

# Matrices With Displacement Structure

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# INTRODUCTION

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- **Structured Matrices**

- Dense
- Only depend on  $O(n)$  parameters
- Want  $LU$ , inverse and solution of  $Ax = b$  faster than  $O(n^3)$  and Matrix-vector multiply faster than  $O(n^2)$ .

- **Examples:**

- Cauchy  $C_{ij} = \frac{1}{x_i - y_j}$ , par.  $x_1, \dots, x_n, y_1, \dots, y_n$ .
- Vandermonde  $V_{ij} = x_i^{j-1}$
- Toeplitz  $T_{ij} = t_{i-j}$

- **Inverses of these Matrices do not have same structure but still "structured" in some sense.**

- **Toeplitz matrix**

$$\begin{bmatrix} t_0 & t_1 & t_2 & \cdots & t_{n-1} \\ t_{-1} & t_0 & t_1 & \cdots & t_{n-2} \\ t_{-2} & t_{-1} & t_0 & \cdots & t_{n-3} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ t_{-n+1} & t_{-n+2} & t_{-n+3} & \cdots & t_0 \end{bmatrix} - \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & t_0 & t_1 & \cdots & t_{n-2} \\ 0 & t_{-1} & t_0 & \cdots & t_{n-3} \\ 0 & \vdots & \vdots & \ddots & \vdots \\ 0 & t_{-n+2} & t_{-n+3} & \cdots & t_0 \end{bmatrix} = \begin{bmatrix} t_0 & t_1 & t_2 & \cdots & t_{n-1} \\ t_{-1} & 0 & 0 & \cdots & 0 \\ t_{-2} & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ t_{-n+1} & 0 & 0 & \cdots & 0 \end{bmatrix}$$

- **We can recover  $T$  from the Right Hand Side**

- **$T - Z_0 \cdot T \cdot Z_0^T = \text{Rank-2-Matrix}$ , where**

$$Z_\phi = \begin{bmatrix} 0 & 0 & 0 & 0 & \phi \\ -1 & 0 & 0 & \cdots & 0 \\ 0 & -1 & 0 & \cdots & 0 \\ 0 & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

- **Better yet to write  $Z_1 \cdot T - T \cdot Z_1 = G \cdot B$ , because  $Z_1$  - diagonalizable.**

# Definition of Displacement Structure

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- Given a matrix  $R$  we have displacement operator

$$\nabla_{F,A}(R) = FR - RA$$

if  $\nabla_{F,A}(R)$  has low rank compared to  $n$  we say that  $R$  has low displacement rank with respect to the operator  $\nabla_{F,A}$  and we write

$$\nabla_{F,A}(R) = FR - RA = GB$$

where  $G$  is  $n \times \alpha$  and  $B$  is  $\alpha \times n$ .

- For Toeplitz matrix  $T$  we have  $\text{rank}(\nabla_{Z_1, Z_{-1}}(R)) = 2$
- All matrices with low  $\{Z_1, Z_{-1}\}$ -rank are called **TOEPLITZ-LIKE**
- Many possible choices of  $F, A$  for a particular matrix, only few useful:
  - Need to be able to recover each entry of  $R$  from  $F, A, G, B$  in  $O(1)$  operations per entry
  - Either  $F$  and  $A^*$  - lower triangular, or diagonalizable (true for all basic classes)

# Basic classes of Structured Matrices

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Toeplitz-like:	$F = Z_1$	$A = Z_{-1}$
Toeplitz-plus-Hankel-like:	$F = Y_{00}$	$A = Y_{11}$
Cauchy-like:	$F = \text{diag}(c_1, \dots, c_n)$	$A = \text{diag}(d_1, \dots, d_n)$
Vandermonde-like:	$F = \text{diag}(\frac{1}{x_1}, \dots, \frac{1}{x_n})$	$A = Z_1$
Chebyshev-Vandermonde-like:	$F = \text{diag}(x_1, \dots, x_n)$	$A = Y_{\gamma, \delta}$
Polynomial-Vandermonde-like:	$F = \text{diag}(x_1, \dots, x_n)$	$A = \textit{ConfederateMatrix}.$

where

$$Z_\phi = \begin{bmatrix} 0 & 0 & \cdots & 0 & \phi \\ 1 & 0 & \cdots & \cdots & 0 \\ 0 & 1 & \ddots & & \vdots \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \end{bmatrix}, \quad *Y_{\gamma, \delta} = \begin{bmatrix} \gamma & 1 & 0 & \cdots & 0 \\ 1 & 0 & 1 & \ddots & \vdots \\ 0 & 1 & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & 0 & 1 \\ 0 & \cdots & 0 & 1 & \delta \end{bmatrix},$$

**i.e.  $Z_\phi$  is the lower shift  $\phi$ -circulant matrix and  $Y_{\gamma, \delta} = Z_0 + Z_0^T + \gamma e_1 e_1^T + \delta e_1 e_1^T.$**

# Displacement Structure is Inherited During

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- Inversion  $FR - RA = GB$  implies  $AR^{-1} - R^{-1}F = (-R^{-1}G)(BR^{-1})$
- Similarity transformations on  $F, A$ :

$$T_1 F_1 T_1^{-1} R - R T_2 A_2 T_2^{-1} = GB$$

implies

$$F_1 (T_1^{-1} R T_2) - (T_1^{-1} R T_2) A_2 = (T_1^{-1} G) (B T_2)$$

- Schur Complementation

Let  $R = \begin{bmatrix} d_1 & u_1 \\ l_1 & R_{22}^{(1)} \end{bmatrix}$  satisfy

$$\nabla_{F_1, A_1}(R_1) = \begin{bmatrix} f_1 & 0 \\ & F_2 \end{bmatrix} \cdot R_1 - R_1 \cdot \begin{bmatrix} a_1 & * \\ 0 & A_2 \end{bmatrix} = G_1 \cdot B_1$$

$$(G_1 \in C^{n \times \alpha}, B \in C^{\alpha \times n}).$$

If  $d_1 \neq 0$  the Schur complement  $R_2 = R_{22}^{(1)} - \frac{l_1 u_1}{d_1}$  satisfies the displacement equation

$$F_2 \cdot R_2 - R_2 \cdot A_2 = G_2 \cdot B_2,$$

where

$$\begin{bmatrix} 0 \\ G_2 \end{bmatrix} = G_1 - \begin{bmatrix} 1 \\ \frac{1}{d_1} l_1 \end{bmatrix} \cdot g_1, \quad \begin{bmatrix} 0 & B_2 \end{bmatrix} = B_1 - b_1 \cdot \begin{bmatrix} 1 & \frac{1}{d_1} u_1 \end{bmatrix},$$

where  $g_1$  and  $b_1$  are the first row of  $G_1$  and the first column of  $B_1$ , respectively.

- Pivoting, if  $F$  is diagonal. Let  $FR - RA = GB, P$ -perm.

$$(PFP^T)(PR) - (PR)A = (PG)B$$

$PFP^T$  is still diagonal and pivoting translates into swapping 2 diagonal entries of  $F$  and 2 rows of  $G$ .

## Fast $O(n^2)$ inversion

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We can solve  $Rx = b$  in  $O(n^2)$  time if  $FR - RA = GB$ .  
The inverse satisfies

$$AR^{-1} - R^{-1}F = -(R^{-1}G)(BR^{-1})$$

- Solve  $Rx = g_i$  where  $g_i$  are the columns of  $G$ . This is  $R^{-1}G$
- Solve  $R^T x = b_i^T$  where  $b_i$  are the rows of  $B$ . The solutions are  $R^{-T}B^T$ . Transpose and we get  $BR^{-1}$ .
- Then we know the  $\{A, F\}$ -generators of  $R^{-1}$  and can recover  $R^{-1}$  in  $O(n^2)$  time
- Total cost =
  - Solving  $2\alpha$  linear equations with  $\alpha$ -small =  $O(n^2)$ .
  - $O(n^2)$  for recovery of  $R^{-1}$  from its displacement equation.
  - Total Cost =  $O(n^2)$ .

# Other Classes of Structured Matrices

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where

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i.e.  $Z_\phi$  is the lower shift  $\phi$ -circulant matrix and  $Y_{\gamma, \delta} = Z_0 + Z_0^T + \gamma e_1 e_1^T + \delta e_1 e_1^T$ .

- **All  $F$ 's and  $A$ 's diagonal or diagonalizable**  
(using Fast Trigonometric Transforms, diagonal scaling and products thereof.)
- **We can transform  $R$  into Cauchy-like matrix**
- **Then we can apply Fast GEPP**
- **If  $ARB$  is Cauchy-like and  $ARB = PLU$**   
then  $R = A^{-1}PLUB^{-1}$ .
- **Applying  $A$  and  $B$  to the right hand side of  $Rx = b$  costs**  
 $\leq O(n \log n)$
- **Overall cost of solving  $Rx = b$  is still  $O(n^2)$ .**