

Map Making and Power Spectrum Estimation of the CMB Radiation

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OVERVIEW

- **MAP MAKING**

- Why Preconditioned Conjugate Gradients Work?
- Best Diagonal Preconditioner
- Implementation of PCG - speed and parallelism

- **WHY STOP THERE?**

 - POWER SPECTRUM ESTIMATION**

 - The Monte Carlo Approach
 - Map Cutting - The Killer of All Matrix Structure

- **CONCLUSIONS**

- **RECOMMENDATION FOR SCAN STRATEGIES**

MAP MAKING

Is the process of solving

$$N^{-1}x = A^T T d_t, \quad \text{where } N^{-1} = A^T T A$$

- **Structure Exploiting Direct Methods**
e.g. Displacement Structure approach
can't handle rectangular matrices
- **Second Best: Iterative Methods**
 N^{-1} - positive definite, so PCG
- **PCG converges very fast:**
about 30 iterations for rel. res. 10^{-4}

THE POINTING MATRIX A

– Has orthogonal columns:

$$A^T A = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 4 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 3 \end{bmatrix} = D_1^2.$$

– AD_1^{-1} - orthogonal

– Equivalent to making all pixels “equally” important

– Equivalent to diagonal preconditioning with D_1

● D_1 – “correct” diagonal preconditioner because

$$\text{cond}((D_1^{-1} A^T)T(AD^{-1})) = \text{cond}(T) \approx 10^{-2}$$

● Saves 1 iteration from PCG.

IMPLEMENTING PCG

- Converges fast, bottleneck – multiplication by N^{-1}
- Issues:
 - Can't store vector of size 10^9 in RAM - need parallelism to avoid disk access
 - Exploit structure to gain speed

$$N^{-1}v = A^T \cdot T \cdot A \cdot v$$

$$= \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 9 & 5 & 1 & & & & & & \\ 5 & 9 & 5 & 1 & & & & & \\ 1 & 5 & 9 & 5 & 1 & & & & \\ & 1 & 5 & 9 & 5 & 1 & & & \\ & & 1 & 5 & 9 & 5 & 1 & & \\ & & & 1 & 5 & 9 & 5 & 1 & \\ & & & & 1 & 5 & 9 & 5 & 1 \\ & & & & & 1 & 5 & 9 & 5 \\ & & & & & & 1 & 5 & 9 \end{bmatrix} \cdot \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix} \cdot v.$$

- $w = Av$ =observing ... OK
- $u = Tw$... band Toeplitz product
- $A^T u$ =repointing on the sky ... OK

Best: Block-Toeplitz. Instead of:

$$\begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ u_4 \\ u_5 \\ u_6 \\ u_7 \\ u_8 \\ u_9 \end{bmatrix} = \begin{bmatrix} 9 & 5 & & & & & & & \\ & 5 & 9 & 5 & & & & & \\ & & 5 & 9 & 5 & & & & \\ & & & 5 & 9 & 5 & & & \\ & & & & 5 & 9 & 5 & & \\ & & & & & 5 & 9 & 5 & \\ & & & & & & 5 & 9 & 5 \\ & & & & & & & 5 & 9 & 5 \\ & & & & & & & & 5 & 9 \end{bmatrix} \cdot \begin{bmatrix} w_1 \\ w_2 \\ w_3 \\ w_4 \\ w_5 \\ w_6 \\ w_7 \\ w_8 \\ w_9 \end{bmatrix}$$

Do this:

$$\begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ u_4 \\ u_5 \\ * \end{bmatrix} = \begin{bmatrix} 9 & 5 & & & 5 \\ & 5 & 9 & 5 & & \\ & & 5 & 9 & 5 & \\ & & & 5 & 9 & 5 \\ & & & & 5 & 9 & 5 \\ 5 & & & & & 5 & 9 \end{bmatrix} \cdot \begin{bmatrix} w_1 \\ w_2 \\ w_3 \\ w_4 \\ w_5 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} * \\ u_5 \\ u_6 \\ u_7 \\ u_8 \\ u_9 \end{bmatrix} = \begin{bmatrix} 9 & 5 & & & 5 \\ & 5 & 9 & 5 & & \\ & & 5 & 9 & 5 & \\ & & & 5 & 9 & 5 \\ & & & & 5 & 9 & 5 \\ 5 & & & & & 5 & 9 \end{bmatrix} \cdot \begin{bmatrix} 0 \\ w_5 \\ w_6 \\ w_7 \\ w_8 \\ w_9 \end{bmatrix}$$

BLOCK TOEPLITZ MULTIPLICATION

- **Cost** $O(n \log \text{band})$
(vs $O(n \log n)$)
- **Embarrassingly parallel**
- **Can select 2^k size blocks**
- **FFTW more efficient for $n = 65536$ than $n = 10^9$**
- **FFTW on consecutive vectors faster than consecutive FFTs**
- **Experimental result:**
Optimal Block Size ≈ 65536 for bandwidth of $T = 10000$.

THE POWER SPECTRUM ESTIMATION

$$\text{maximize } f(c) = d^T \cdot D \cdot d + \text{Tr}[\ln D]$$

$$D = \mathcal{N} + S = \mathcal{N} + c_1 S_1 + c_2 S_2 + \dots + c_l S_l$$

c_i - multipole powers we want

$$l = 20..1000$$

d - pixel map

\mathcal{N} - noise correlations of the “cut map”:

$$N = \left[\begin{array}{c|c} \mathcal{N} & * \\ \hline * & N_c \end{array} \right]$$

- **Newton's Method:**

$$\frac{\partial f}{\partial c_i} = \frac{1}{2} \left(d^T D^{-1} S_i D^{-1} d - \text{Tr} [D^{-1} S_i] \right)$$

$$\frac{\partial^2 f}{\partial c_i \partial c_j} = -d^T D^{-1} S_i D^{-1} S_j D^{-1} d + \frac{1}{2} \text{Tr} [D^{-1} S_i D^{-1} S_j]$$

NEWTON'S METHOD CONTINUED

- **Newton's Method:**

$$\frac{\partial f}{\partial c_i} = \frac{1}{2} \left(d^T D^{-1} S_i D^{-1} d - \text{Tr} [D^{-1} S_i] \right)$$

$$\frac{\partial^2 f}{\partial c_i \partial c_j} = -d^T D^{-1} S_i D^{-1} S_j D^{-1} d + \frac{1}{2} \text{Tr} [D^{-1} S_i D^{-1} S_j]$$

- **We need quantities like $D^{-1}z$ and $\text{Tr}(D^{-1}S_i)$, $\text{Tr}(D^{-1}S_i D^{-1}S_j)$**
- **Monte Carlo Estimator for $\text{Tr}(A)$**

$$\text{Tr}(A) = \langle y^T A y \rangle, \quad y = \{\pm 1\}, \text{ uniformly distributed}$$

- **Works for $\text{Tr}(D^{-1}S) = \langle y D^{-1}(S_i y) \rangle$**
- **Rough $\text{Tr}(A)$ OK**
- **All comes down to $D^{-1}z$, i.e. solving**

$$(\mathcal{N} + S)x = z \quad \text{i.e.} \quad (\mathcal{N} + c_1 S_1 + \dots + c_l S_l)x = z$$

- **(Preconditioned) Conjugate Gradients**
- **$S_i x$ in $O(n \log^2 n)$ - on processor $i = 1, \dots, l$**
- **$\mathcal{N}x$... we don't even have N or N^{-1} explicitly**

SOLVING $(\mathcal{N} + S)x = z$

$$N = \left[\begin{array}{c|c} \mathcal{N} & * \\ \hline * & N_c \end{array} \right]$$

- If we don't cut maps ($\mathcal{N} = N$) solve

$$(I + N^{-1}S)x = N^{-1}z$$

using PCG. Easy

- Computing $\mathcal{N}x$ if we cut maps using just $N^{-1} = A^T T A$:

- We solve $N^{-1} \begin{bmatrix} u \\ v \end{bmatrix} = \begin{bmatrix} x \\ 0 \end{bmatrix}$ using PCG
(same as in the map-making)

- Equivalent to:

$$\begin{bmatrix} u \\ v \end{bmatrix} = N \cdot \begin{bmatrix} x \\ 0 \end{bmatrix} = \left[\begin{array}{c|c} \mathcal{N} & * \\ \hline * & N_c \end{array} \right] \begin{bmatrix} x \\ 0 \end{bmatrix}$$

- Therefore $u = \mathcal{N}x$

- Convergence reasonable $\text{cond}(\mathcal{N} + S) \approx 10^4$, 200 or so iterations
- no known preconditioner for $(\mathcal{N} + S)$

CONCLUSIONS

- Overall cost $O(N_t \log \text{band} \cdot l^2)$... almost optimal provided PCGs converge fast
- Can handle 10^9 timesamples in days vs hundreds of years
- Recommendations for Scan Strategies:
 - Observe all pixels the same number of times
 - * Good or bad!
 - * Map making OK either way
 - * Keeps $\text{cond}(\mathcal{N} + S)$ small
 - Design instruments with smaller cross-corellation of timesamples (i.e. make T more diagonal)

FUTURE WORK:

- Fast spherical harmonic transforms

SLIDES AT:

<http://www.math.berkeley.edu/~plamen/pisa.ps>