Accurate and Efficient Matrix Computations with Totally Positive Generalized Vandermonde Matrices Using Schur Functions

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GOALS

- Accurate (Small relative error) and Efficient $(O(n^3)$ or perhaps $O(n^p)$, independent of condition number) Linear Algebra
 - $-A^{-1}$
 - -Ax = b
 - LDU from GENP, GEPP, GECP
 - -SVD
- Can't be done for general matrices, must be "structured"
 - Certain sparsity patterns
 - Cauchy
 - Vandermonde
 - **—** ...
- Goal of this talk: Accurate and Efficient Linear Algebra for Generalized Vandermonde Matrices

| | | | | | | | Small Forward | Small Backward |
|------------------|-----------|----------|-------|------|------|-----|------------------|-------------------|
| Type of | | | Any | GENP | | | Error in | Error in |
| Matrix | $\det(A)$ | A^{-1} | minor | GEPP | GECP | SVD | Ax = b | Ax = b |
| Cauchy | | | | | | | | |
| | | | | | | | | |
| Totally Positive | | | | | | | | |
| Cauchy | | | | | | | | |
| Vandermonde | | | | | | | | |
| | | | | | | | | |
| Totally positive | | | | | | | | |
| Vandermonde | | | | | | | | |
| Polynomial | | | | | | | | |
| Vandermonde | | | | | | | | |
| Poly. Vand. | | | | | | | | |
| Orth. poly. | | | | | | | | |
| Generalized | | | | | | | | |
| Vandermonde | | | | | | | | |
| TP Generalized | | | | | | | | |
| Vandermonde | | | | | | | | |

Totally Positive = Matrix with all minors > 0

OUTLINE

- Model of arithmetic
- Classical method for achieving the goals for simple examples The Björck-Pereyra Method for Vandermonde Matrices
- How and why it works?
- ullet Application to TP Generalized Vandermonde matrices

How can we lose accuracy in computing in floating point?

- $f(a \otimes b) = (a \otimes b)(1 + \delta)$ model of arithmetic with no over/underflow
- OK to multiply, divide, add positive numbers Proof: $1 + \delta$ factors can be factored out
- $x_i \pm x_j$, where x_i and x_j are initial data (so exact)
- $\bullet (x_i + y_j)(x_i y_{j-1})x_{i+1}/(x_{i-1} y_j) \mathbf{OK}$
- Cancellation when subtracting approximate results dangerous:

• We will compute everything using only allowable expressions

Classical Example: A Vandermonde Linear System

• Solve Vy = b, where V is Vandermonde:

$$\begin{bmatrix} 1 & x_1 & \dots & x_1^{n-1} \\ 1 & x_2 & \dots & x_2^{n-1} \\ 1 & x_3 & \dots & x_3^{n-1} \\ & \ddots & & \\ 1 & x_n & \dots & x_n^{n-1} \end{bmatrix} \cdot \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_n \end{bmatrix} = \begin{bmatrix} + \\ - \\ + \\ \vdots \\ - \end{bmatrix}$$

and $0 < x_1 < \ldots < x_n$.

- Equivalent to interpolation
- The Björck-Pereyra method solves Vy = b
 - In $O(n^2)$ time
 - With small forward error: $|y_i \hat{y}_i| \leq O(\epsilon)|y_i|$
 - With small backward error: If $\hat{V}\hat{y} = b$ then $|V_{ij} \hat{V}_{ij}| \leq O(\epsilon)|V_{ij}|$.
- How does it work?

The Björck-Pereyra Method

• If $(x_1, x_2, x_3) = (1, 2, 3)$ and $b = (2, -1, 14)^T$ then using BP to solve

$$\begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 2^2 \\ 1 & 3 & 3^2 \end{bmatrix} \cdot y = \begin{bmatrix} 2 \\ -1 \\ 14 \end{bmatrix}$$
 means

$$y = V^{-1}b = \begin{bmatrix} 1 & -1 \\ & 1 & -1 \\ & & 1 \end{bmatrix} \cdot \begin{bmatrix} 1 \\ & 1 & -2 \\ & & 1 \end{bmatrix} \cdot \begin{bmatrix} 1 \\ & 1 \\ & -\frac{1}{2} & \frac{1}{2} \end{bmatrix} \cdot \begin{bmatrix} 1 \\ & -1 & 1 \\ & & -1 & 1 \end{bmatrix} \cdot \begin{bmatrix} 2 \\ & -1 \\ & 14 \end{bmatrix} = \begin{bmatrix} 23 \\ & -30 \\ & 9 \end{bmatrix}$$

• Notice:

- Bidiagonal Decomposition of V^{-1} (accurate)
- Checkerboard sign pattern
- \Rightarrow No subtractive cancellation
- \Rightarrow High relative accuracy

• Questions:

- Which matrices have bidiagonal decomposition of their inverses?
- Checkerboard signs?
- Accurate?

The Björck-Pereyra Method Dissected

• Questions:

- Which matrices have bidiagonal decomposition of their inverses?
- Checkerboard signs?
- Accurate?

• Answers:

All nonsingular matrices do
 This is Neville elimination in matrix form:

$$\begin{bmatrix} 1 \\ -1 & 1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 4 \\ 1 & 3 & 9 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 3 \\ 0 & 1 & 5 \end{bmatrix}; \qquad \begin{bmatrix} 1 \\ 1 \\ -\frac{1}{2} & \frac{1}{2} \end{bmatrix} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 3 \\ 0 & 1 & 5 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 3 \\ 0 & 0 & 1 \end{bmatrix}$$
$$\begin{bmatrix} 1 & -1 \\ 1 & -1 \\ 1 & 1 \end{bmatrix} \cdot \begin{bmatrix} 1 \\ 1 \\ 1 -\frac{1}{2} & \frac{1}{2} \end{bmatrix} \cdot \begin{bmatrix} 1 \\ 1 \\ -1 & 1 \\ 1 & 3 \end{bmatrix} \cdot \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 4 \\ 1 & 3 & 9 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

- Checkerboard sign pattern \iff Total positivity $(A \text{ is TP} \iff \text{all minors} > 0)$
- Accurate? Yes.

ACCURACY OF THE BJÖRCK-PEREYRA METHOD

$$\begin{pmatrix}
\begin{bmatrix} 1 & x_1 & x_1^2 & x_1^3 \\ 1 & x_2 & x_2^2 & x_2^3 \\ 1 & x_3 & x_3^2 & x_3^3 \\ 1 & x_4 & x_4^2 & x_4^3
\end{bmatrix}
\end{pmatrix}^{-1} = \begin{bmatrix} 1 & -x_1 & & & \\ & 1 & -x_1 & & \\ & & 1 & -x_1 \\ & & & 1
\end{bmatrix}
\begin{bmatrix} 1 & & & \\ & 1 & -x_2 \\ & & 1 & -x_2 \\ & & & 1
\end{bmatrix}
\begin{bmatrix} 1 & & & \\ & 1 & -x_3 \\ & & & 1
\end{bmatrix} \times$$

$$\begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & \frac{-1}{x_4 - x_1} & \frac{1}{x_4 - x_1} \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & \frac{-1}{x_3 - x_1} & \frac{1}{x_3 - x_1} & \\ & & \frac{-1}{x_4 - x_2} & \frac{1}{x_4 - x_2} \end{bmatrix} \begin{bmatrix} 1 & & & \\ \frac{-1}{x_2 - x_1} & \frac{1}{x_2 - x_1} & \\ & & \frac{-1}{x_3 - x_2} & \frac{1}{x_3 - x_2} \\ & & & \frac{-1}{x_4 - x_3} & \frac{1}{x_4 - x_3} \end{bmatrix}$$

Other TP matrices? ... Yes

TP Cauchy matrices $x_1 > \ldots > x_n > y_1 > \ldots > y_n$

$$\left(\begin{bmatrix} \frac{1}{x_1 - y_1} & \frac{1}{x_1 - y_2} & \frac{1}{x_1 - y_3} \\ \frac{1}{x_2 - y_1} & \frac{1}{x_2 - y_2} & \frac{1}{x_2 - y_3} \\ \frac{1}{x_3 - y_1} & \frac{1}{x_3 - y_2} & \frac{1}{x_3 - y_3} \end{bmatrix} \right)^{-1} = \begin{bmatrix} 1 & \frac{-(x_1 - y_1)}{y_1 - y_2} & 0 \\ 0 & \frac{x_1 - y_2}{y_1 - y_2} & \frac{-(x_1 - y_2)}{y_2 - y_3} \\ 0 & 0 & \frac{x_1 - y_3}{y_2 - y_3} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & \frac{-(x_2 - y_1)}{y_1 - y_3} \\ 0 & 0 & \frac{x_2 - y_3}{y_1 - y_3} \end{bmatrix} \begin{bmatrix} x_1 - y_1 \\ x_2 - y_2 \\ x_3 - y_3 \end{bmatrix} \times$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & \frac{-(x_1 - y_2)}{x_3 - x_1} & \frac{x_3 - y_2}{x_3 - x_1} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ \frac{-(x_1 - y_1)}{x_2 - x_1} & \frac{x_2 - y_1}{x_2 - x_1} & 0 \\ 0 & \frac{-(x_2 - y_1)}{x_3 - x_2} & \frac{x_3 - y_1}{x_3 - x_2} \end{bmatrix}$$

Unifying Characteristic?

The Connection with Minors

- Which TP matrices permit *accurate* bidiagonal decomposition?
- Each entry is *product of quotients of minors*

$$L_{i+1,i}^{(k)} = -\frac{\det(A(i-k+2:i+1,1:k))}{\det(A(i-k+2:i,1:k-1))} \cdot \frac{\det(A(i-k+1:i-1,1:k-1))}{\det(A(i-k+1:i,1:k))}$$

- Specifically: Initial minors
 - Contiguous
 - Include first row and column
- Initial minors of Cauchy:

$$\det(C) = \frac{\prod_{i < j} (x_j - x_i)(y_j - y_i)}{\prod_{i,j} (x_i + y_j)}$$

• Initial minors of Vandermonde:

$$\det V = \prod_{i>j} (x_i - x_j)$$

- How did we think of minors?
- Gaussian Elimination and Neville Elimination Each entry of V = LDU is a quotient of minors, so not surprising

New results: Generalized Vandermonde Matrices

• TP Matrices with initial minors that are easy to compute accurately Vandermonde and Generalized Vandermonde

$$V = \begin{bmatrix} 1 & x_1 & \dots & x_1^{n-1} \\ 1 & x_2 & \dots & x_2^{n-1} \\ & \ddots & & \\ 1 & x_n & \dots & x_n^{n-1} \end{bmatrix}, \qquad G_{\lambda} = \begin{bmatrix} x_1^{\lambda_1} & x_1^{1+\lambda_2} & \dots & x_1^{n-1+\lambda_n} \\ x_2^{\lambda_1} & x_2^{1+\lambda_2} & \dots & x_2^{n-1+\lambda_n} \\ & & \ddots & & \\ x_n^{\lambda_1} & x_n^{1+\lambda_2} & \dots & x_n^{n-1+\lambda_n} \end{bmatrix},$$

where $x_1 > x_2 > \cdots > x_n > 0$, $\lambda_n \geq \lambda_{n-1} \geq \cdots \geq \lambda_1 \geq 0$, $|\lambda| = \lambda_1 + \ldots + \lambda_n$

• Initial Minors for G_{λ} ?

$$\det(G_{\lambda}) = \det(V) \cdot s_{\lambda}(x_1, x_2, \dots, x_n)$$

- s_{λ} called Schur function
 - Polynomial with positive integer coefficients
 - Widely studied in combinatorics [MacDonald], group representation theory
- Example:

$$\det \left(\begin{bmatrix} 1 & x_1^2 & x_1^4 \\ 1 & x_2^2 & x_2^4 \\ 1 & x_3^2 & x_3^4 \end{bmatrix} \right) = \det \left(\begin{bmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \end{bmatrix} \right) \cdot (2x_1x_2x_3 + x_1^2x_2 + x_1x_2^2 + x_1^2x_3 + x_1x_3^2 + x_2^2x_3 + x_2x_3^2)$$

Accuracy and Efficiency for Generalized Vandermonde Matrices

• Example:

$$\det \left(\begin{bmatrix} 1 & x_1^2 & x_1^4 \\ 1 & x_2^2 & x_2^4 \\ 1 & x_3^2 & x_3^4 \end{bmatrix} \right) = \det \left(\begin{bmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \end{bmatrix} \right) \cdot (2x_1x_2x_3 + x_1^2x_2 + x_1x_2^2 + x_1^2x_3 + x_1x_3^2 + x_2^2x_3 + x_2x_3^2)$$

• Accuracy?

$$-\det(V) = \prod_{i>j} (x_i - x_j) - YES.$$

 $-s_{\lambda}$ - polynomials with > 0 coefficients - YES.

• Efficiency?

$$-\det(V) = \prod_{i>j} (x_i - x_j) - \mathbf{OK}.$$

$$-s_{\lambda}(x_1,x_2,\ldots,x_n)$$
?

- * Traditional algorithm exponential $-n^{|\lambda|}$
- * Now exponential speedup: Linear complexity in n. Idea:

$$s_{(2)}(x_1, \dots, x_n) = \sum_{i < j} x_i x_j = (x_1 + \dots + x_n) x_1 + (x_2 + \dots + x_n) x_2 + \dots + (x_{n-1} + x_n) x_{n-1} + x_n x_n$$

cost: 3n, although n^2 terms.

| Type of | | | Any | GENP | | | | Ax = b | Ax = b |
|---------------------------|-------------------|---------------------|----------------|---------------|-------|-------|---------------|---------------|---------------|
| Matrix | $\det(A)$ | A^{-1} | minor | GEPP | GECP | SVD | NENP | Frwrd* | Bckwrd* |
| Cauchy | n^2 | n^2 | n^2 | n^3 | n^3 | n^3 | n^2 | n^2 | |
| | | | | | | | | | |
| TP Cauchy | n^2 | n^2 | n^2 | n^3 | n^3 | n^3 | n^2 | n^2 | n^2 |
| | | | | | | | | | |
| Vandermonde | n^2 | | | | | n^3 | n^2 | n^2 | |
| | | | | | | | | | |
| TP | n^2 | n^3 | EXP | n^3 | EXP | n^3 | n^2 | n^2 | n^2 |
| Vandermonde | | | | | | | | | |
| Polynomial | | | | | | | | | |
| Vandermonde | n^2 | | | n^3 | | n^3 | | | |
| Orth. Poly. | | | | | | | | | |
| Poly. Vand. | | | | | | | | | |
| Orth. poly. ¹⁾ | n^2 | n^3 | \mathbf{EXP} | n^3 | EXP | n^3 | | | |
| $0 < x_1 < \ldots < x_n$ | | | | | | | | | |
| Generalized | | | | | | | | | |
| Vandermonde | | | | | | | | | |
| TP Generalized | $\Lambda n + n^2$ | $\Lambda n^2 + n^3$ | EXP | Λn^2 | EXP | EXP | Λn^2 | Λn^2 | Λn^2 |
| Vandermonde | | | | | | | | | |

Big-O sense

*FORWARD BOUND: $|x - \hat{x}| \leq O(\epsilon)|A^{-1}||b|$, implying $|x - \hat{x}| \leq O(\epsilon)|x|$ for x checkerboard BACKWARD BOUND: $|A - \hat{A}| \leq O(\epsilon)|A|$, where $\hat{A}\hat{x} = b$.

 $^{1)}$ + Other conditions on the signs of the three-term recurrence

$$\Lambda \leq (\lambda_1 + 1)(\lambda_2 + 1)^2 \dots (\lambda_p + 1)^2 p$$
, where $\lambda = (\lambda_1, \dots, \lambda_p)$.

Conclusions

- TP Structured linear systems can be solved very accurately, if initial minors factor
- Implies accurate A^{-1}
- New application: Generalized Vandermonde Matrices
- Accurate SVD of some Polynomial Vandermonde Matrices
- Sometimes the SVD is easier than the inverse

Open Problems

- Totally Positive Matrices in general appear impossible. Proof?
- Characterize which structured matrices permit accurate and efficient linear algebra

Resources

• These slides: www.math.berkeley.edu/~plamen/bascd02.pdf

• Reports:

- J. Demmel and P. Koev, Necessary and sufficient conditions for accurate and efficient rational function evaluation and factorizations of rational matrices. In *Structured matrices in mathematics, computer science, and engineering. II (Boulder, CO, 1999)*, pages 117–143. Amer. Math. Soc., Providence, RI, 2001.
- J. Demmel and P. Koev, Accurate and Efficient Matrix Computations with Totally Positive Generalized Vandermonde Matrices Using Schur Functions, www.math.berkeley.edu/~plamen/hagen.ps
- J. Demmel, M. Gu, S. Eisenstat, I. Slapničar, K. Veselić, and Z. Drmač. Computing the singular value decomposition with high relative accuracy. *Lin. Alg. Appl.*, 299(1-3):21-80, 1999.

 www.cs.berkeley.edu/~demmel/DGESVD.ps
- J. Demmel, Accurate SVDs for Structured Matrices, Accurate SVDs of structured matrices. SIAM J. Mat. Anal. Appl., 21(2):562-580, 1999. www.netlib.org/lapack/lawns/lawn130.ps