

Lecture # 4

- Outline:
- 1) Review from last time
 - 2) Colorings of Maps Theorem
 - 3) The connect sums of surfaces
 - 4) Oriented surfaces and curves

Section 1 : Review

Definition : A surface is space that locally looks like \mathbb{R}^2
↳ ie, zoom in close it just looks like a "piece of paper."

Definition : A polygonal complex is a space obtained by gluing together polygons, edges, and vertices, where by glue we mean that we identify edges w/ edges and vertices w/ vertices (could glue polygon to self)

Definition: Let $X =$ polygonal complex w/

- $V(X) = \#$ of vertices
- $E(X) = \#$ of edges
- $F(X) = \#$ of faces

The Euler characteristic of X is

$$\chi(X) = V(X) - E(X) + F(X)$$

Examples: 1) Random Polygonal Complex

$$\chi \left(\text{Diagram} \right) = 11 - 14 + 3 = 0$$

2) Torus

$$\chi \left(\text{Diagram} \right) = 1 - 2 + 1 = 0$$

3) Sphere

$$\chi \left(\text{Diagram} \right) = 3 - 2 + 1 = 2$$

Proposition: Let X and Y be polygonal complexes that are homeomorphic to the same surface. Then their Euler characteristics agree.

$$\chi(X) = \chi(Y)$$

Definition: The Euler characteristic of a surface Σ is the Euler characteristic of any polygonal cpx that is homeomorphic to Σ .

Remark: To compute $\chi(\Sigma)$, break Σ up into regions and count the # of vertices, edges, and faces.

Examples:

$$1) \chi(S^2) = 2$$

$$2) \chi(T^2) = 0$$

$$3) \chi(\text{Klein bottle}) = 0$$

$$4) \chi(\text{genus } g \text{ surface}) = 2 - 2g.$$

↳ g donut holes

Section 2: Colorings of Maps (4 Colors Theorem)

- Question:
- What is the minimum number of colors needed to color any map of the globe so that no two adjacent regions are colored the same color?
 - What is the minimum number of colors needed to color any map of a surface so that no two adjacent regions are colored the same color?

Definition:

A surface is compact if it admits a polygonal complex structure w/ a finite # of vertices, edges, and faces.

↳ Secretly, we needed to assume that our surfaces were compact when we defined their Euler characteristics.

Example:

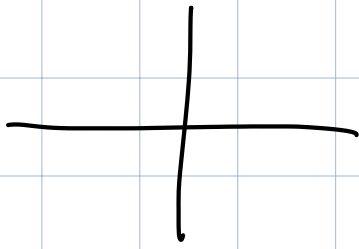
1) Sphere, torus, multiple donut holes

Compact



2)

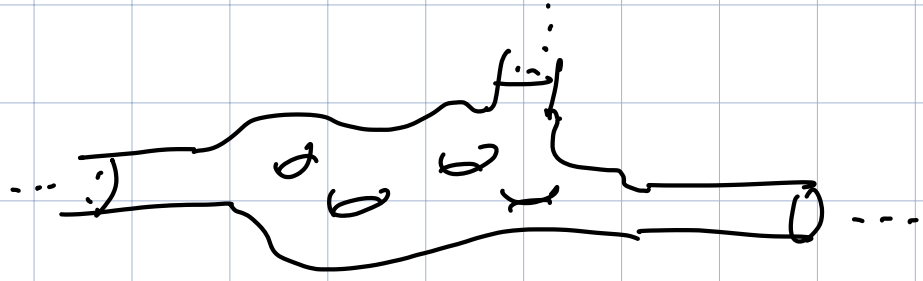
\mathbb{R}^2



non compact:

3)

Infinite Ends



Definition:

A geographic complex associated to a compact surface Σ is a polygonal complex that is homeomorphic to Σ and satisfies:

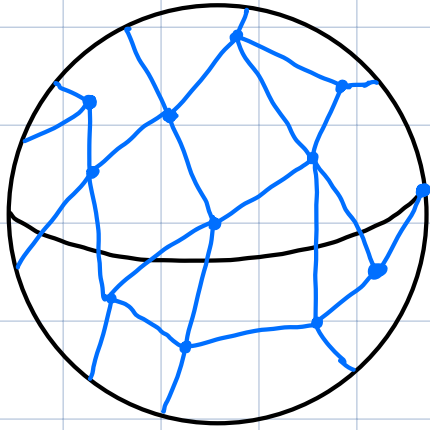
- 1) Every face does not meet itself
- 2) Any two faces that meet/share a unique edge
- 3) At least three faces meet at each vertex.

Remark:

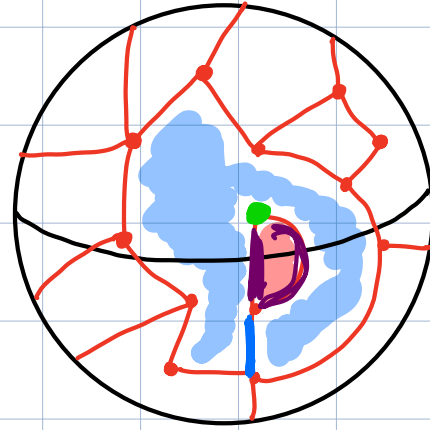
Intuitively, a geographic cpx is a map of the surface that satisfies

- 1) A region cannot boarder itself
- 2) Two regions can only share one unique boarder
- 3) Vertices are where 3 or more regions meet

Example :



Geo cp^x .



Not geo cp^x

- 1) A region cannot boarder itself
- 2) Two regions can only share one unique boarder
- 3) Vertices are where 3 or more regions meet

Definition: • A legal coloring of a geo. cpx. is an assignment of a color to each face st no two adjacent faces have the same color.

• The coloring number of a geo cpx X
 $N(X) =$ minimum # of colors needed to produce a legal coloring of X .

• The coloring number of a compact surface Σ is
 $N(\Sigma) =$ minimum # of colors needed to produce a legal coloring of all geo. cpx associated to Σ .

Theorem:

$$N(\Sigma) \leq \frac{7 + \sqrt{49 - 24 \cdot \chi(\Sigma)}}{2}$$

Corollary:

Any geo. cpx on Σ can be legally colored using

$$\frac{7 + \sqrt{49 - 24 \cdot \chi(\Sigma)}}{2}$$

colors.

Example:

1) $\Sigma = S^2 \Rightarrow$ need at most 4 colors

2) $\Sigma = T^2 \Rightarrow$ " " " 7 "

3) $\Sigma =$ genus 4 surface \Rightarrow need at most 10 colors

Remark:

To prove the theorem for $\Sigma = S^2$ is extremely difficult.

We will prove it for $\chi(\Sigma) \leq 1$.

Notation:

Let X be the geo. cpx associated to Σ that satisfies:

1) $N(X) = N(\Sigma)$

2) If Y is another geo. cpx associated to Σ w/ $N(Y) = N(X)$, then $F(X) \leq F(Y)$.

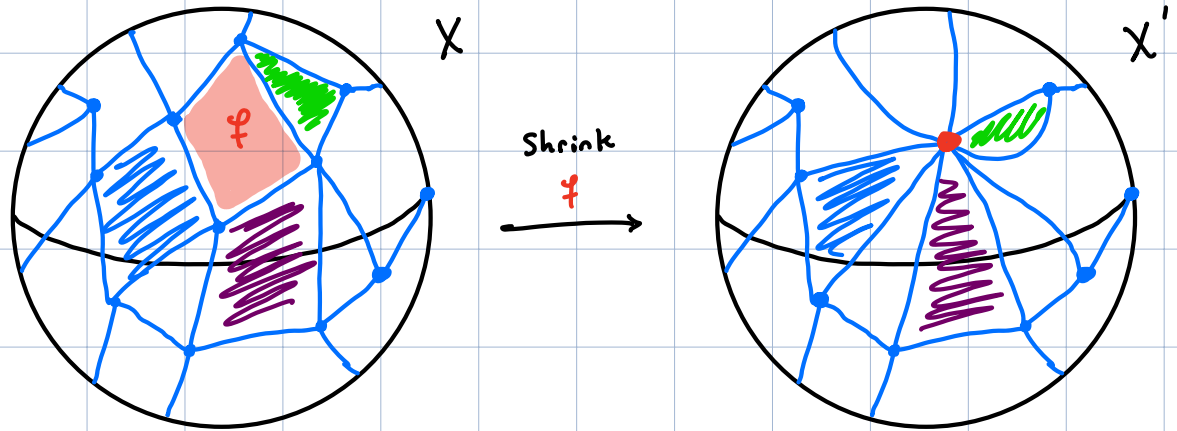
Lemma 0:

Every face of X has at least $N(X) - 1$ edges.

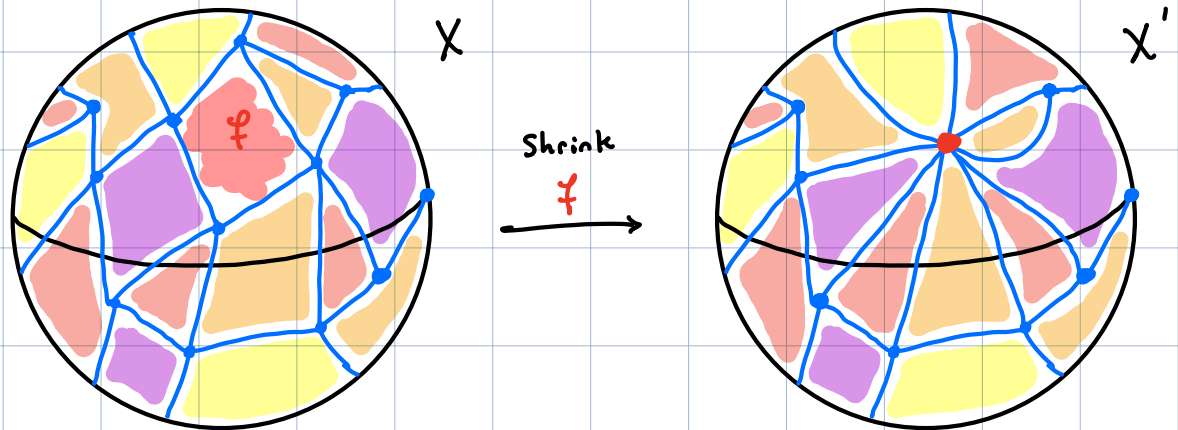
Proof:

- 1) We suppose by way of contradiction that there exists a face in X w/ strictly less than $N(X) - 1$ edges.
- 2) Denote this face by f .
- 3) Shrink f and all of its edges down to a single vertex.

This produces a new geo. cpx X' .



4) Since this procedure does not produce any new edges any coloring of X gives rise to a coloring of X'



5) So $N(X') \leq N(X)$

6) If $N(X') = N(X)$, then by assumption on X ,
 $F(X) \leq F(X') = F(X) - 1$

\Rightarrow we actually must have $N(X') < N(X)$

7) So we may color X' w/ $N(X) - 1$ colors.

But this allows us to color X w/ $N(X) - 1$ colors.

Namely, we color X' , then since f has less than $N(X) - 1$ edges, it has at most $N(X) - 2$ adjacent faces. So we can always pick one of the $N(X) - 1$ colors to color f differently than all its adjacent faces. $\Rightarrow N(X) \leq N(X) - 1$, a contradiction.

8) \Rightarrow Every face of X has at least $N(X) - 1$ edges

Lemma 1: $(N(X) - 1) \cdot F(X) \leq 2 \cdot E(X)$

$$N(X) \leq \frac{2E}{F} + 1$$

Proof:

1) Every edge touches two unique faces.

\Rightarrow Average # of edges per face is $2E(X)/F(X)$

2) By Lemma 0, each face has at least $N(X) - 1$ edges

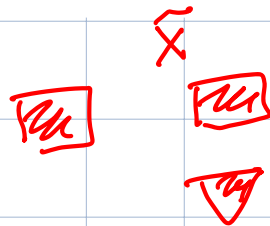
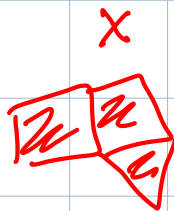
\Rightarrow Average # of edges per face $\geq N(X) - 1$

3) Combining these inequalities,

$$N(X) - 1 \leq 2E(X)/F(X)$$

Lemma 2 :

$$3V(X) \leq 2E(X)$$



Proof :

1) Let \tilde{X} = preglued collection of polygons that we glue together to produce X .

2) Note, $2E(X) = E(\tilde{X})$

3) Since at least 3 faces meet at each vertex,

$$3V(X) \leq V(\tilde{X})$$

4) Since \tilde{X} is disjoint collection of polygons,

$$E(\tilde{X}) = V(\tilde{X})$$

5) Combining,

$$2E(X) = E(\tilde{X}) = V(\tilde{X}) \geq 3V(X)$$

□

Lemma 3: $N(X) = N(\Sigma) \leq 7 - 6 \cdot \chi(\Sigma) / F(X)$

Proof:
$$\begin{aligned} N(\Sigma) &= N(X) \\ &\leq 1 + 2E(X) / F(X) && \uparrow \text{Lemma 1} \\ &\leq 1 + (6E(X) - 6V(X)) / F(X) && \uparrow \text{Lemma 2} \\ &= 1 + (-6\chi(X) + 6F(X)) / F(X) \\ &= 7 - 6 \cdot \chi(X) / F(X) \\ &= 7 - 6 \cdot \chi(\Sigma) / F(X) \end{aligned}$$

□

$$\begin{aligned} 6\chi &= 6V - 6E + 6F \\ 6E - 6V &= 6F - \chi \cdot 6 \end{aligned}$$

Proof: $(\chi(\Sigma) = 1) :$

$$N(\Sigma) \leq 7 - 6/F(x)$$

$$\leq 6$$

$$= \frac{7 + \sqrt{49 - 24 \cdot 1}}{2}$$

$$= \frac{7 + \sqrt{49 - 24 \cdot \chi(\Sigma)}}{2}$$

□

Proof: $(\chi(\Sigma) \leq 0)$

$$1) N(\Sigma) = N(x) \leq F(x)$$

$$2) N(\Sigma) \leq 7 - 6 \cdot \chi(\Sigma) / F(x)$$

$$\leq 7 - 6 \cdot \chi(\Sigma) / N(\Sigma)$$

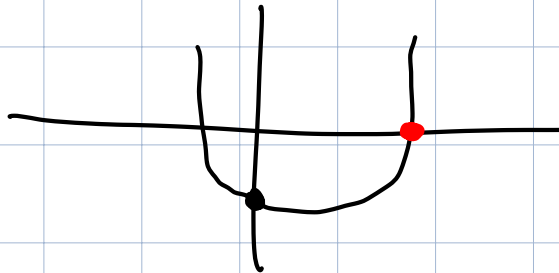
$$3) \Rightarrow N(\Sigma)^2 - 7N(\Sigma) + 6 \cdot \chi(\Sigma) \leq 0$$

4) This polynomial in $N(\Sigma)$ is upwards opening w/ at least one point on $N(\Sigma)$ -axis.

5) \Rightarrow Largest $N(\Sigma)$ for which this holds is largest zero of poly.

$$6) \Rightarrow N(\Sigma) \leq \frac{7 + \sqrt{49 - 24 \cdot \chi(\Sigma)}}{2}$$

□

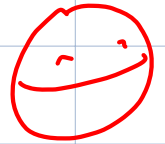
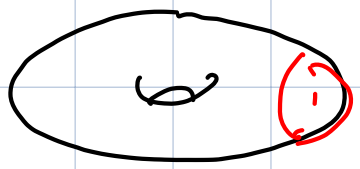


Section: The connect sum of surfaces

Definition: Given two surfaces X and Y , the connect sum of X and Y , denoted $X \# Y$, is obtained via

- 1) Remove an open disk from both X and Y to create two surfaces w/ "boundaries"
- 2) Glue the resulting boundaries together to create the new surface $X \# Y$.

Picture :



↓
cut out
disks



↓
glue along
boundaries



Example:

1) $T^2 \# T^2 = 2$ holed surface = genus 2 surface.

2) $S^2 \# S^2 = S^2$

3) $S^2 \# T^2 = T^2$

4) $T^2 \# \dots \# T^2 \left. \vphantom{\dots} \right\} g\text{-times} = \text{genus } g \text{ surface.}$

Proposition:

$$\chi(X \# Y) = \chi(X) + \chi(Y) - 2$$

Ex:

1) $\chi(T^2 \# T^2) = \chi(T^2) + \chi(T^2) - 2 = 0 + 0 - 2 = -2$

2) $2 = \chi(S^2) = \chi(S^2) + \chi(S^2) - 2 = 2 + 2 - 2 = 2.$ ✓

3) ✓

4) $\chi(T^2 \# \dots \# T^2) = 2 - 2g.$