# Affine Grassmannians

# Nicholas Lai

#### March 6th, 2020

# Intuition

Let G be a connected reductive algebraic group, k a field.

Affine Grassmannian "=" {trivialisable G-bundles on formal punctured disc}/{triv G-bundles on formal disc

If  $k = \mathbb{C}$ , *G*-bundles are

- Trivialisable after finite number of punctures
- Then, bundles "is" taking disc over each punctures as trivialising cover, and specify transition maps.

 $\{\text{triv } G\text{-bundle over Spec} A, \text{trivialisation}\} \leftrightarrow \{\text{triv of trivial bundle}\} = \max(\text{Spec} A, G) = G(A).$ 

punctured disc  $\leftrightarrow$  Speck((t)) disc  $\leftrightarrow$  Speck[[t]]

Naively, we want

$$\operatorname{Gr}_G = G(k((t)))/G(k[[t]]).$$

## Affine Grassmannians for $GL_n$ .

We will do everything over  $\mathbbm{Z}$  and base change to a field later.

**Definition.** For any ring R, an R-family of lattice is a finite locally free R[[t]]-submodules  $\Lambda \subseteq R((t))^n$  such that  $\Lambda \otimes_{R[[t]]} R((t)) = R((t))^n$ . e.g. the standard lattice  $\Lambda_0(R) = R[[t]]^n$ . Def. An affine Grassmannian is the functor

$$\operatorname{Gr}_{\operatorname{GL}_n} : \operatorname{\mathbf{Rings}} \to \operatorname{\mathbf{Sets}}$$

$$R \mapsto \{R - \text{lattice}\}.$$

Here **Rings** are commutative. We would like this functor to be a scheme, but it's not.

However, it is an ind-scheme, a colimit of schemes.

The affine Grassmannian can be decomposed into a "system" of schemes, and we can study the geometry piecewise.

We want to say  $\operatorname{Gr}_{\operatorname{GL}_n} = \bigcup_N \operatorname{Gr}_{\operatorname{GL}_n}^{(N)}$ .

**Def.** For integers  $a \leq b$ , define

$$\operatorname{Gr}_{[a,b]}(R) = \{\Lambda \in \operatorname{Gr}_{\operatorname{GL}_n}(R) : t^b \Lambda_0(R) \subseteq \Lambda \subseteq t^a \Lambda_0(R) \}.$$

This gives a filtered system, and as functors,

$$\operatorname{Gr}_{\operatorname{GL}_n}(R) = \operatorname{colim}_{a \le b} \operatorname{Gr}_{[a,b]}(R).$$

(finite locally free over R implies finitely generated)

The main point is the following:

**Theorem.** The functor  $\operatorname{Gr}_{[a,b]}$  is represented by a closed subscheme of

$$Grass(M_{[a,b]}) : \mathbf{Rings} \to \mathbf{Sets}$$

 $R \mapsto \{N \subseteq t^a \Lambda_0(R)/t^b \Lambda_0(R) : t^a \Lambda_0(R)/t^b \Lambda_0(R) \text{ is f.g. loc free } R - mod\}.$ 

**Rem.** Grass $(M_{[a,b]})$  is represented by a smooth proper scheme over  $\mathbb{Z}$  and is the finite disjoint union over  $0 \le k \le \mathrm{rk} t^a \Lambda_0(R)/t^b \Lambda_0(R)$  of classical Grassmannians of rank k.

What is  $Gr_{[a,b]}$  exactly?

Define  $\operatorname{Grass}^{t}(M_{[a,b]})(R) = \{N \in \operatorname{Grass}(M_{[a,b]})(R) : tN \subset N\}.$ 

This is a closed subscheme of  $Grass(M_{[a,b]})$ .

Remark  $\operatorname{Grass}^{t}(M_{[a,b]})$  is a union of Springer fibres.

Claim:  $\operatorname{Gr}_{[a,b]} \to \operatorname{Grass}^t(M_{[a,b]})$ 

$$\Lambda \mapsto \Lambda/t^b \Lambda_0(R).$$

Key idea:  $\operatorname{Grass}^t(M_{[a,b]}) \to \operatorname{Spec}\mathbb{Z}$  of finite type.

 $N \in \text{Grass}^t(M_{[a,b]})$  is defined over finitely-generated Z-algebra. So we can ssume that R is Noetherian.

 $R[t] \rightarrow R[[t]]$  is flat (Noetherian assumption is necessary).

Therefore,  $\Lambda = \ker(t^a \Lambda_0(R) \to t^a \Lambda_0(R)/t^b \Lambda_0(R))$ 

 $\Lambda_f = \operatorname{Ker}(t^a R[t] \to t^a R[t]/t^b R[t])$ 

 $\Lambda_f$  is finite locally free, so the map is surjective.

**Remark.** The idea of replacing R[[t]] with R[t] is a special case of Beauville-Laszlo's realisation of Gr via a global curve.

Digression. (strict)-Ind Schemes.

Example:  $\mathbb{A}^{\infty} = \bigcup_{i>0} \mathbb{A}^i$ ,  $\mathbb{A}^i \subseteq \mathbb{A}^{i+i}$  on the first *i*-coordinates.

Definition. A (strict) ind-schemes is a functor

 $X : \mathbf{Aff} - \mathbf{sch}^{op} \to \mathbf{Sets}.$ 

Which can be written as  $X \approx \operatorname{colim}_{i \in I} X_i$  as a filtered colimit of schemes where all transition maps  $X_i \to X_j$ ,  $i \leq j$  are closed immersions.

**Remark.** We will identify  $\mathbf{AffSch}^{op} = \mathbf{Rings}$ .

e.g.  $\mathbb{A}^{I}_{\mathbb{Z}}: T \mapsto \bigoplus_{i \in I} P(T, \mathcal{O}_{T}) = \operatorname{colim}_{J \subset I} \mathbb{A}^{J}_{\mathbb{Z}} \to \mathbb{A}^{|J|}.$ 

Remark. We will not reference the fpqc topology on AffSch. However,

**Lemma.** Every ind-scheme satisfies the sheaf condition for the fpqc topology on **AffSch**.

- This works over colimits of schemes over filtered index category
- Ran spaces are not ind-schemes.

**Lemma.** Let  $X \to Y$  be a map of functors **AffSch**<sup>op</sup>  $\to$  **Sets**.

Suppose for all affine schemes  $T \to Y$  the fibre product  $X \times_Y T$  is a scheme. Then if Y is an ind-scheme, so is X.

**Cor.**  $X_i \subset X$  is representable by close immersion.

Lemma. The category of ind-schemes IndSch has the properties:

- SpecZ is the final object
- Closed under fibre products (admits finite limits)
- Directed limits with affine transitions
- Arbitrary disjoint union

Definition. For a local property P of schemes, an Ind-scheme has P if there exists a representation  $X = \operatorname{colim}_{i \in I} X_i$  such that every  $X_i$  has P.

Lemma. Scheme has property P if and only if it has P as an ind-scheme.

**Def.** Same for morphism

Lemma. The same lemma is true but only for quasi-compact map of schemes.

Base change. Can do all of this over scheme S instead  $Spec\mathbb{Z}$ , we get  $IndSch_S$ , (called a slice category).

- AffSch<sub>S</sub> has a notion of fpqc topology.
- $S = \operatorname{Spec} R$ ,  $\operatorname{AffSch}^{op} = R \operatorname{Alg}$ .

Can run the whole machinery verbatim.

Algebraic spaces can be extended to IndAlgSp.

Lemma.  $AlgSp \cap IndSch = Sch.$ 

## Back to affine Grassmannians

Base change  $\operatorname{Gr}_{\operatorname{GL}_n,S}$  : AffSch<sup>op</sup>  $\rightarrow$  Set.

 $T = \operatorname{Spec} R \mapsto \operatorname{Gr}_{\operatorname{GL}_n, \mathbb{Z}}(R)$  is representable by Ind-proper, ind-scheme

 $\operatorname{Gr}_{\operatorname{GL}_n} \times_{\operatorname{Spec}\mathbb{Z}} S \to S$ 

#### Some geometry

Want to construct ind-affine open covers of  $\operatorname{Gr}_{\operatorname{GL}_n}$  for  $\mu \in \mathbb{Z}^n$ , let

$$t^{\mu} = \operatorname{diag}(t^{\mu_1}, \cdots, t^{\mu_n}) \in \operatorname{GL}_n(\mathbb{Z}((t))).$$

For a ring R, let

$$\Lambda_{\mu}^{-}(R) = t^{\mu}(t^{-1}R[t^{-1}]^{n})$$

considered as  $R[t^{-1}]$ -submodule of R((t)).

Define  $U_{\mu}(R) = \{\Lambda \in \operatorname{Gr}_{\operatorname{GL}_n}(R) : \Lambda_{\mu}(R) \oplus \Lambda \equiv R((t))^n \text{ as } R \operatorname{-mod} \}.$ 

Note that  $\Lambda_0 \subset U_0$ ,  $U_{\mu}$  is  $t^{\mu}$ -translate of  $U_0$  under the action of  $\operatorname{GL}_n(\mathbb{Z}((t)))$  on  $\operatorname{Gr}_{\operatorname{GL}_n}$ .

This is related to Birkhoff decomposition for Kac-Moody algebras and double Tits-systems (twin buildings).

 $\operatorname{Gr}_{\operatorname{GL}_n} = \bigcup_{\mu \in \mathbb{Z}^n} U_{\mu}$ , in particular,  $U_{\mu} \subseteq \operatorname{Gr}$  is represented by open immersions on schemes.

Idea:  $U_{\mu} \cap \operatorname{Gr}_{[a,b]}$  is pullbacks of standard open covers of classical Grassmannians. Cor.  $U_{\lambda} \cap U_{\mu} \neq \emptyset \Leftrightarrow |\mu| = |\lambda|$ . This this case, all fivers of  $U_{\lambda} \cap U_{\mu} \to \operatorname{Spec}\mathbb{Z}$  are nonempty.