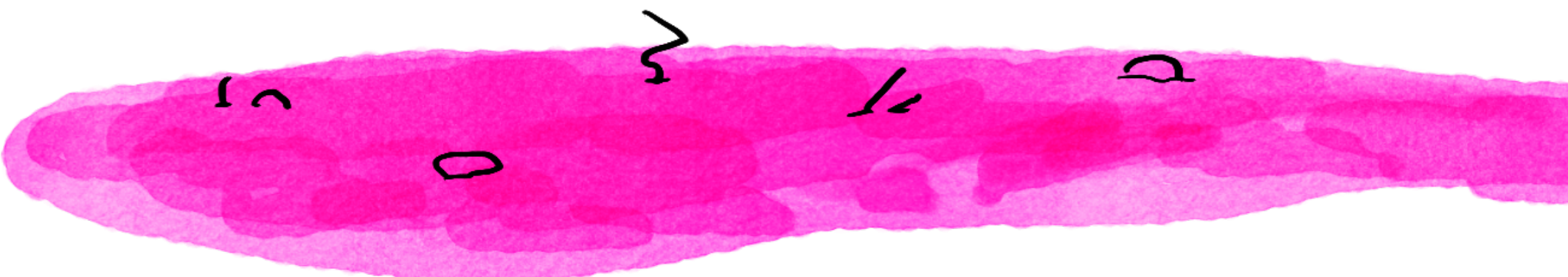
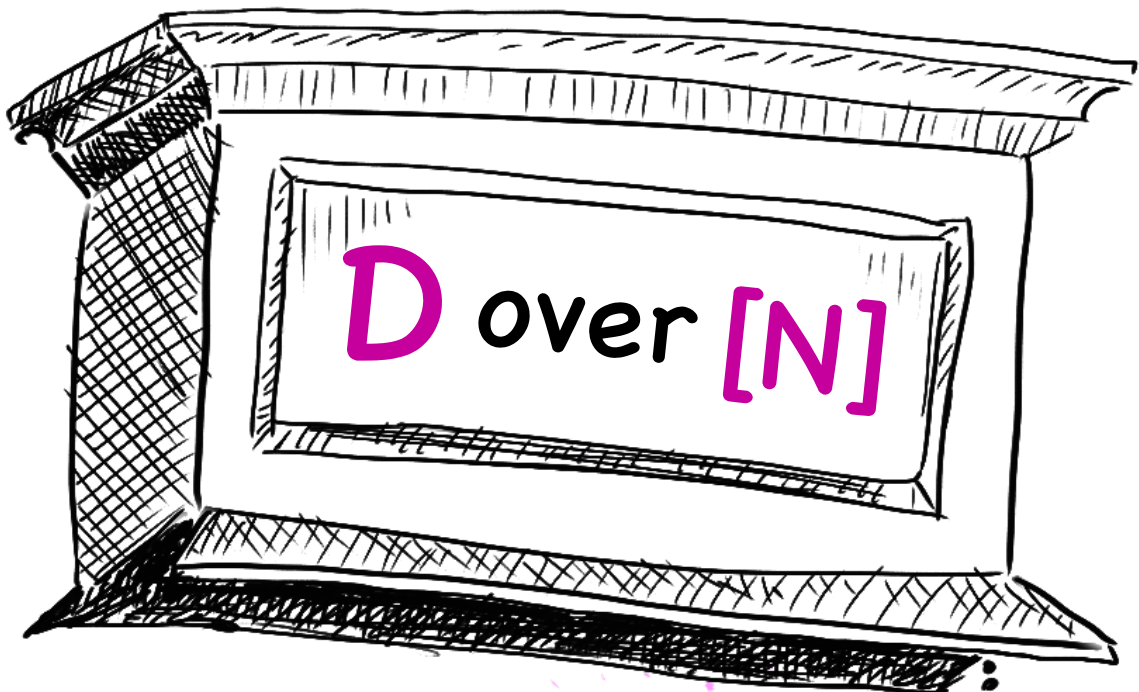
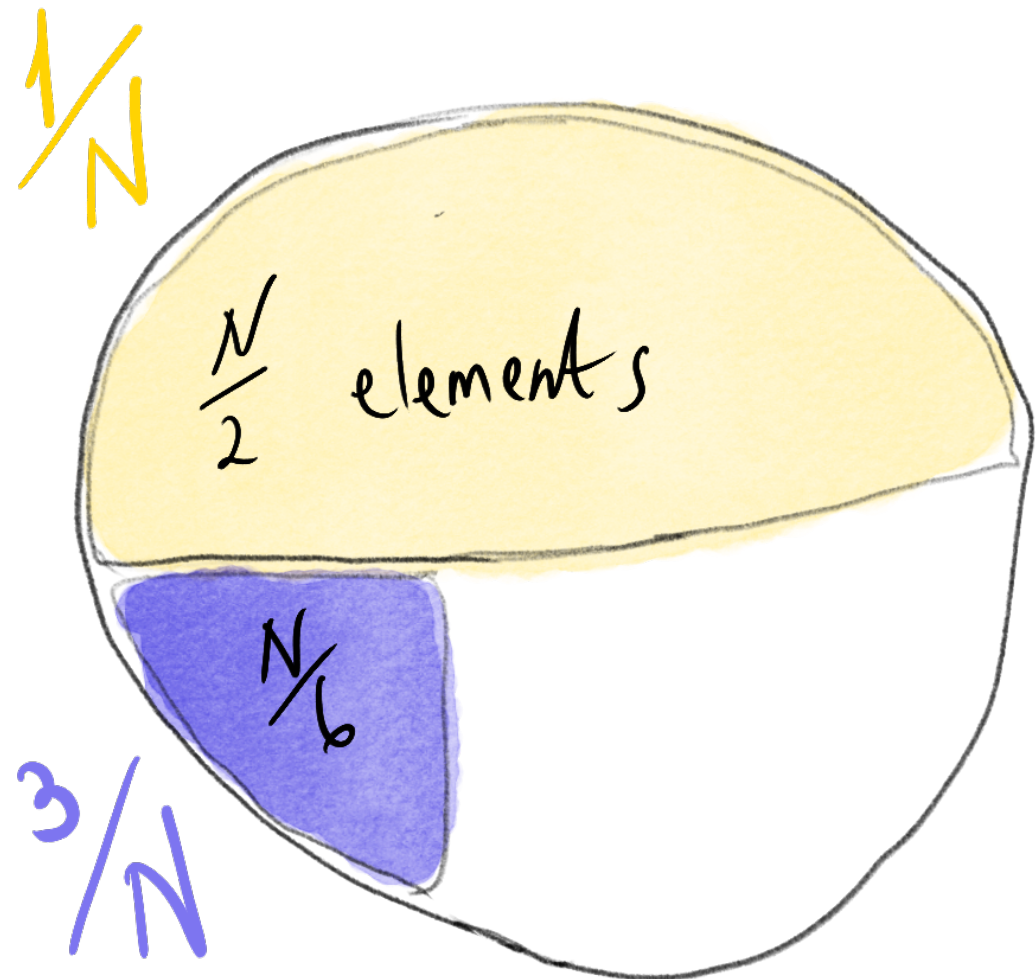
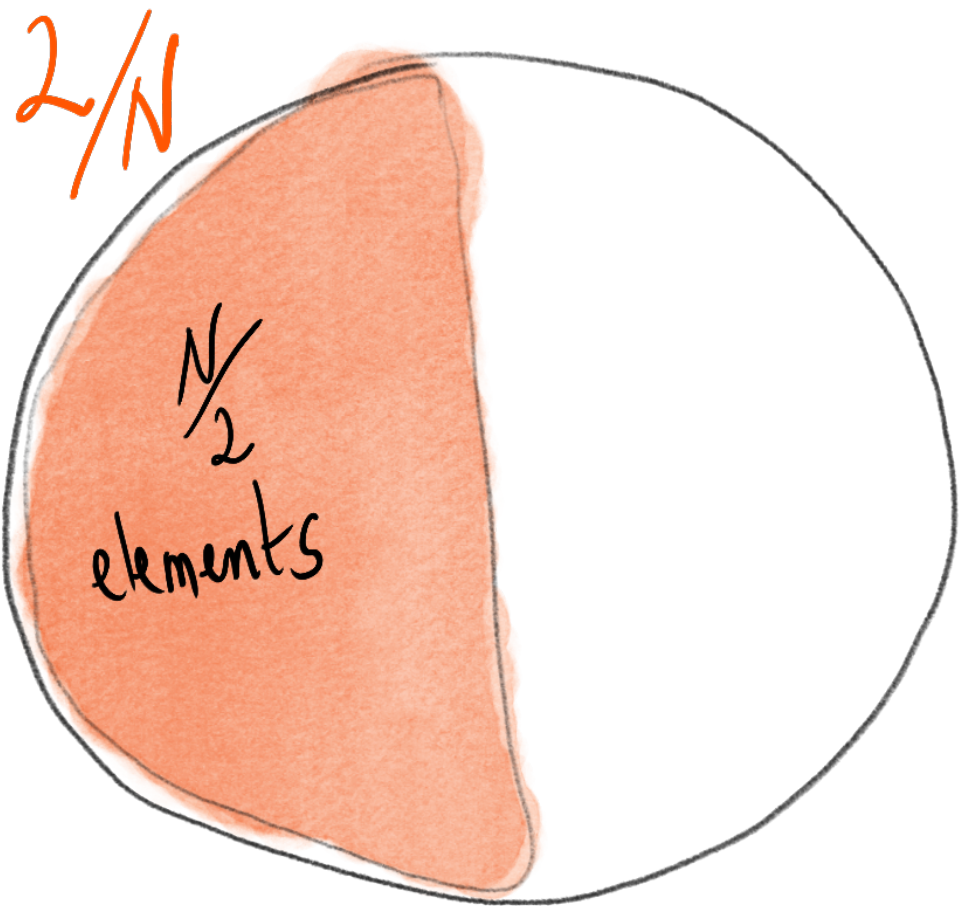
Testing the property:

- $\Theta(N^{2/3})$ samples. [BC'17]
- (Tolerantly) $\Theta(N/\log N)$ samples. [VV'11]



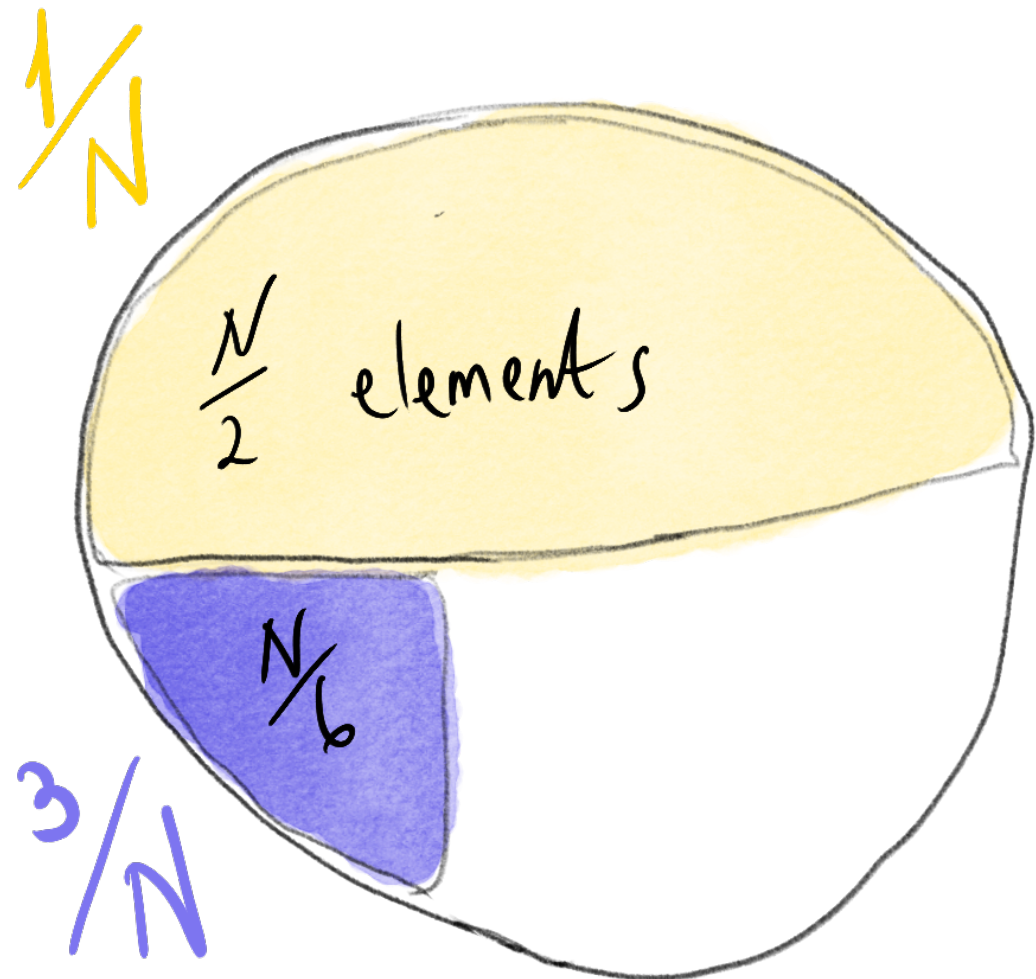
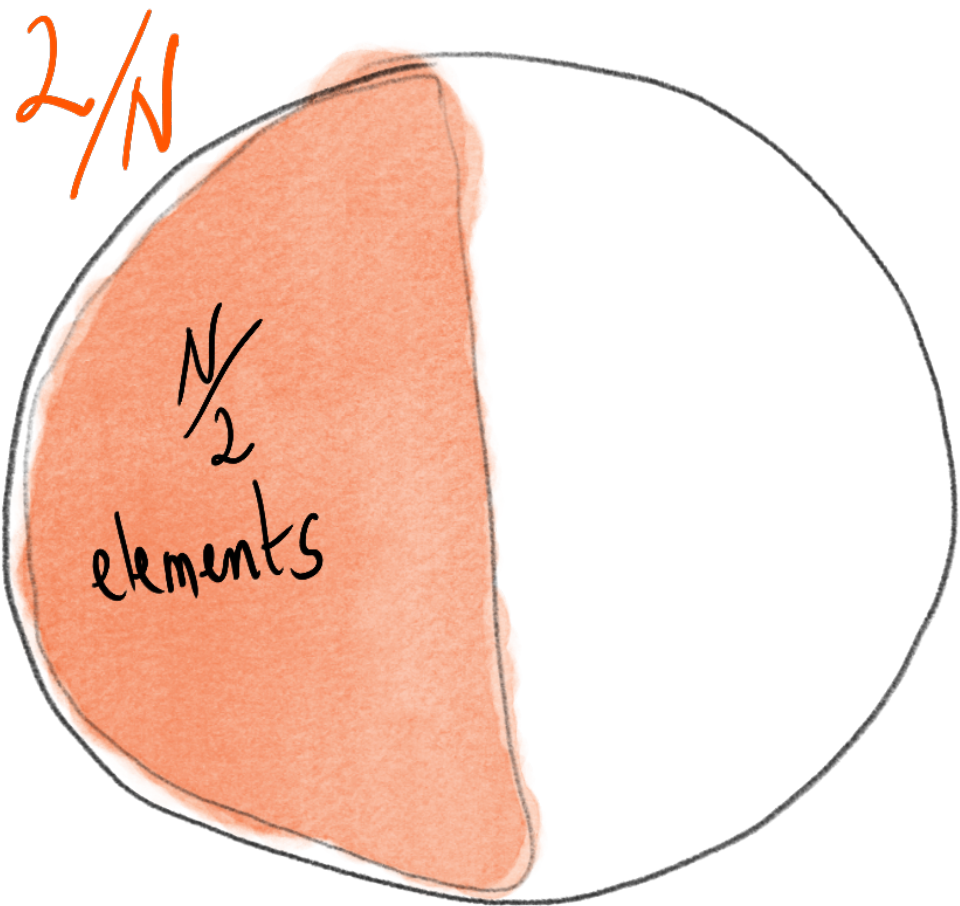
Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.

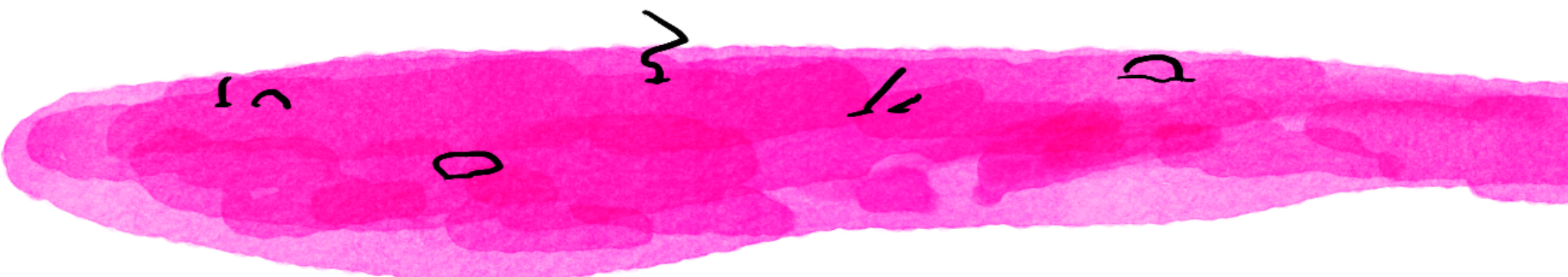
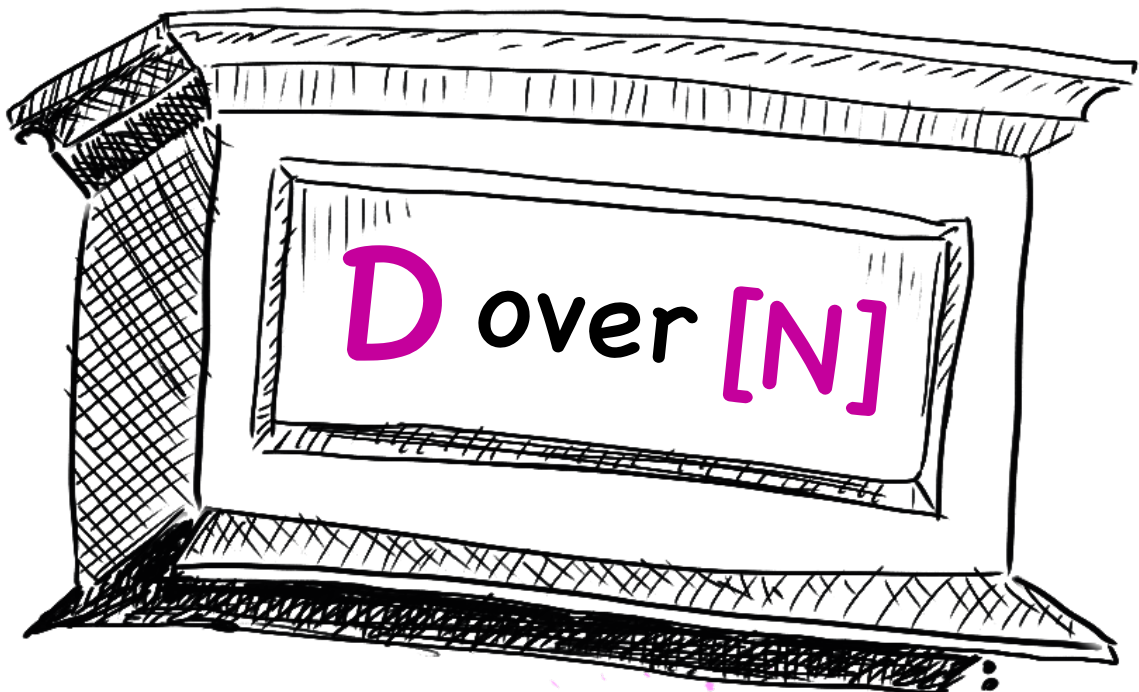


Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.

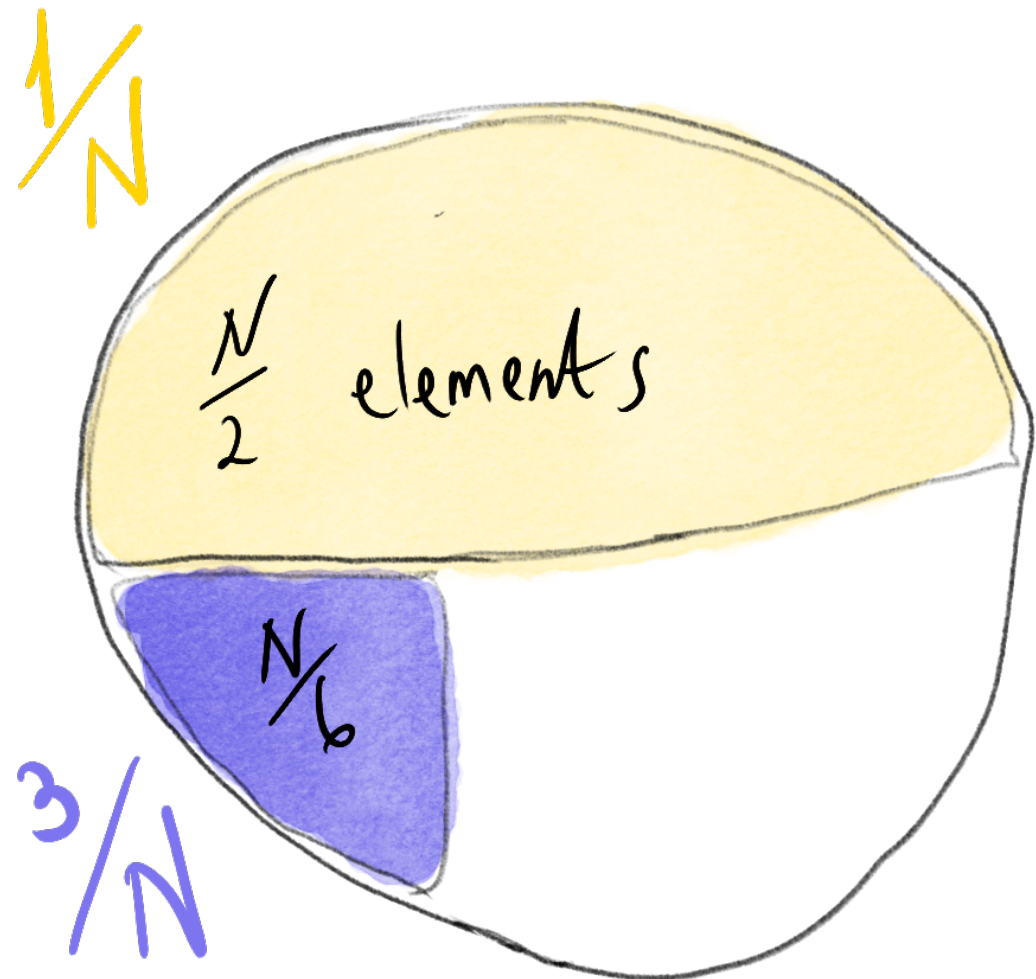
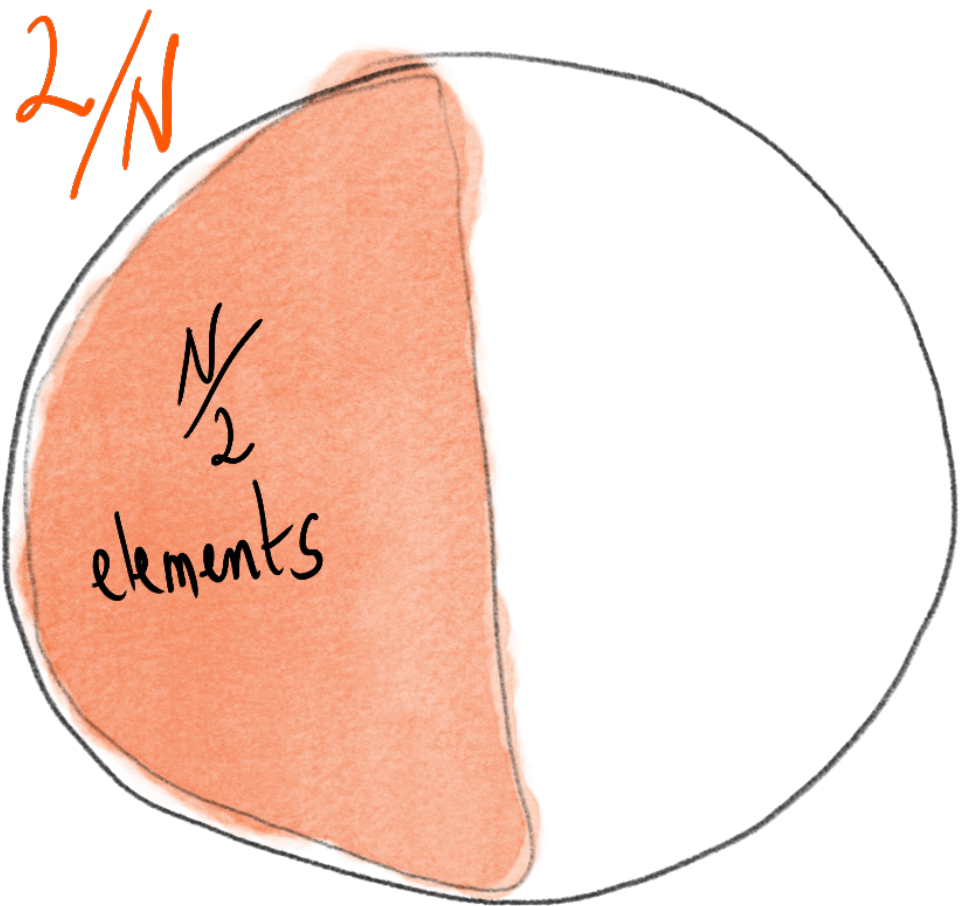


Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x)$

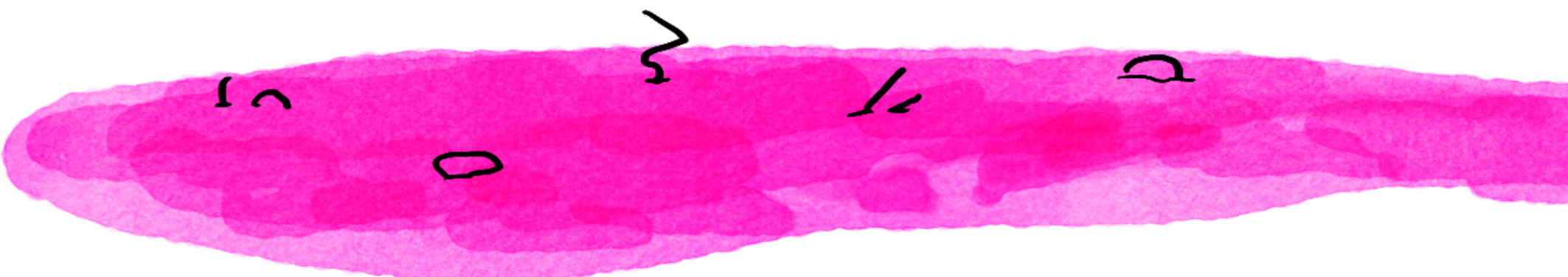
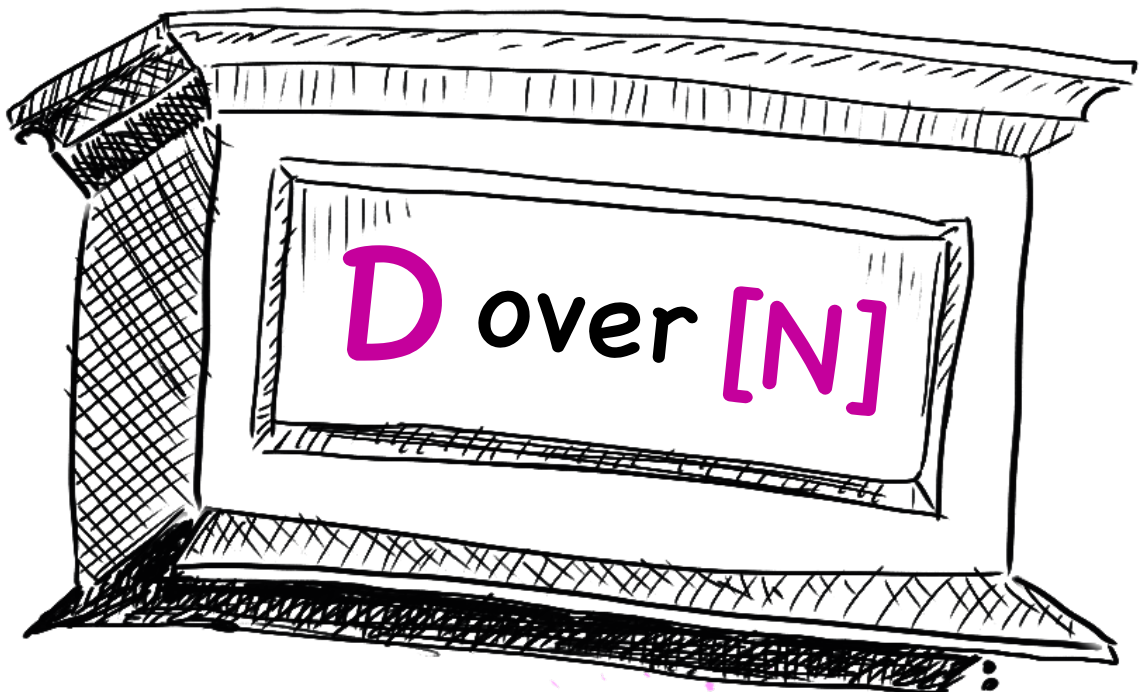


Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.

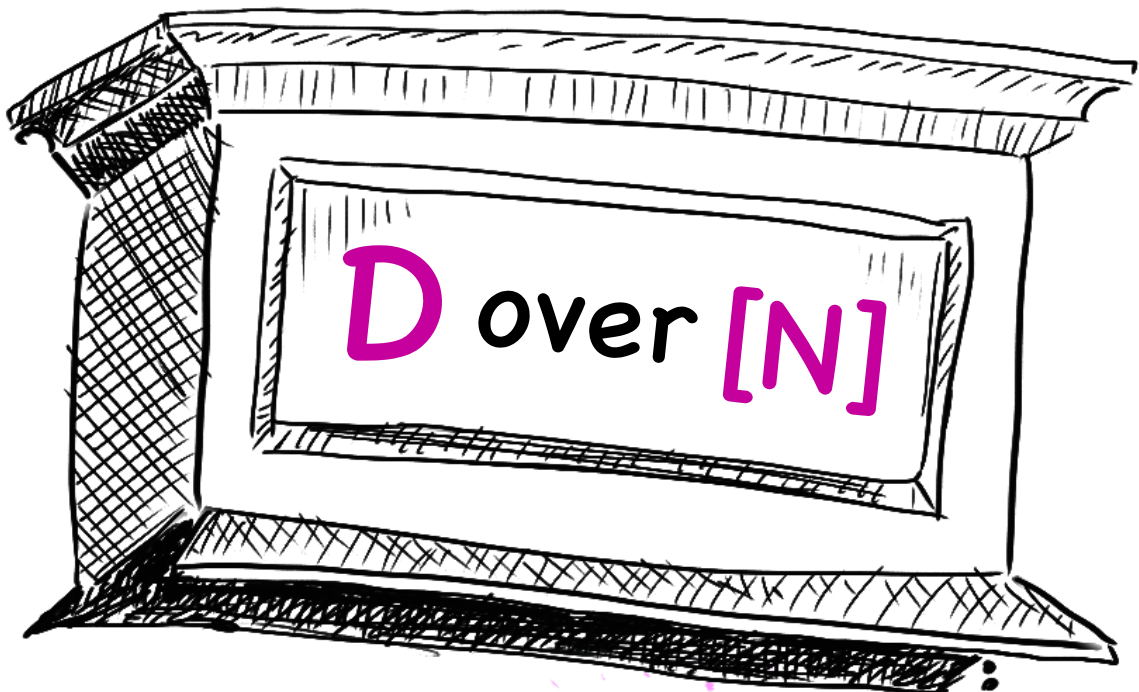
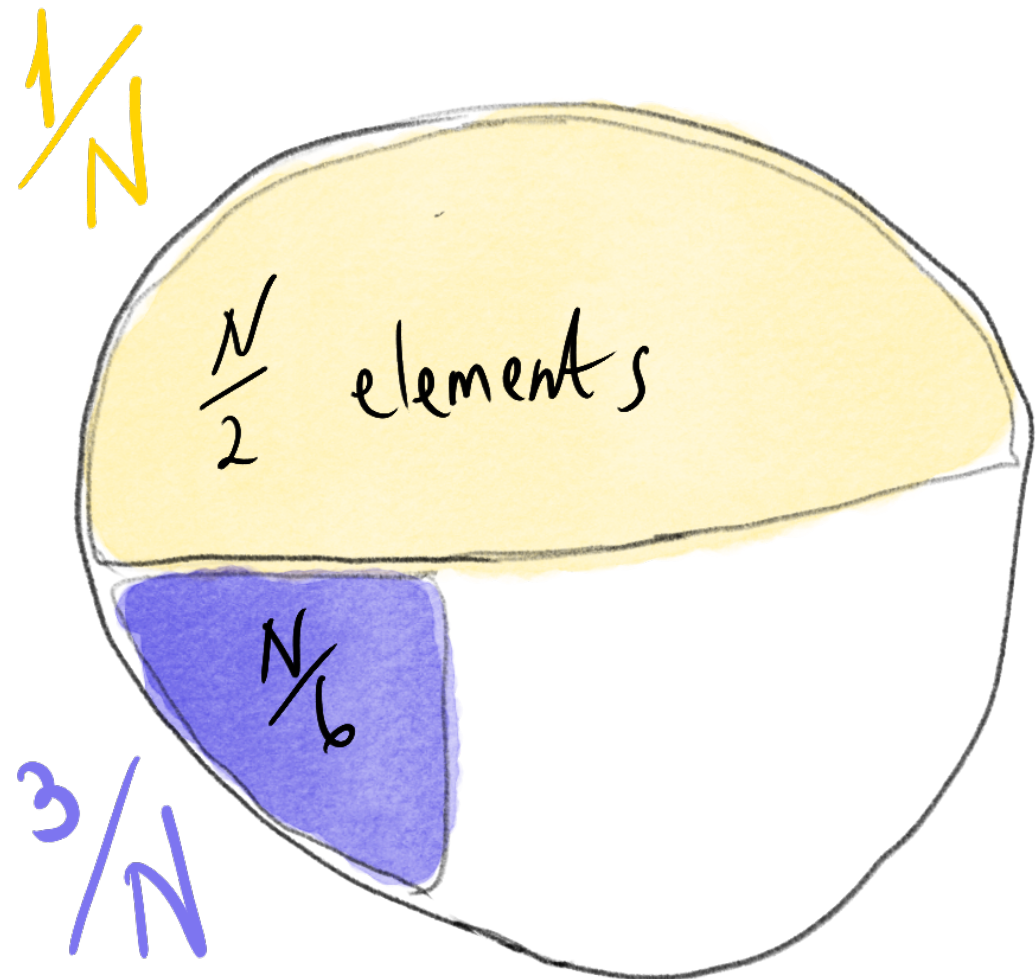
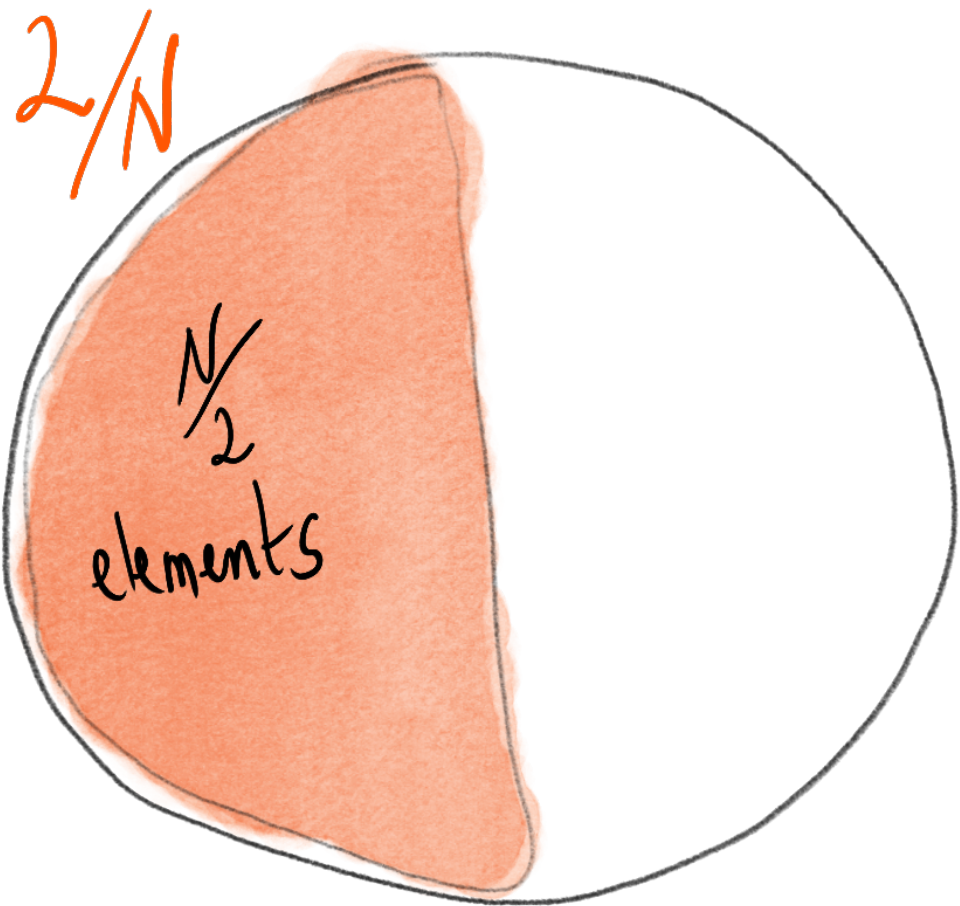


Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x) = \frac{2}{N}$



Example: Verifying Uniformity Over Half Domain

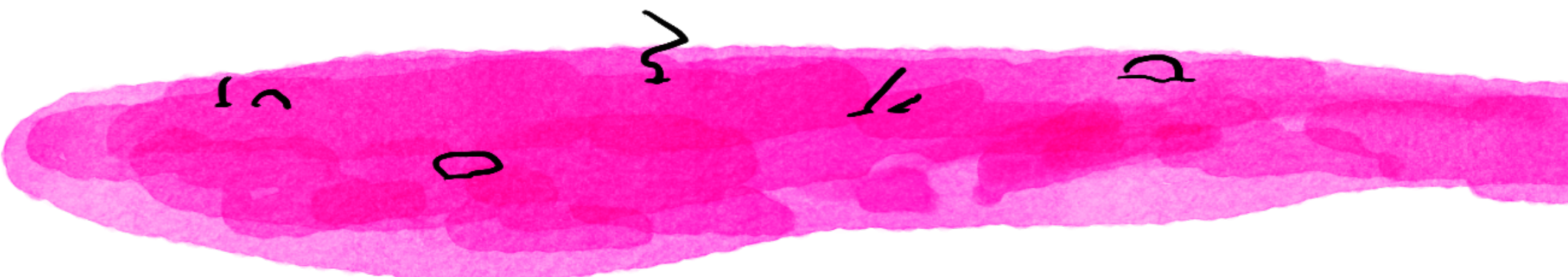
U_{half} All distributions uniform over some half of their domain.



18
34
72

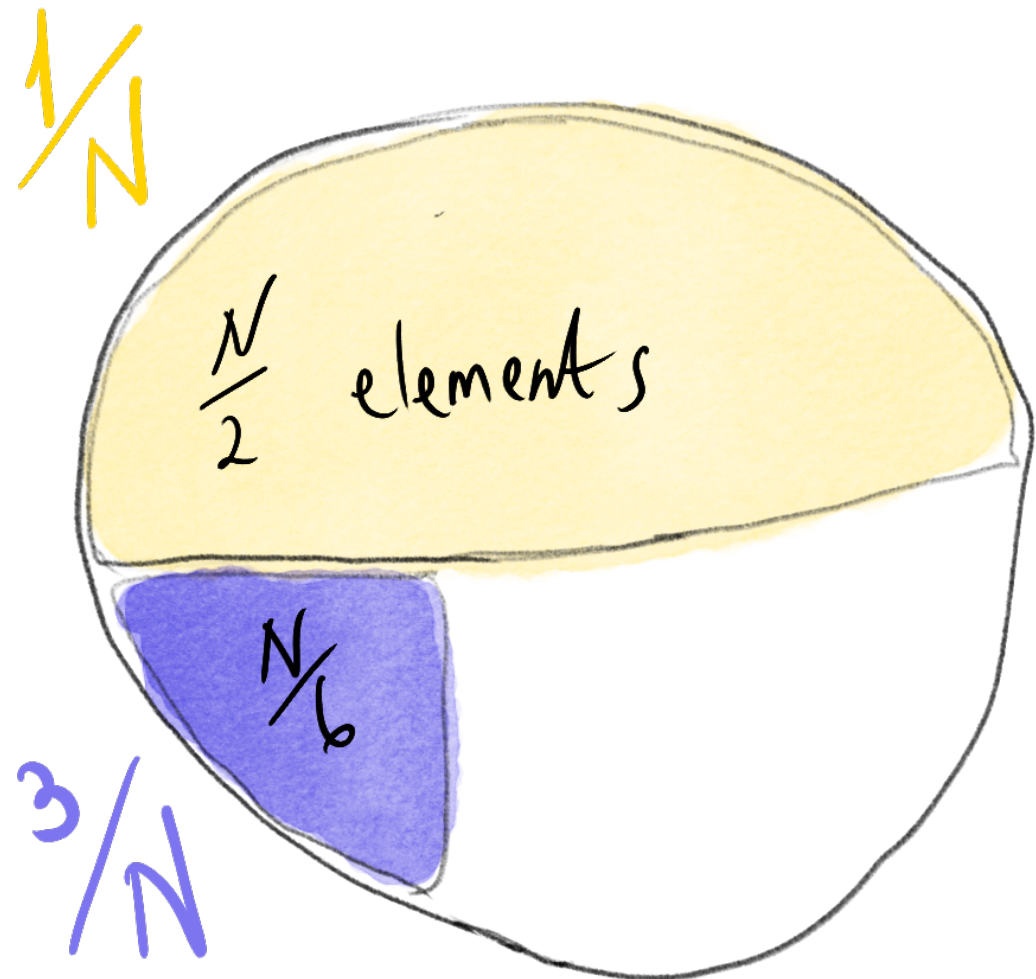
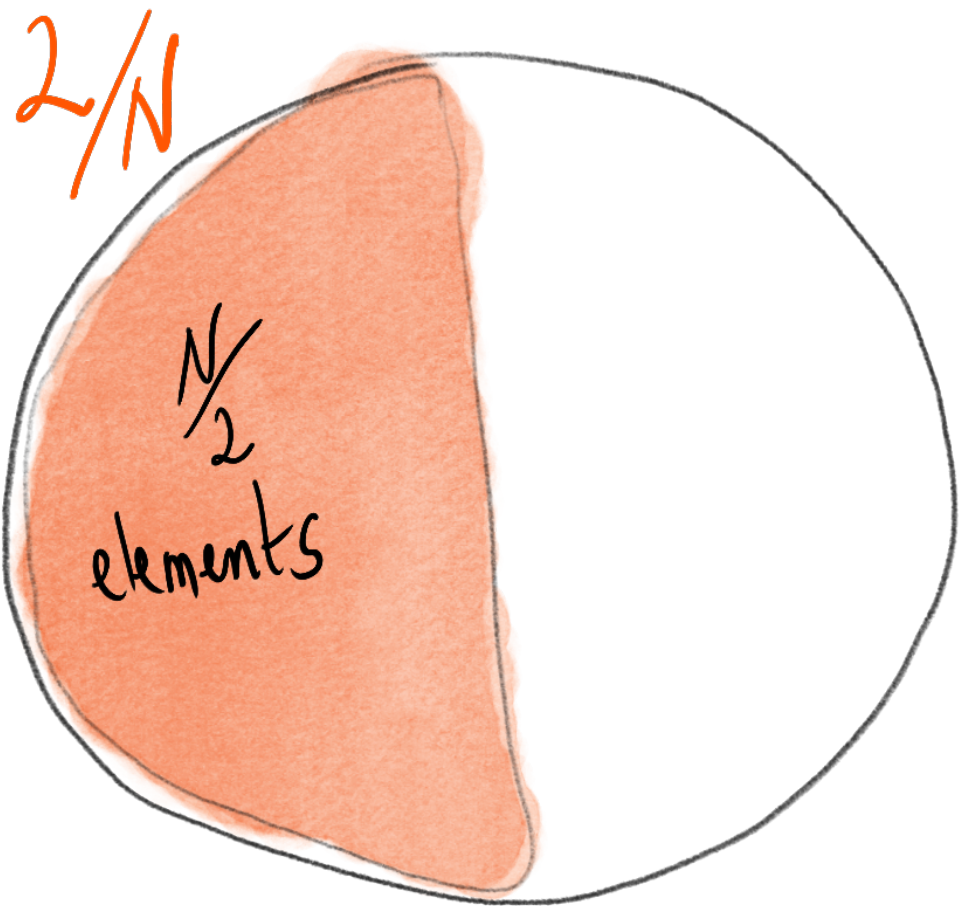
Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x) = \frac{2}{N}$

Can't test with $o(N^{2/3})$



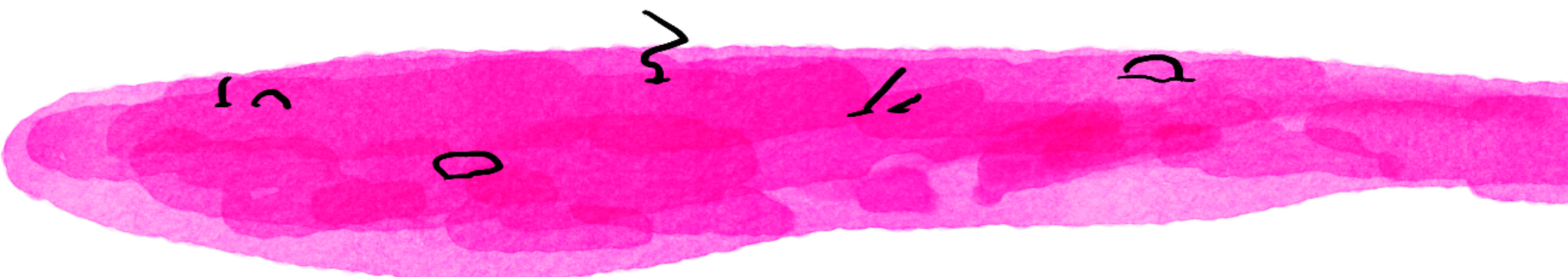
Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.



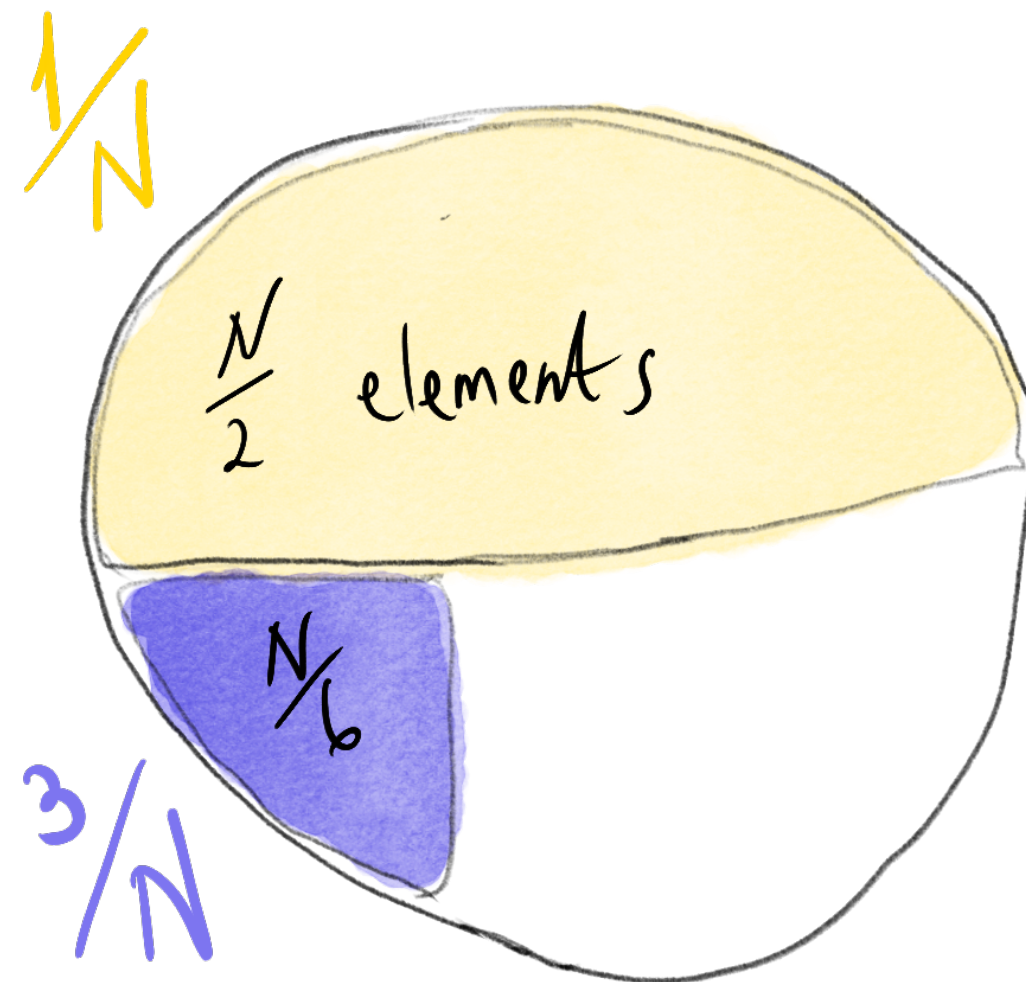
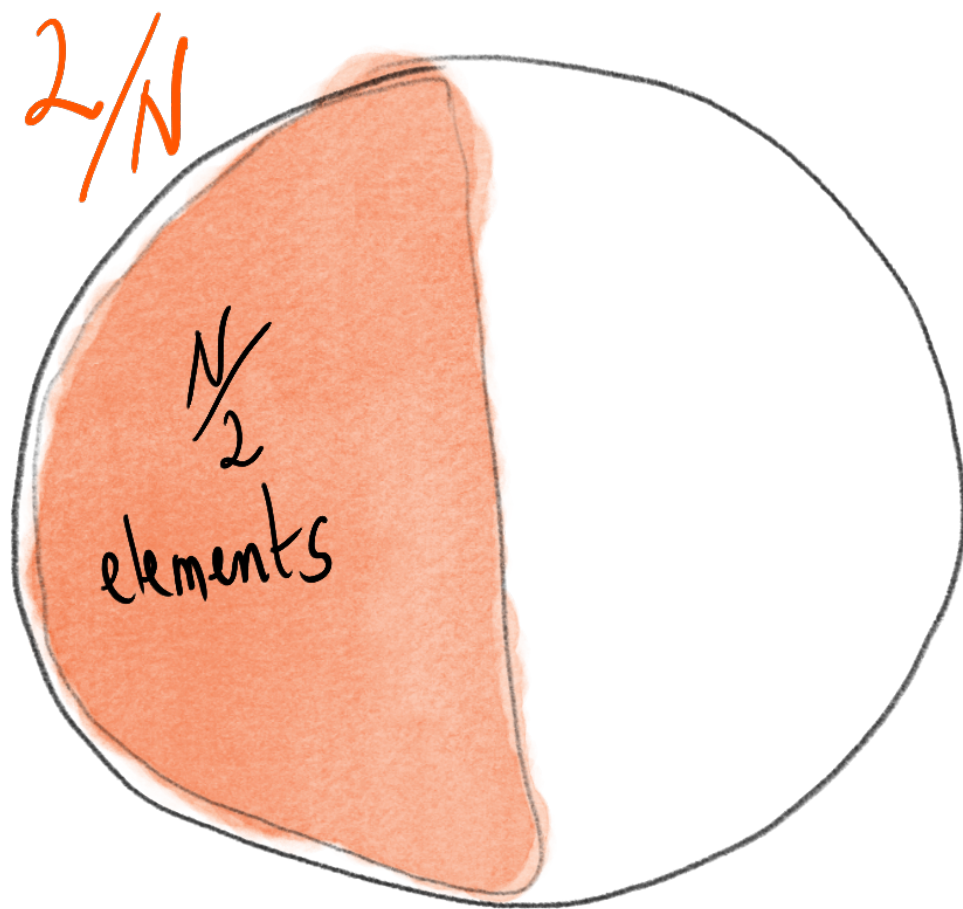
Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x) = \frac{2}{N}$

With $O(N^{2/3})$ samples exhibit 3-collisions & decide!



Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.

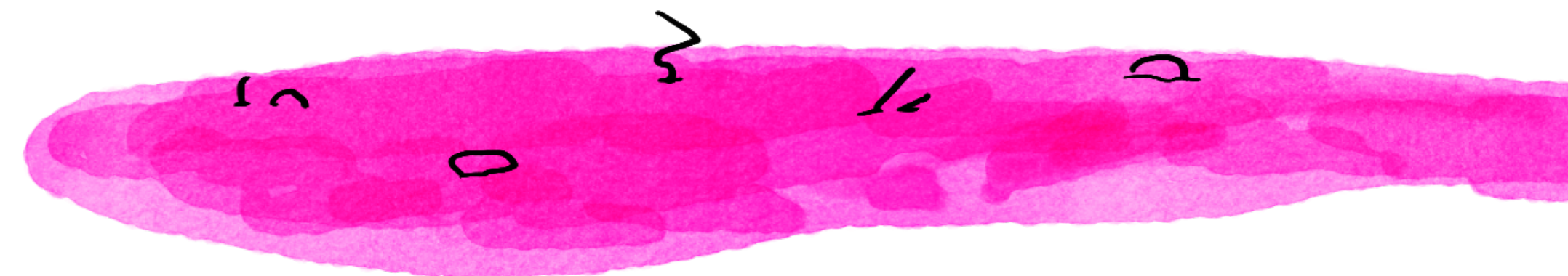


Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x) = \frac{2}{N}$

With $O(N^{2/3})$ samples exhibit 3-collisions & decide!

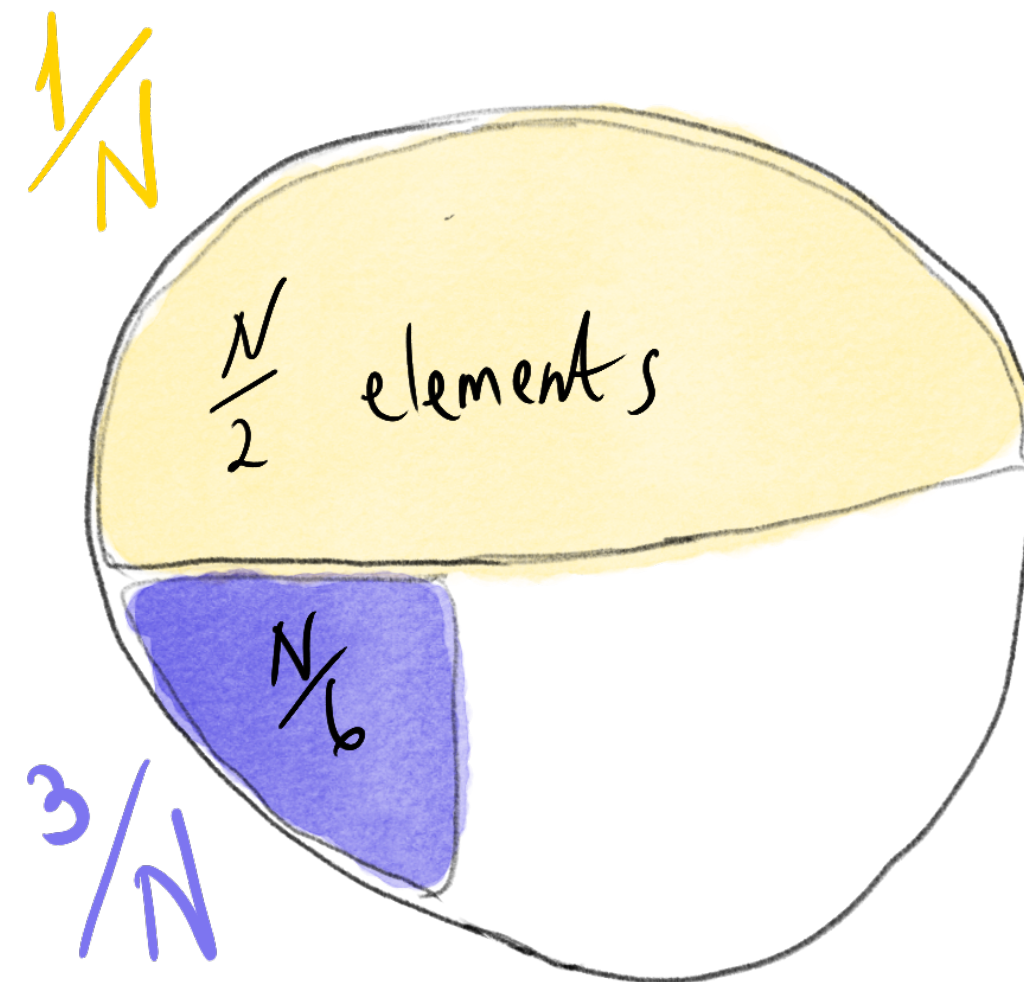
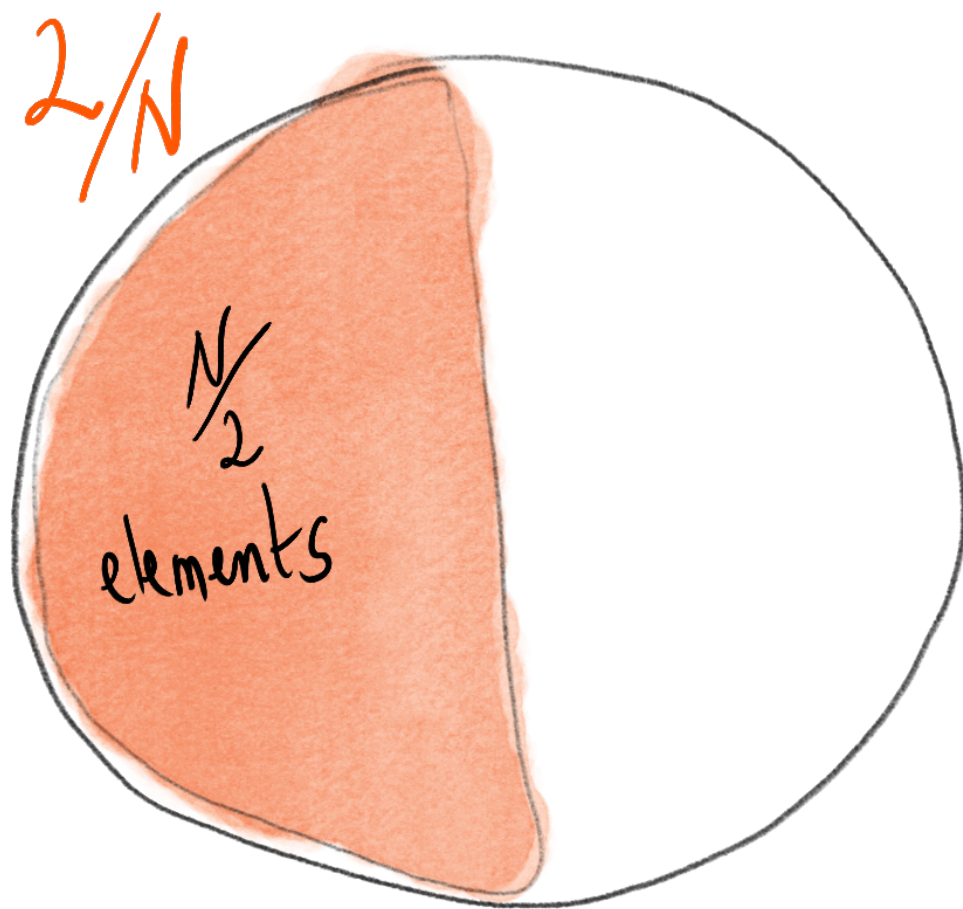
Observation:

Support sizes are different.



Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.



Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x) = \frac{2}{N}$

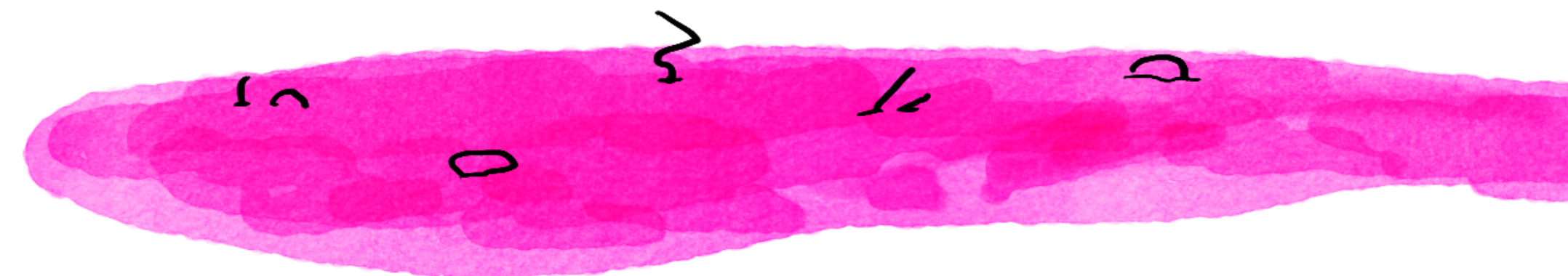
With $O(N^{2/3})$ samples exhibit 3-collisions & decide!

Observation:

Support sizes are different.

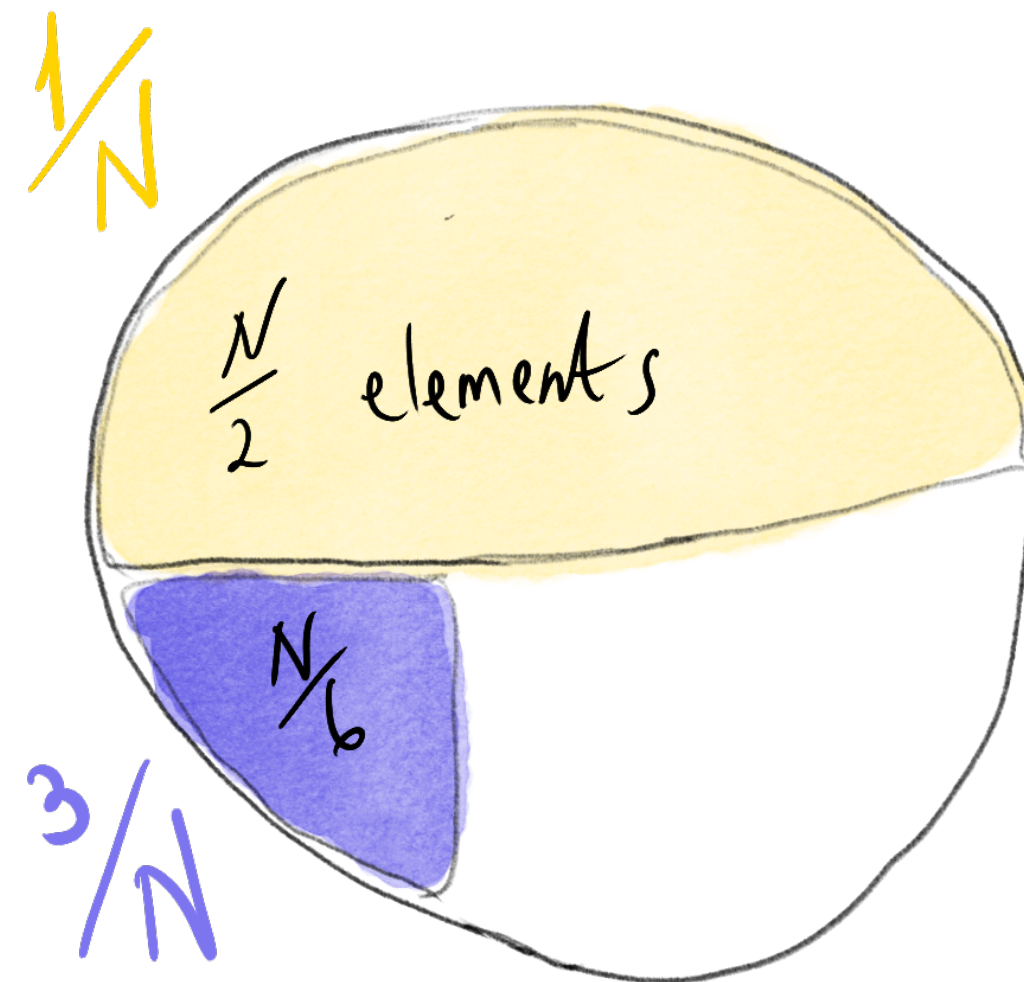
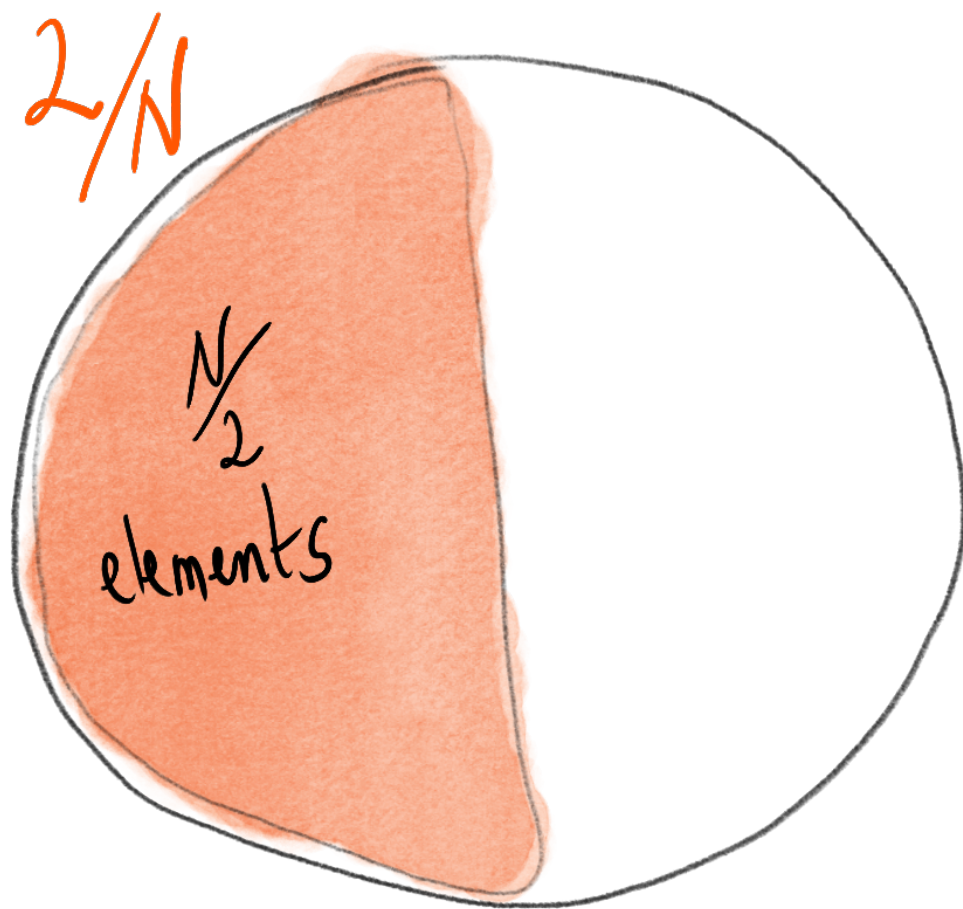
$$\sum_x D^2(x) = \frac{2}{N} \text{ AND } \delta_{TV}(D, U_{half}) \geq \epsilon$$

$$|\text{Supp}(D)| \geq \frac{N}{2} (1 + \Theta(\epsilon^2))$$



Example: Verifying Uniformity Over Half Domain

U_{half} All distributions uniform over some half of their domain.



Probability of seeing a 2-collision : $\sum_x D(x) \cdot D(x) = \frac{2}{N}$

With $O(N^{2/3})$ samples exhibit 3-collisions & decide!

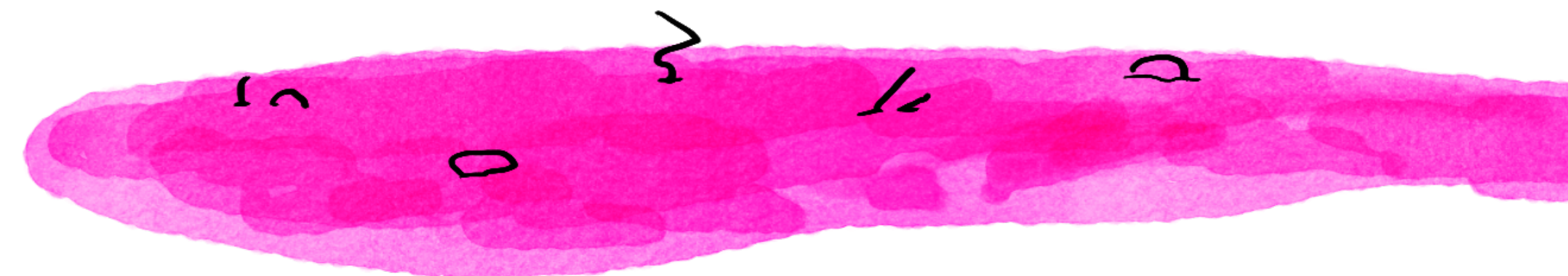
Observation:

Support sizes are different.

Every NO-case either has:

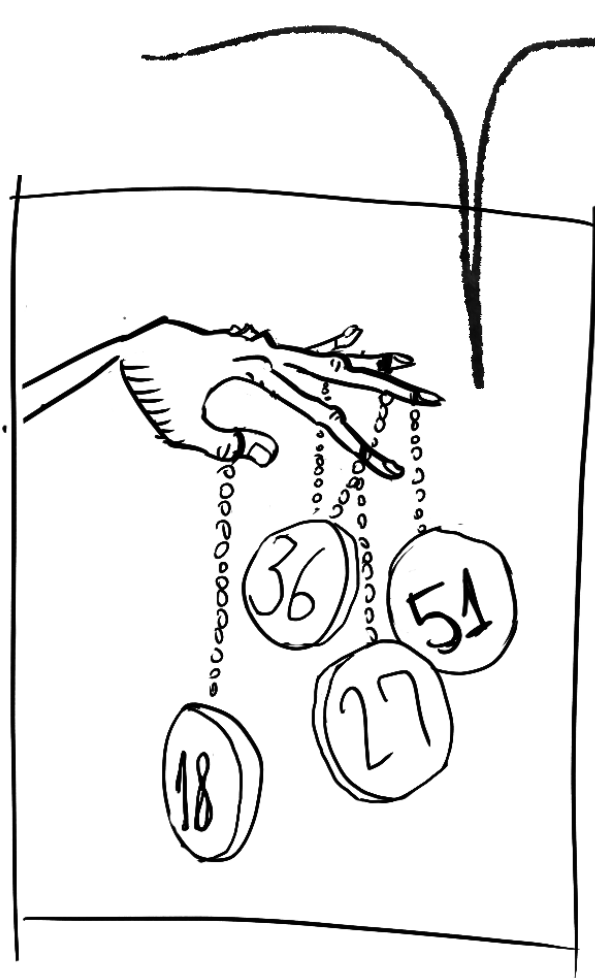
- 2-collision prob. $\neq 2/N$;
- or, support larger than.

$$\frac{N}{2} \left(1 + \Theta(\epsilon^2) \right)$$

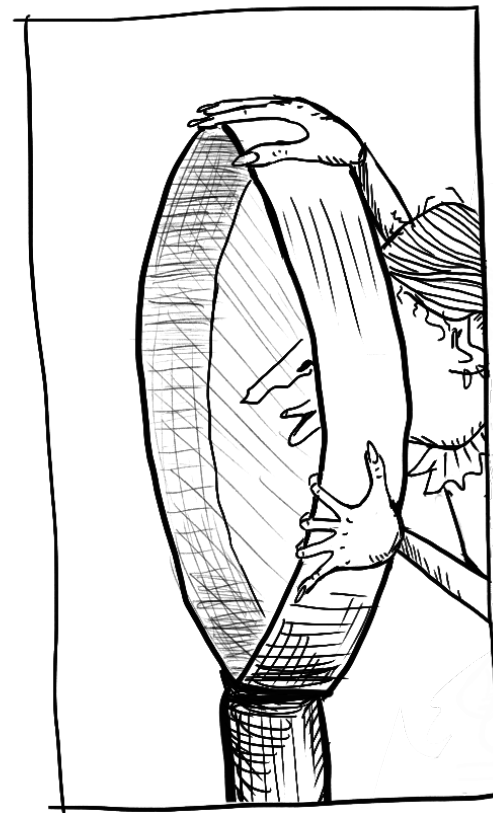


Verifying Uniformity Over Half Domain

$$\mathbf{D} \in \mathbf{U}_{N/2}$$



Prover

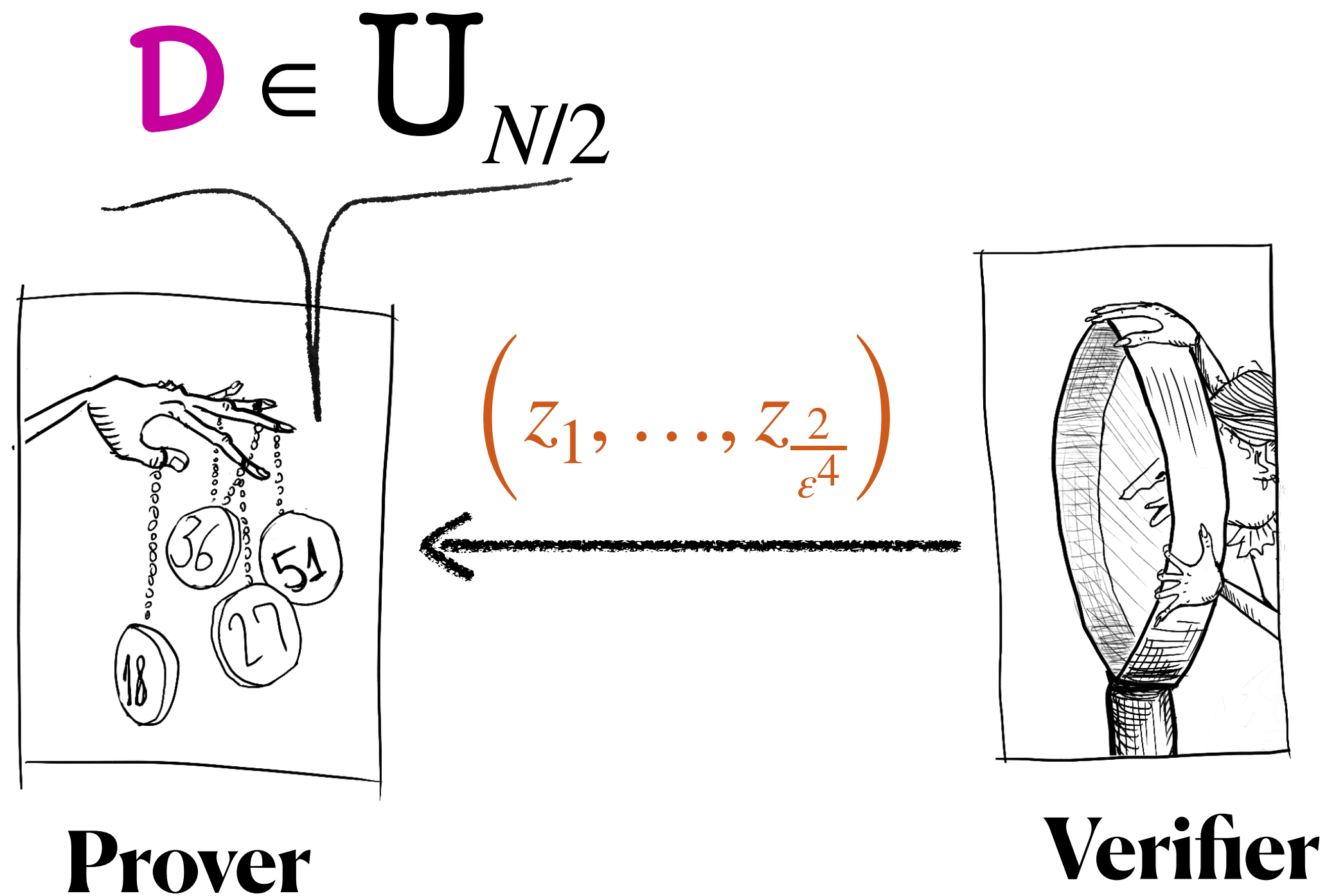


Verifier

1. Draw $s = O(\sqrt{N})$ samples.

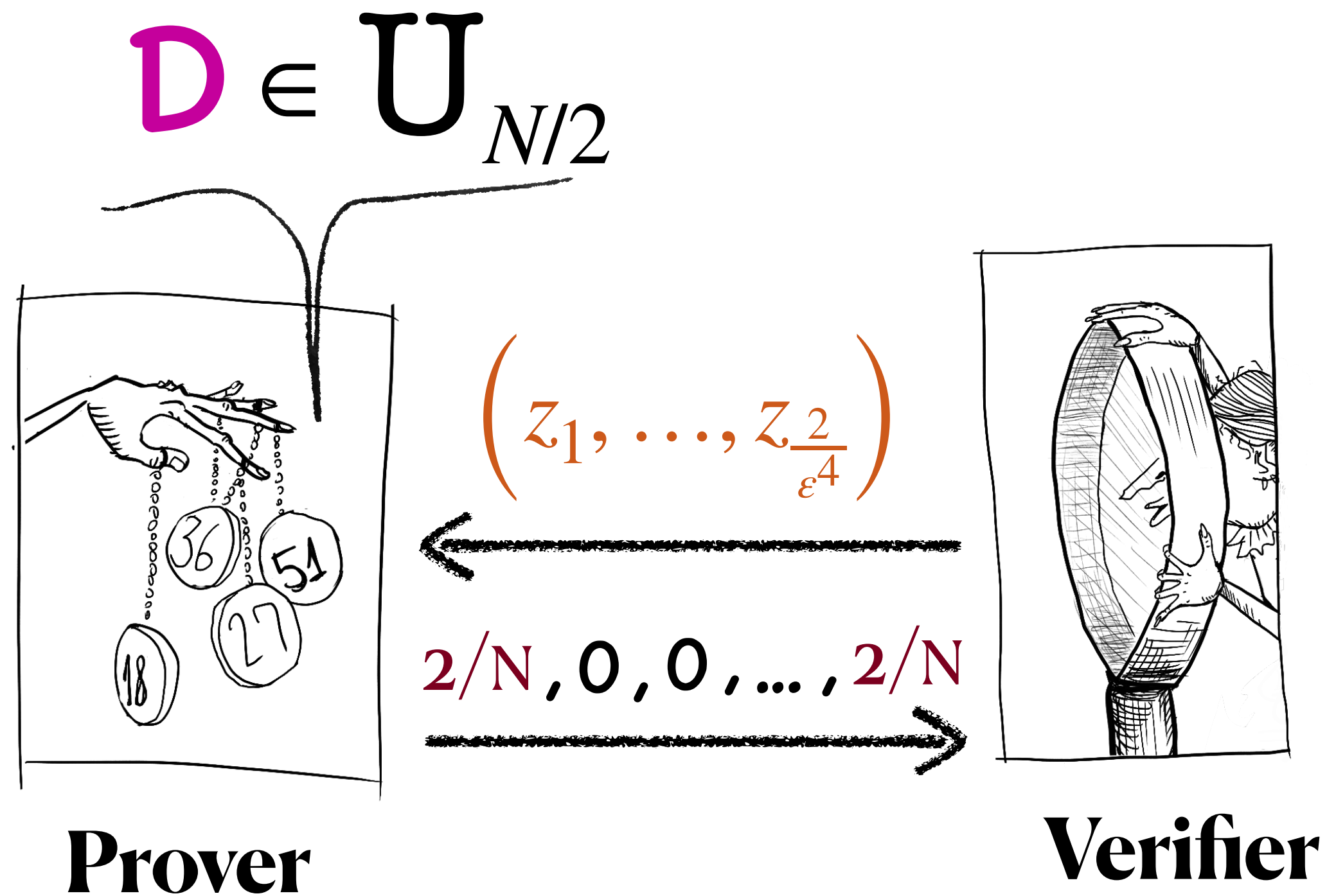
Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$

Verifying Uniformity Over Half Domain



1. Draw $s = O\left(\sqrt{N}\right)$ samples.
Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$
2. Draw ϵ^{-4} samples from $U_{[N]}$.
Draw ϵ^{-4} samples from D .
Randomly mix: $\left(z_1, \dots, z_{\frac{2}{\epsilon^4}}\right)$

Verifying Uniformity Over Half Domain



1. Draw $s = O(\sqrt{N})$ samples.

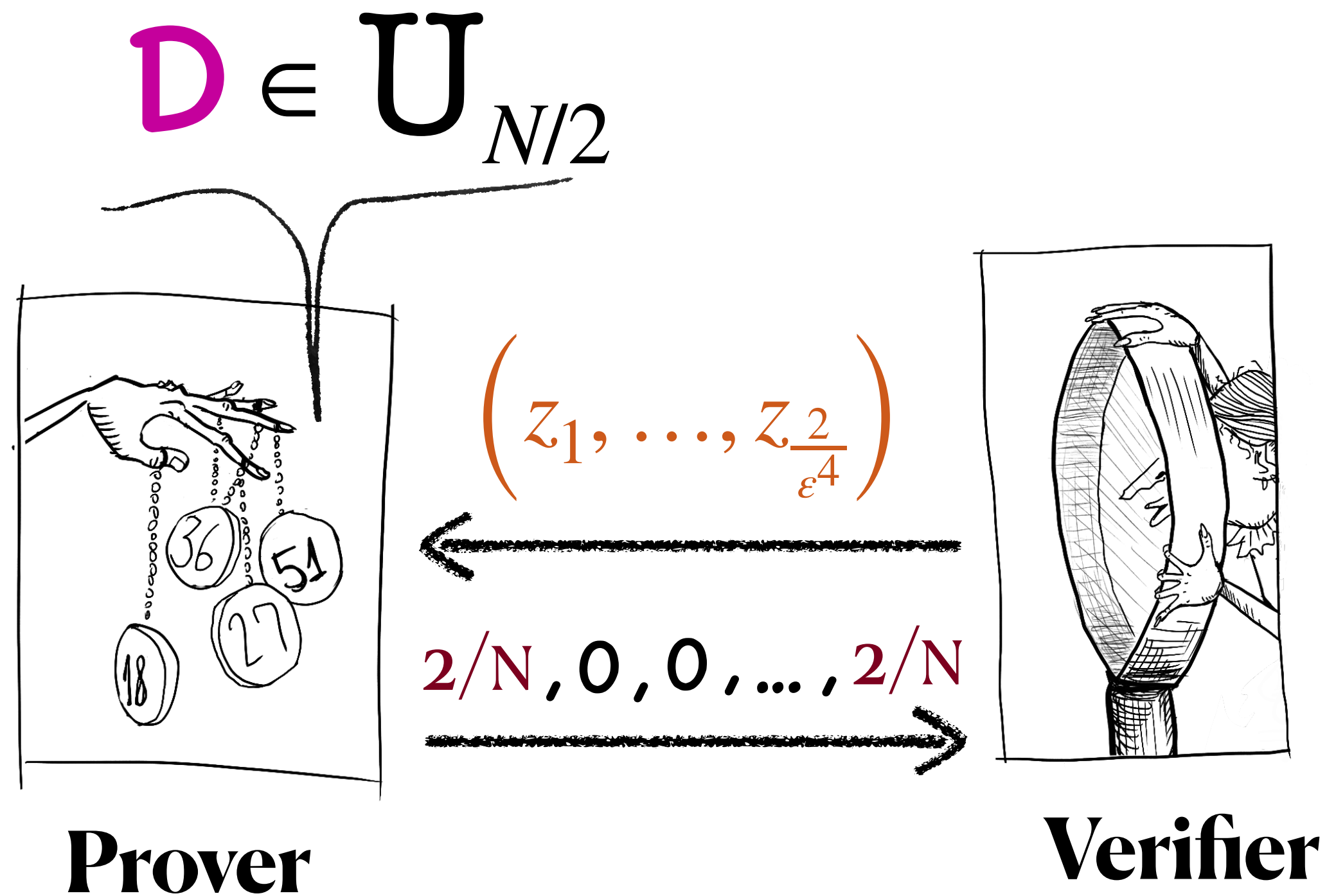
Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$

2. Draw ϵ^{-4} samples from $\mathcal{U}_{[N]}$.

Draw ϵ^{-4} samples from D .

Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$

Verifying Uniformity Over Half Domain



1. Draw $s = O(\sqrt{N})$ samples.

Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$

2. Draw ϵ^{-4} samples from $U_{[N]}$.

Draw ϵ^{-4} samples from D .

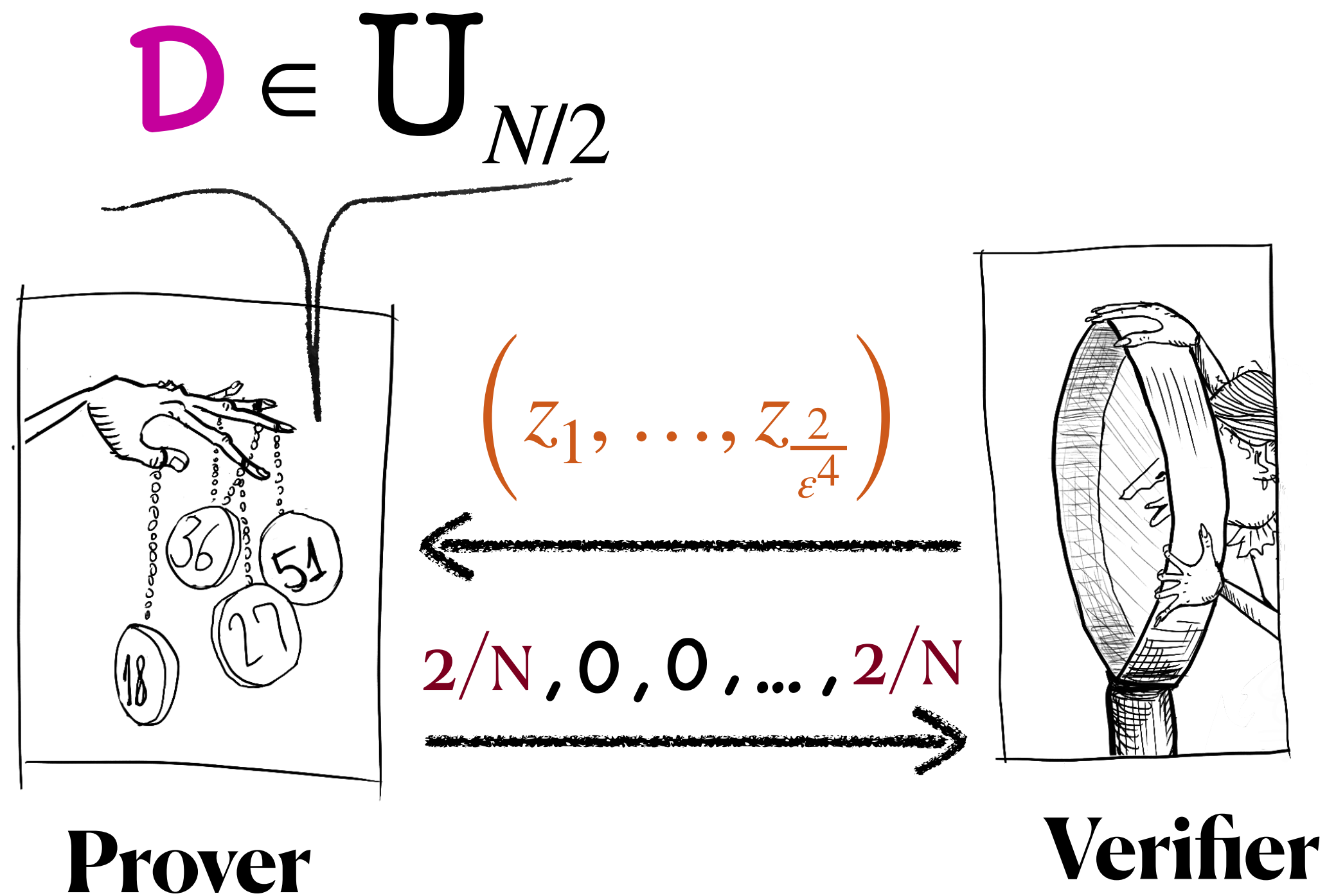
Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$

Test 2:

$\approx$ *half* the $U_{[N]}$ -samples have tag $2/N$

ALL D -samples have tag $2/N$

Verifying Uniformity Over Half Domain



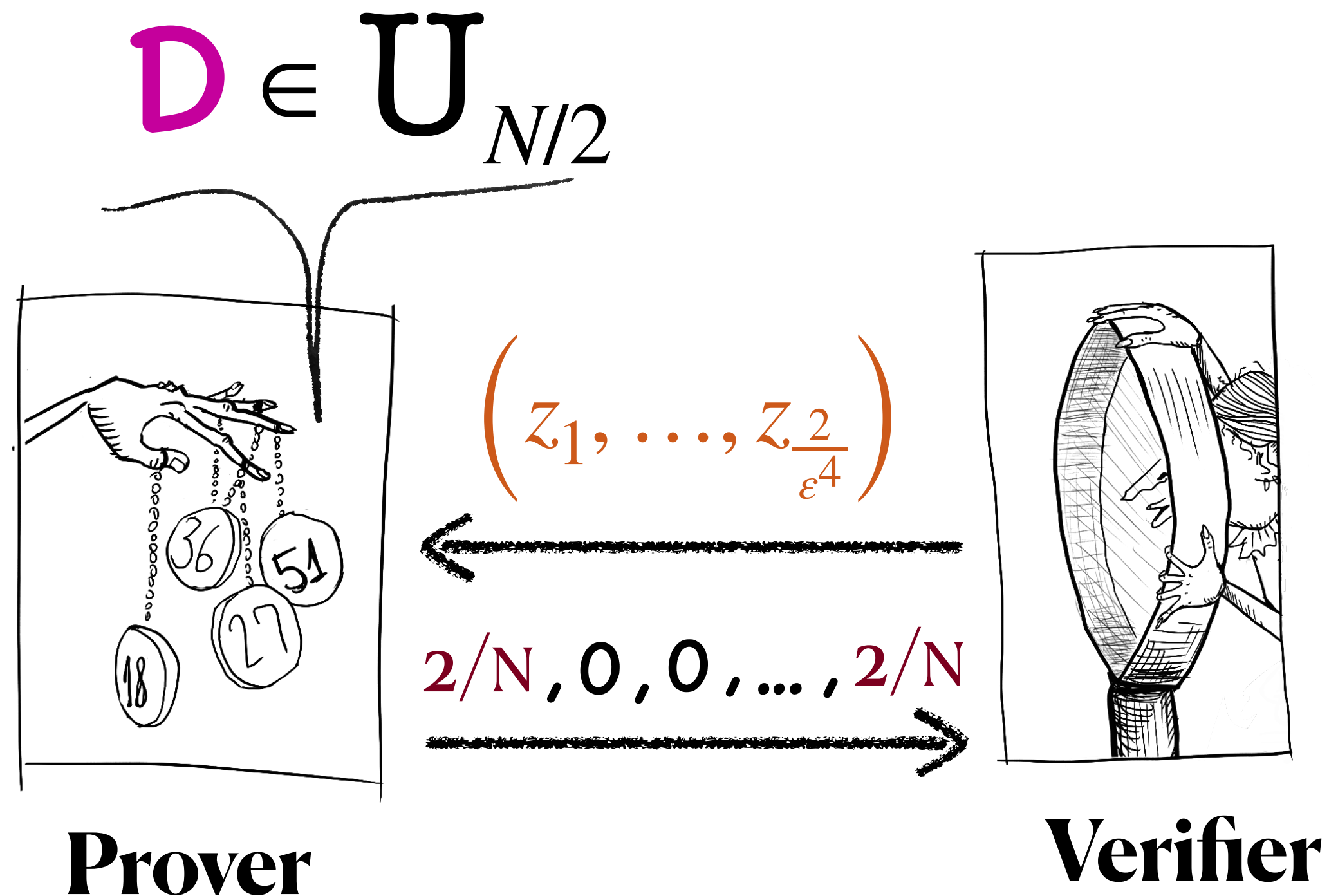
Completeness:

Soundness:

1. Draw $s = O(\sqrt{N})$ samples.
Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$

2. Draw ϵ^{-4} samples from $\mathcal{U}_{[N]}$.
 Draw ϵ^{-4} samples from D .
 Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$
Test 2:
 $\approx \textit{half}$ the $\mathcal{U}_{[N]}$ -samples have tag $2/N$
 $\textit{ALL } D$ -samples have tag $2/N$

Verifying Uniformity Over Half Domain



Completeness:

Soundness:

1. Draw $s = O(\sqrt{N})$ samples.

Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$

2. Draw ϵ^{-4} samples from $U_{[N]}$.

Draw ϵ^{-4} samples from D .

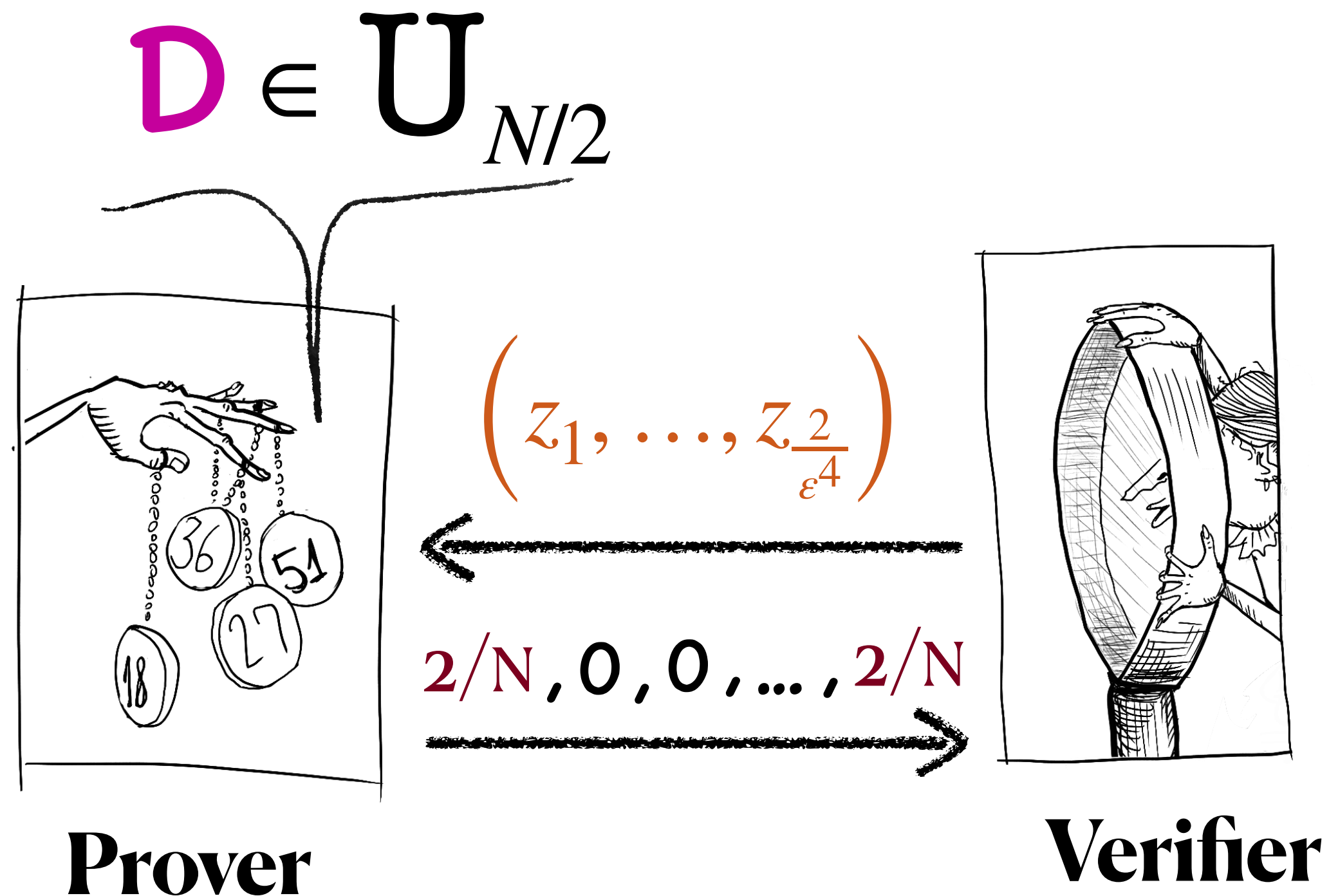
Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$

Test 2:

$\approx$ *half* the $U_{[N]}$ -samples have tag $2/N$

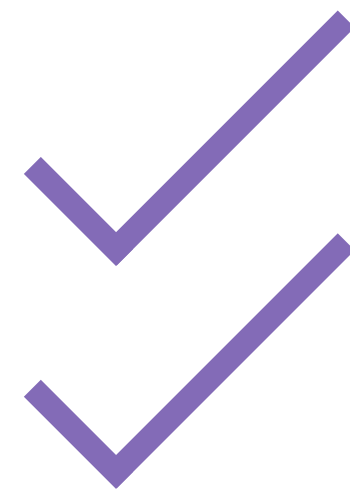
ALL D -samples have tag $2/N$

Verifying Uniformity Over Half Domain



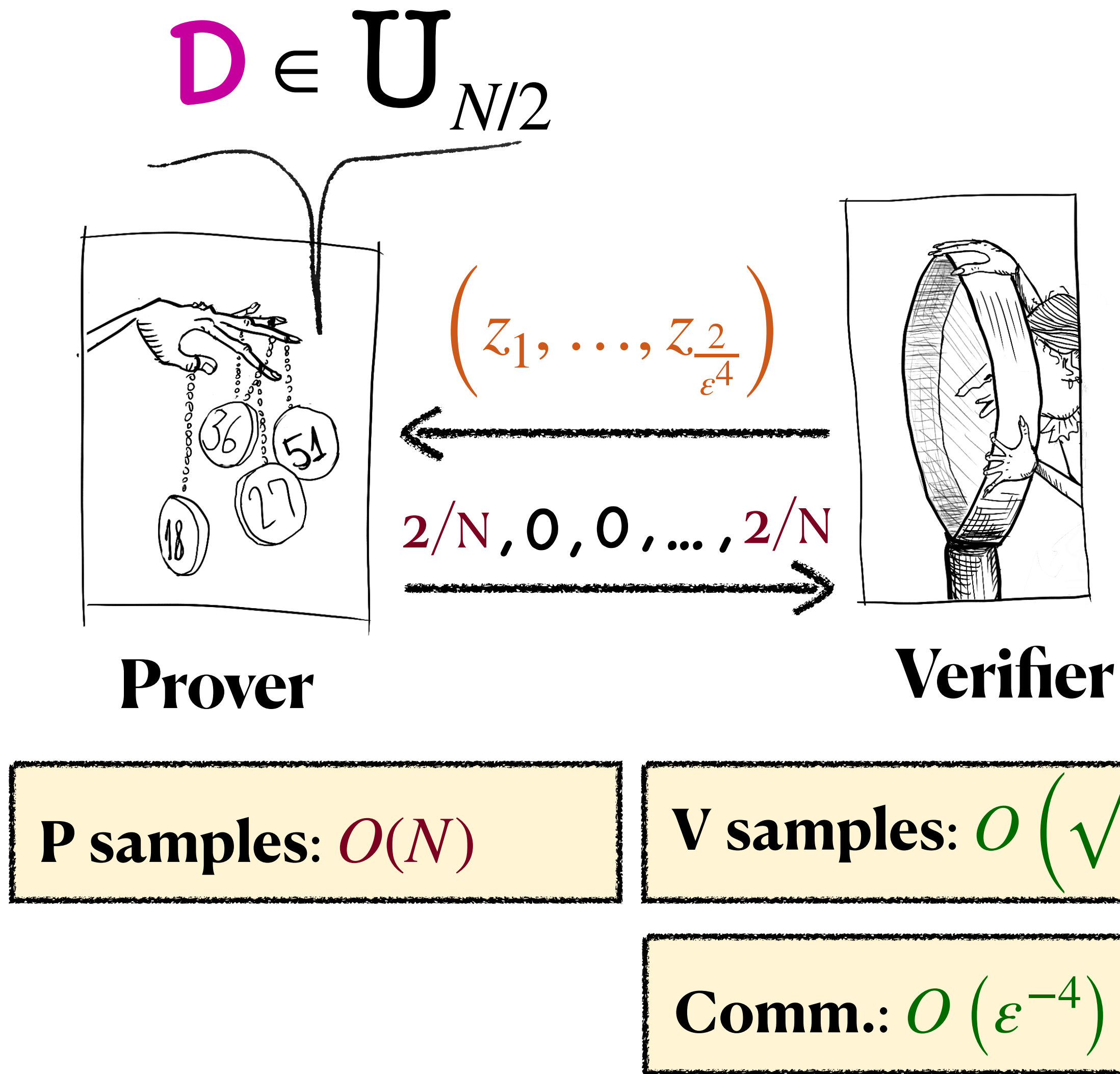
Completeness:

Soundness:



1. Draw $s = O(\sqrt{N})$ samples.
Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$
2. Draw ϵ^{-4} samples from $\mathcal{U}_{[N]}$.
 Draw ϵ^{-4} samples from D .
 Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$
Test 2:
 $\approx \text{half}$ the $\mathcal{U}_{[N]}$ -samples have tag $2/N$
ALL D -samples have tag $2/N$

Verifying Uniformity Over Half Domain



1. Draw $s = O(\sqrt{N})$ samples.

Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{2}{N}$

2. Draw ϵ^{-4} samples from $U_{[N]}$.

Draw ϵ^{-4} samples from D .

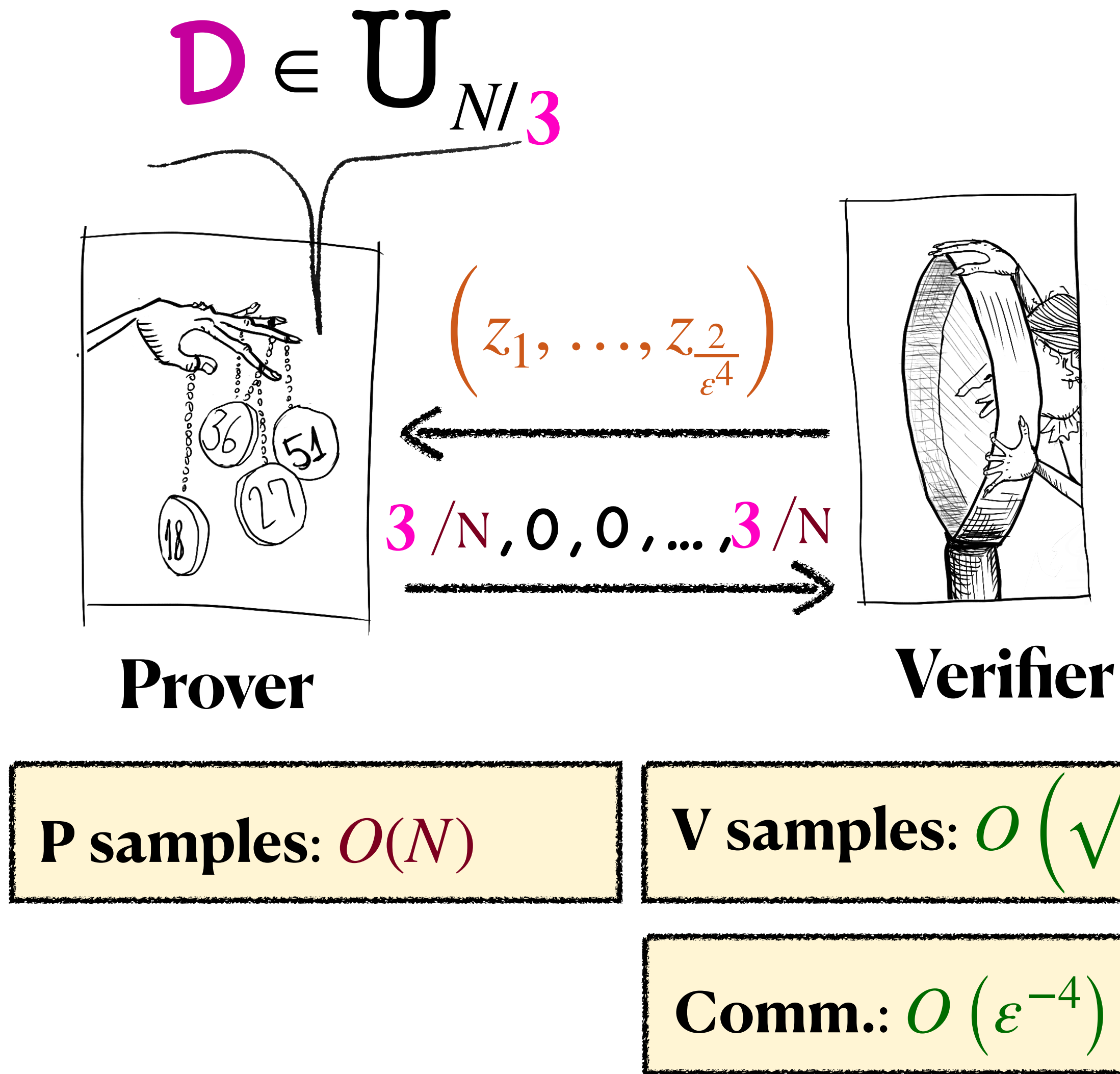
Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$

Test 2:

$\approx$ *half* the $U_{[N]}$ -samples have tag $\frac{2}{N}$

ALL D -samples have tag $\frac{2}{N}$

Verifying Uniformity Over **Third** Domain



1. Draw $s = O(\sqrt{N})$ samples.

Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{3}{N}$

2. Draw ϵ^{-4} samples from $U_{[N]}$.

Draw ϵ^{-4} samples from D .

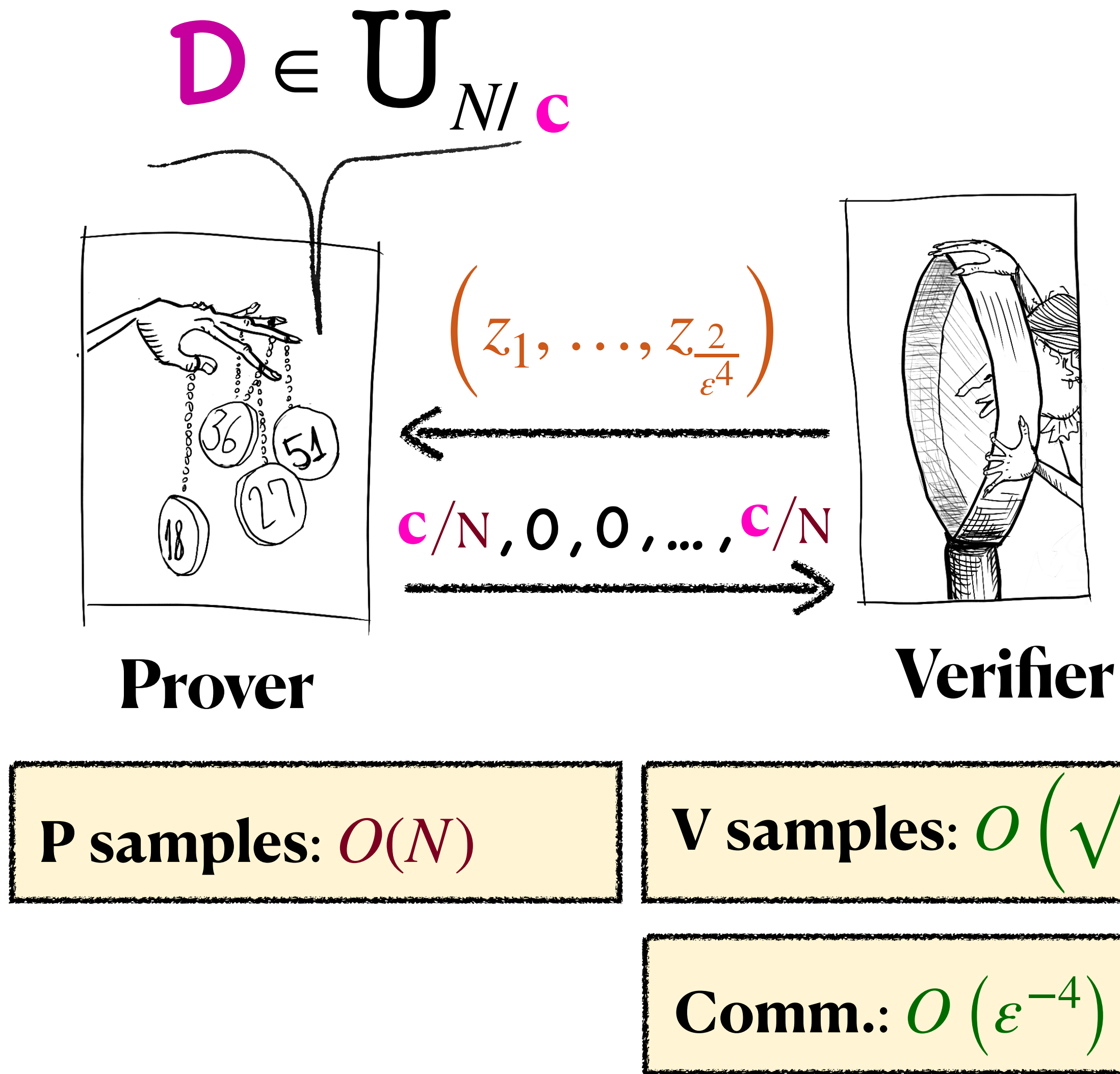
Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$

Test 2:

$\approx$ **third** the $U_{[N]}$ -samples have tag $3/N$

ALL D -samples have tag $3/N$

Verifying Uniformity Over **Fraction of Domain**



1. Draw $s = O(\sqrt{N})$ samples.

Test 1: #of 2-collisions is $\approx \binom{s}{2} \cdot \frac{\mathbf{c}}{N}$

2. Draw ϵ^{-4} samples from $U_{[N]}$.

Draw ϵ^{-4} samples from D .

Randomly mix: $(z_1, \dots, z_{\frac{2}{\epsilon^4}})$

Test 2:

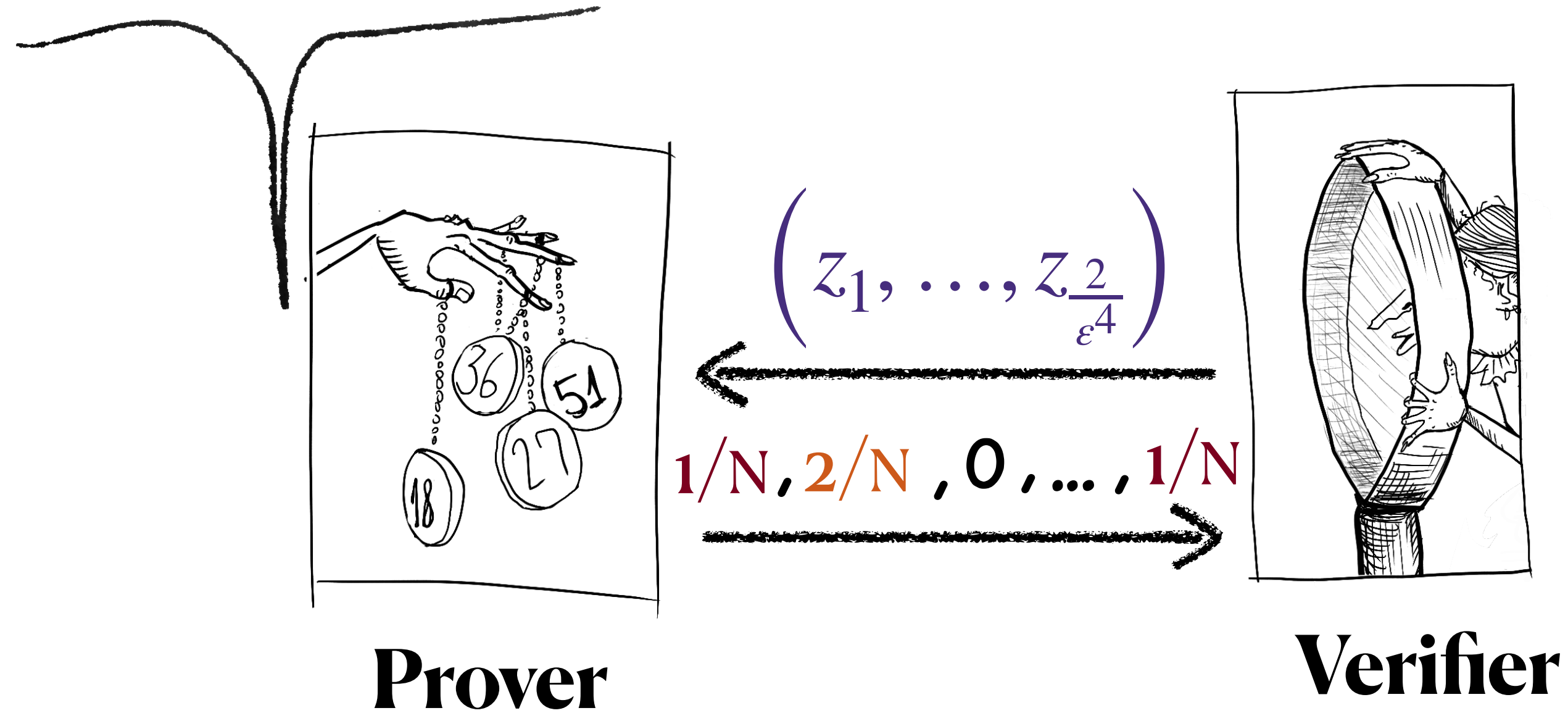
$\approx 1/\mathbf{c}$ the $U_{[N]}$ -samples have tag $\mathbf{c}/N$

ALL D -samples have tag $\mathbf{c}/N$

D is supported over:

S_1 elements of probability $1/N$

S_2 elements of probability $2/N$



P samples: $O(N)$

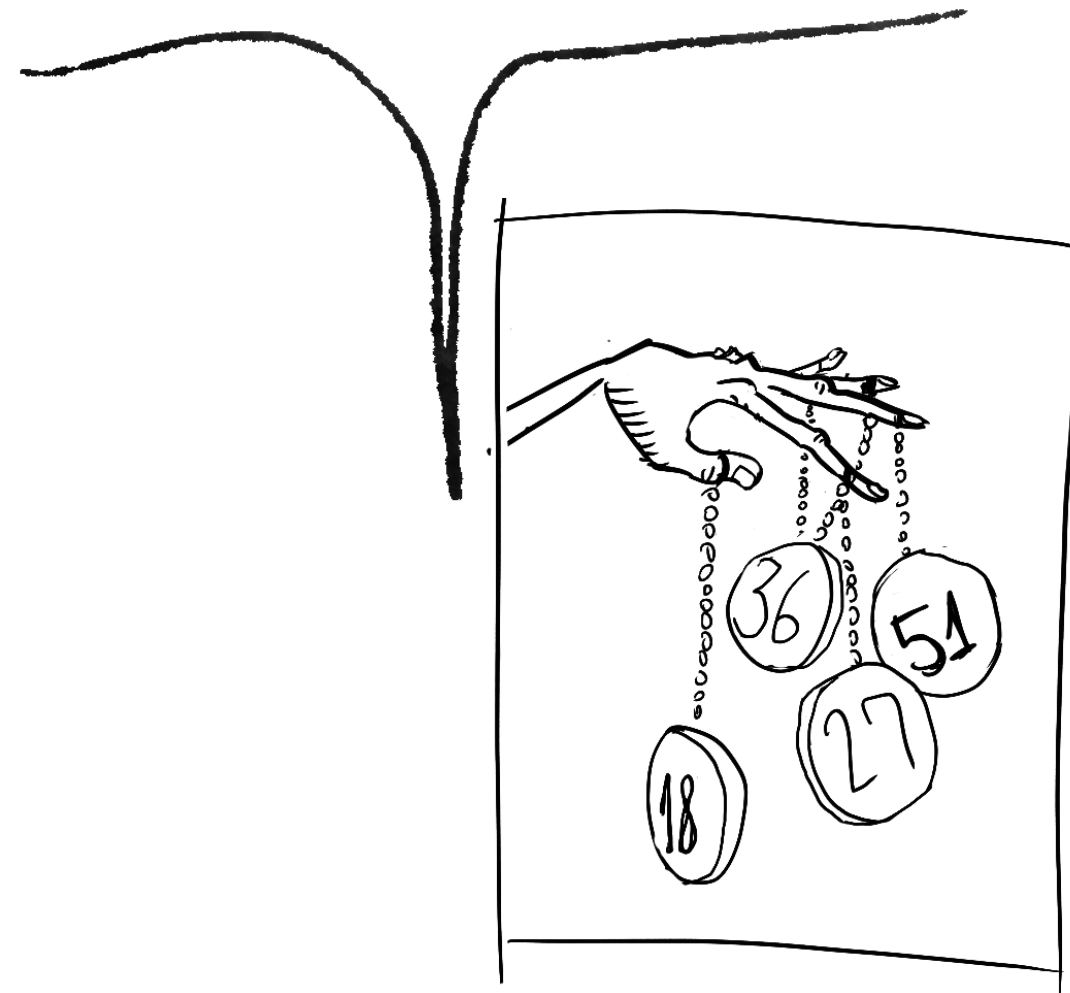
Comm.: $O(\epsilon^{-4})$

V samples: $O(\sqrt{N})$

$$\forall j \in [\log N / \varepsilon]$$

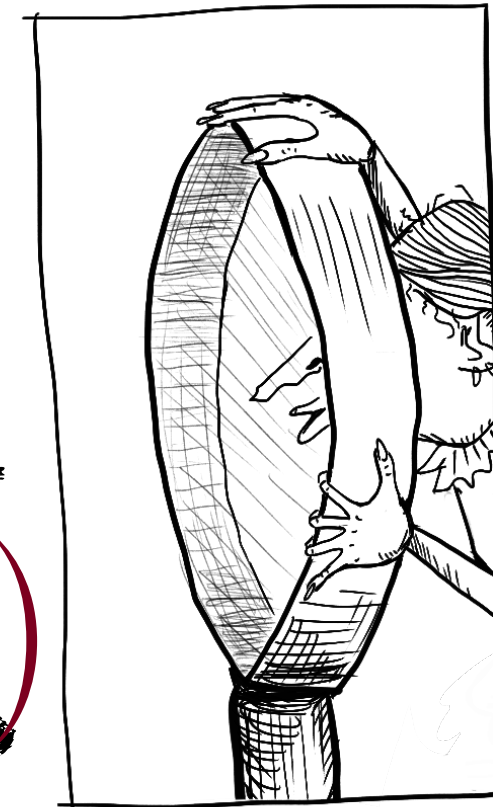
D is supported over:

$$S_j \text{ elements w.p. } \approx \frac{\varepsilon \cdot (1 + \varepsilon)^j}{N}$$



Prover

$$\begin{array}{c} \xleftarrow{(z_1, \dots, z_s)} \\ \xrightarrow{(\bar{D}_{z_1}, \bar{D}_{z_2}, \dots, \bar{D}_{z_s})} \end{array}$$



Verifier

P samples: $O(N)$

Comm.: $O(\sqrt{N})$

V samples: $O(\sqrt{N})$

Partition of $[0,1]$:

$$\left(\frac{\varepsilon}{N}, \frac{\varepsilon}{N} \cdot (1 + \varepsilon) \right]$$

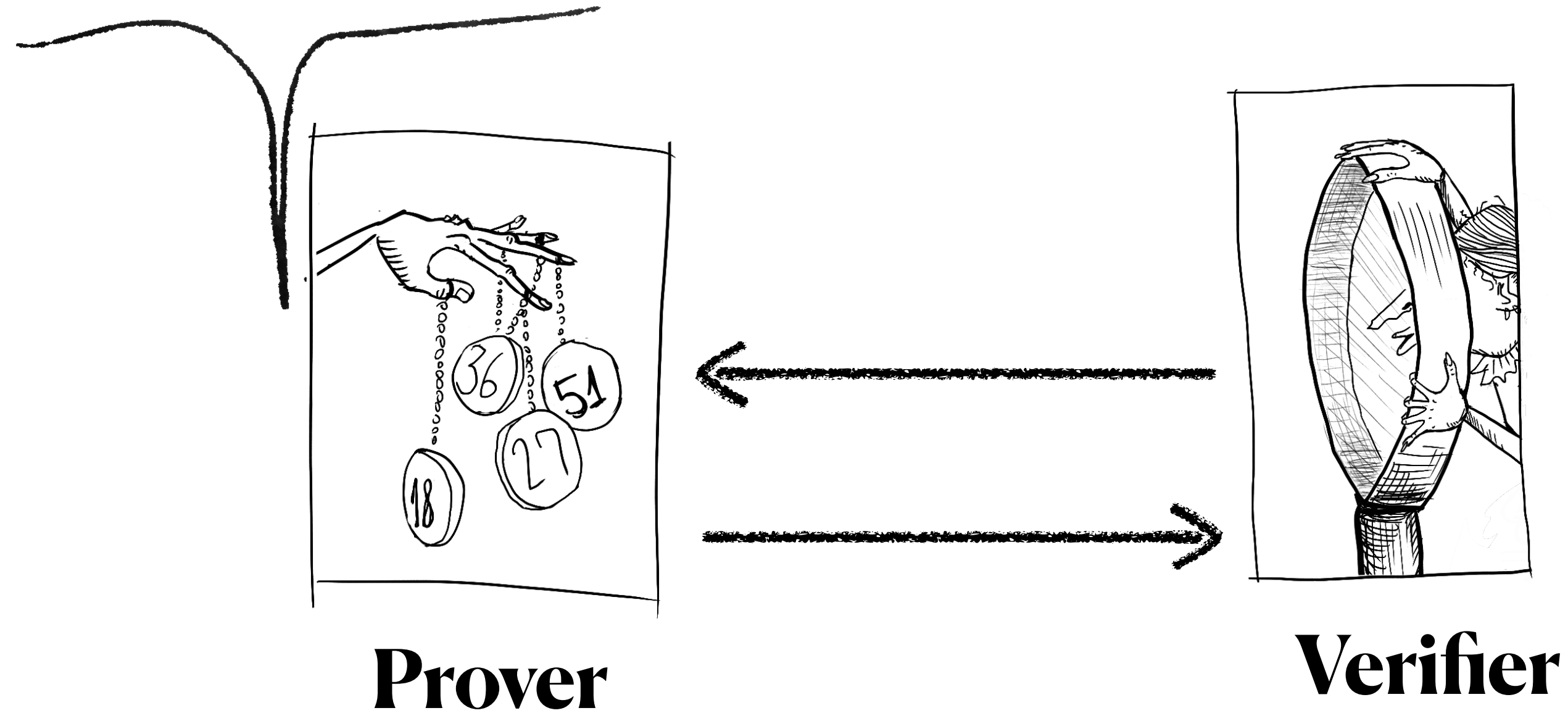
$$\left(\frac{\varepsilon}{N} \cdot (1 + \varepsilon), \frac{\varepsilon}{N} \cdot (1 + \varepsilon)^2 \right]$$

$\vdots$

$$\left(\frac{\varepsilon}{N} \cdot (1 + \varepsilon)^j, \frac{\varepsilon}{N} \cdot (1 + \varepsilon)^{j+1} \right]$$

$\vdots$

D has probability histogram $hist_0$

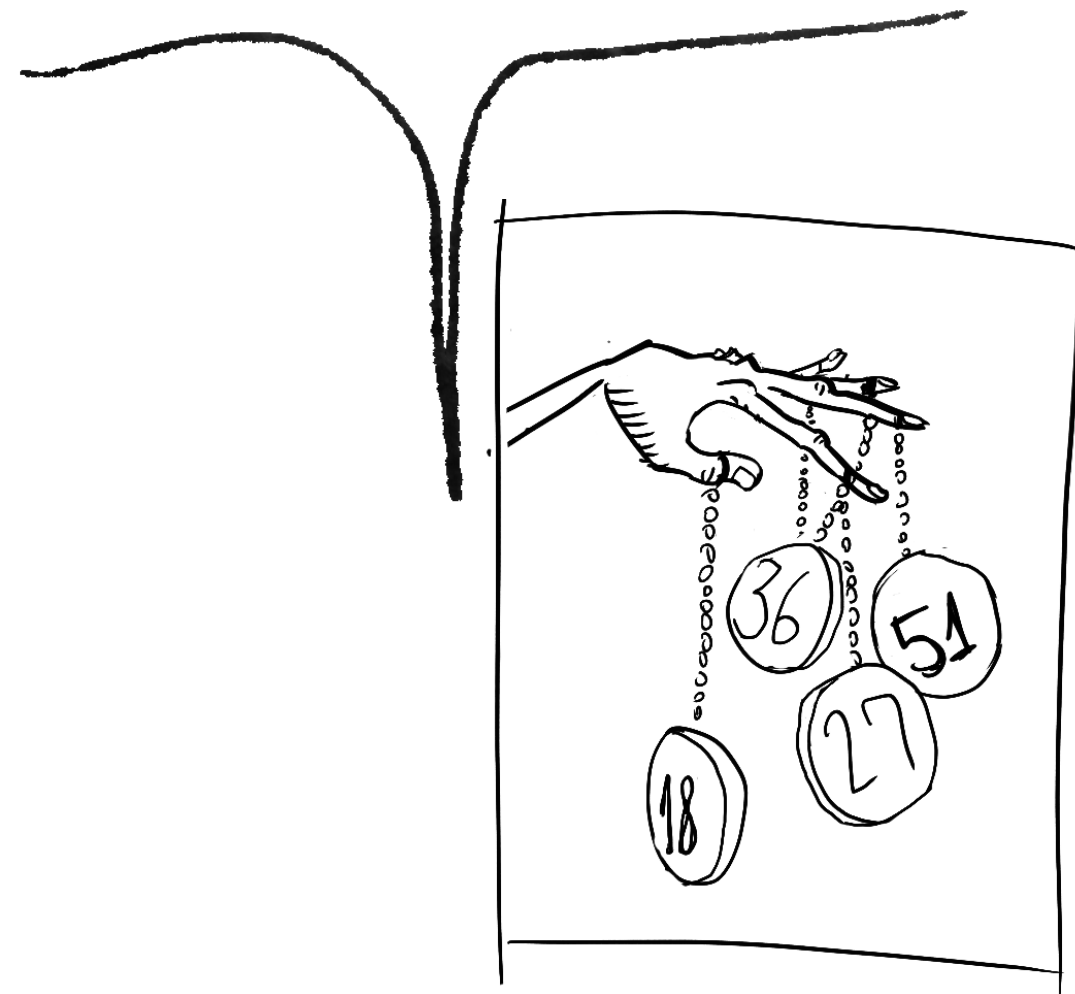


P samples: $O(N)$

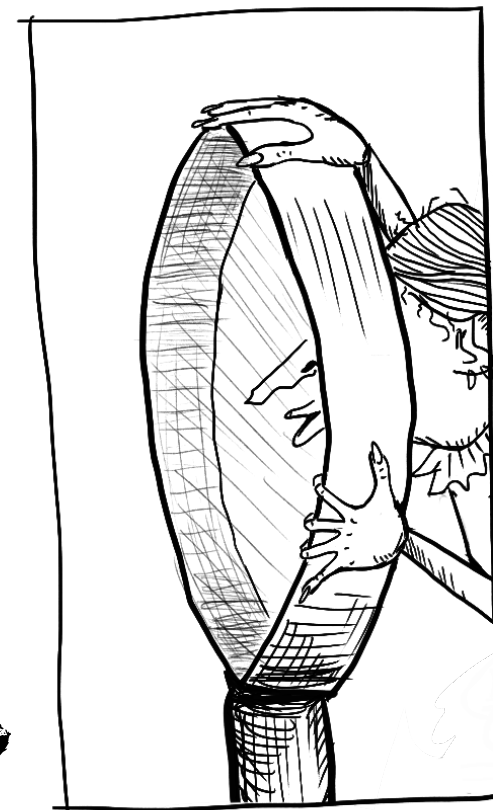
Comm.: $O(\sqrt{N})$

V samples: $O(\sqrt{N})$

D has probability histogram $hist_0$



Prover



Verifier

Verify histogram, then compute:

- * Shannon entropy!
- * Distance from uniformity!
- * Support size!

P samples: $O(N)$

Comm.: $O(\sqrt{N})$

V samples: $O(\sqrt{N})$

Theorem [H-Rothblum'22,H-Rothblum'23]

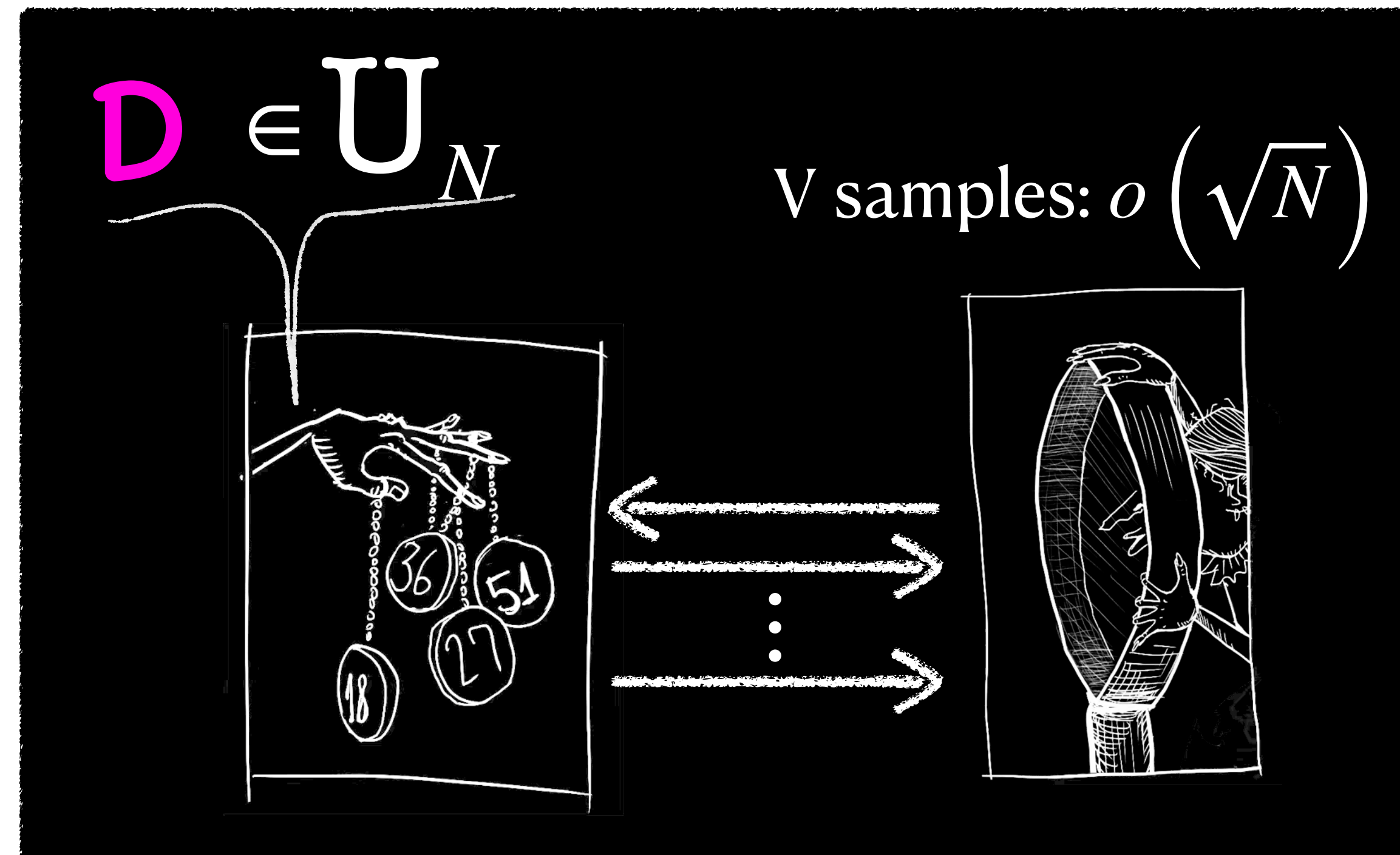
Every **label-invariant** distribution property admits a 2-message tolerant proof system with:*

P samples: $O(N)$ **Comm.:** $O\left(\sqrt{N}\right)$ **V samples:** $O\left(\sqrt{N}\right)$

Can we verify uniformity with $o\left(\sqrt{N}\right)$?

Can we verify uniformity with $o\left(\sqrt{N}\right)$?

Assume...



Can we verify uniformity with $o\left(\sqrt{N}\right)$?

Tester [CG'17]

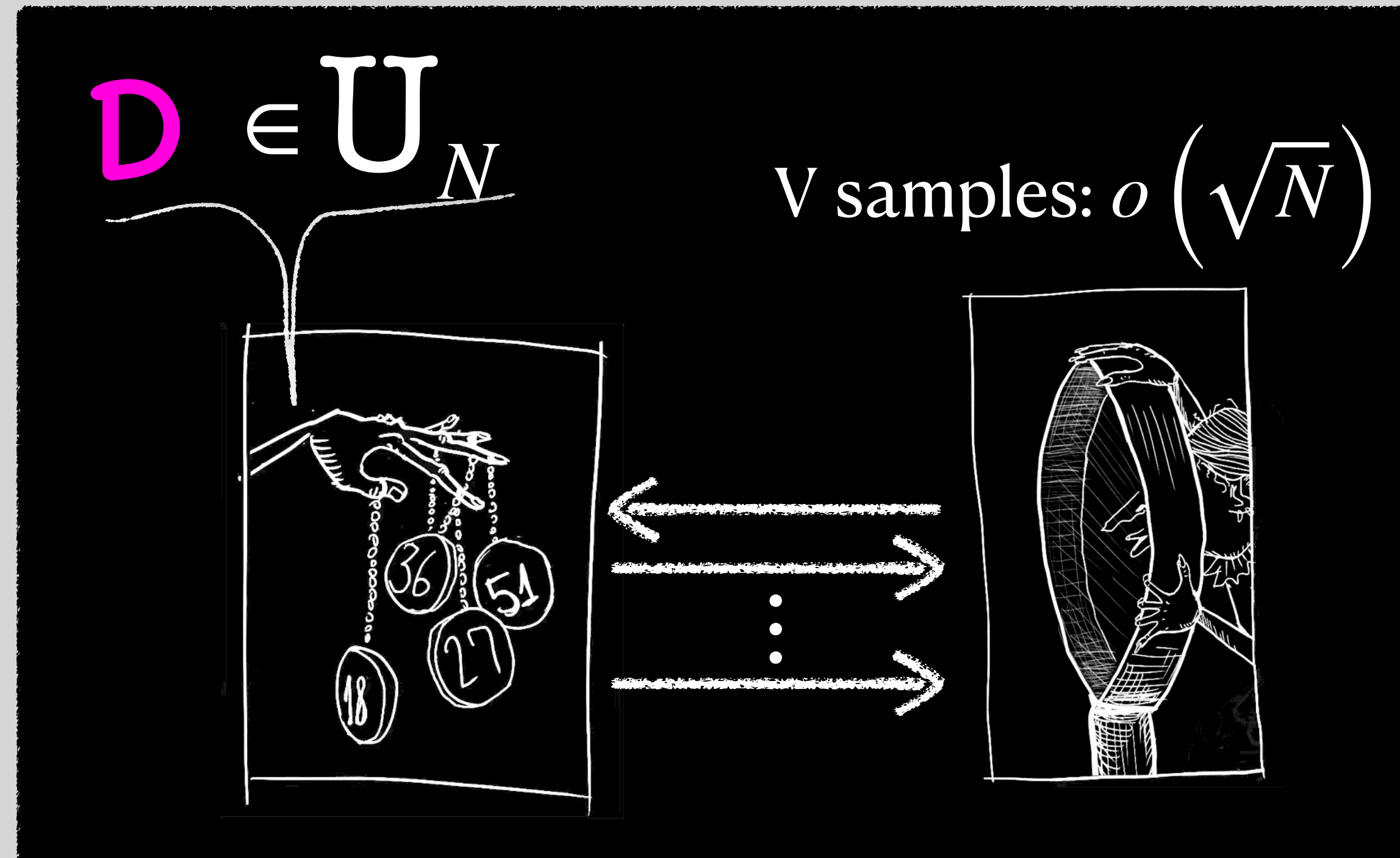
Total samples: $o\left(\sqrt{N}\right)$

Assume...

Simulate by answering
according to $U_{[N]}$

*No samples
needed!*

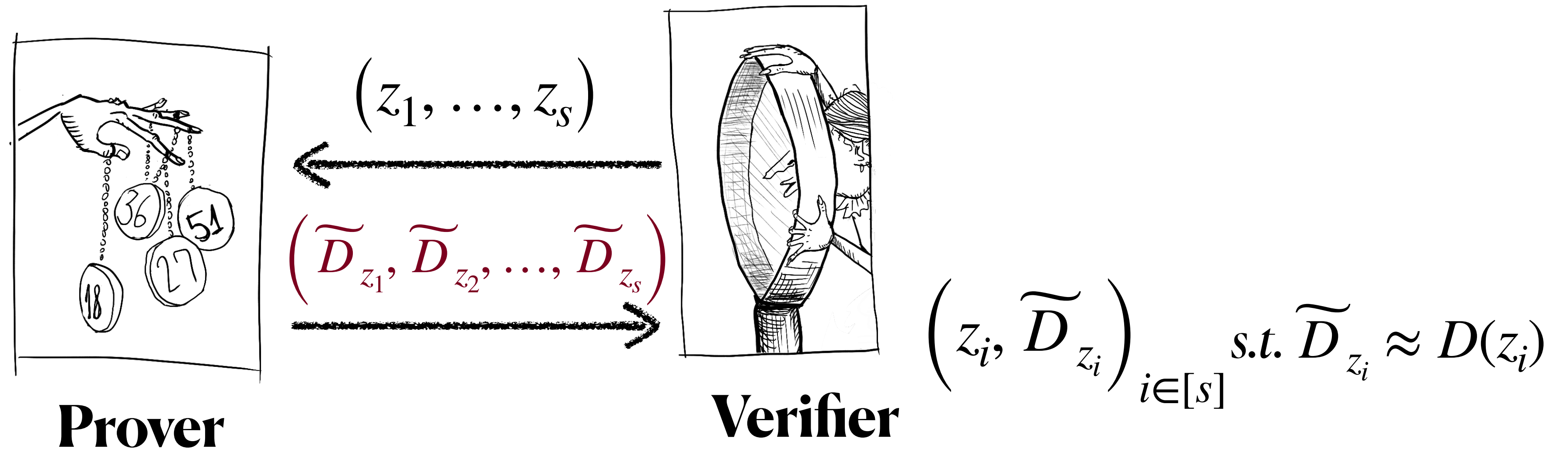
$\mathbf{D} \in \mathbf{U}_N$



Simulate V

$o\left(\sqrt{N}\right)$ samples

Main tool:



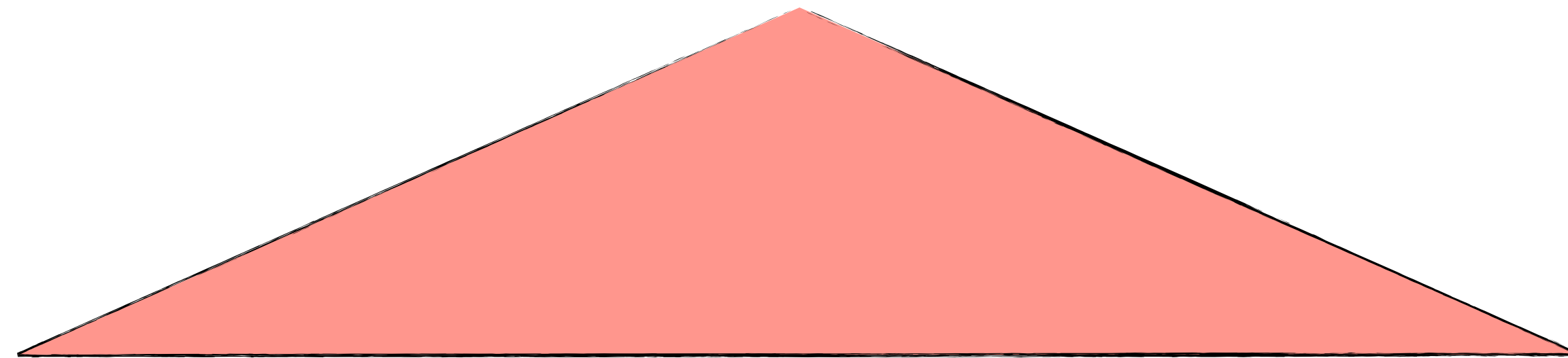
P samples: $O(N)$

Comm.: $O(\sqrt{N})$

V samples: $O(\sqrt{N})$

$$\mathbf{D} \quad \left((1, D(1)) , (2, D(2)) , \dots , (N, D(N)) \right)$$

$$b \in \{0,1\}$$



$$\mathbf{D} \left((1, D(1)) , (2, D(2)) , \dots , (N, D(N)) \right)$$

$$b \in \{0,1\}$$

*Almost **IPP**...*

Sample and not query access

P must know the accessed locations.

D

$\left((1, D(1)) , (2, D(2)) , \dots , (N, D(N)) \right)$

Theorem [H-Rothblum'24]

Every distribution property of *reasonable complexity* admits a

$(\log N)$ -messages tolerant proof system with:

P samples: $O(N^{1.1})$ **Comm.:** $O(N^{0.95})$ **V samples:** $O(N^{0.9})$

$b \in \{0,1\}$

Almost **IPP**...

Sample and not query access

P must know the accessed locations.

D $\left((1, D(1)) , (2, D(2)) , \dots , (N, D(N)) \right)$

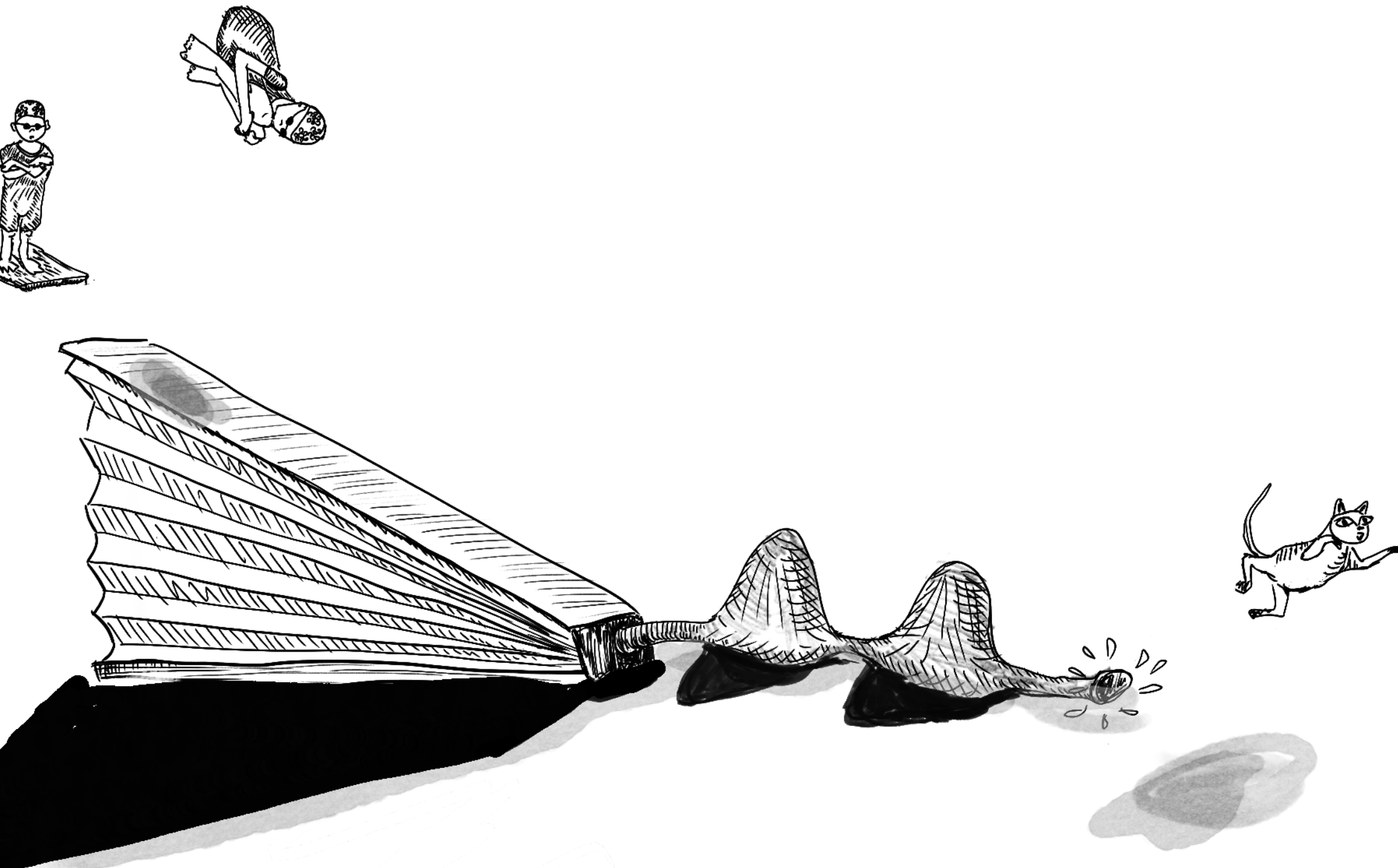
More Results...

- Exploring *features of protocols*: public coin [H'24], differential privacy [DDTZ'26], zero knowledge [CPSTY'26], doubly sublinear proof system [GHR'25])
- *Argument systems* for every* property with $O(N^{1/2})$ V samples & comm. [H-Rothblum'25].
- Exploring *different access models*: PCond verification [BBCS'25]
- Lower bounds:
 - *Proving* membership in distribution properties can be *strictly harder than testing* [H-Rothblum'25].
 - *The comm. complexity* of verifying label-invariant properties is $\widetilde{\Theta}\left(N \cdot \sum_x D^2(x)\right)$ [H'26].

Summary and Future work

Verifying distribution properties can be done more efficiently than testing.

New lens for insights into distribution properties, and a rich area for future work.



Summary and Future work

Verifying distribution properties can be done more efficiently than testing.

New lens for insights into distribution properties, and a rich area for future work.

Open Questions:

Improving parameters, exploring different features (ZK, DP...).

Proofs for different models (continuous distributions, structured domains...)

Applications:

- *Verifiable Data Science, verifiable learning*
- Applications for AI Verification:
 - Unique properties, domains, and access models.

