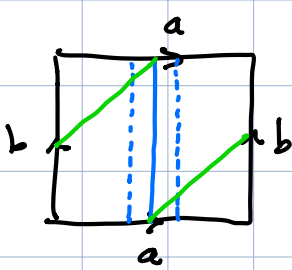


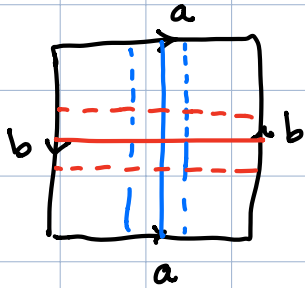
Ex:

i)



2-sided curve

ii)



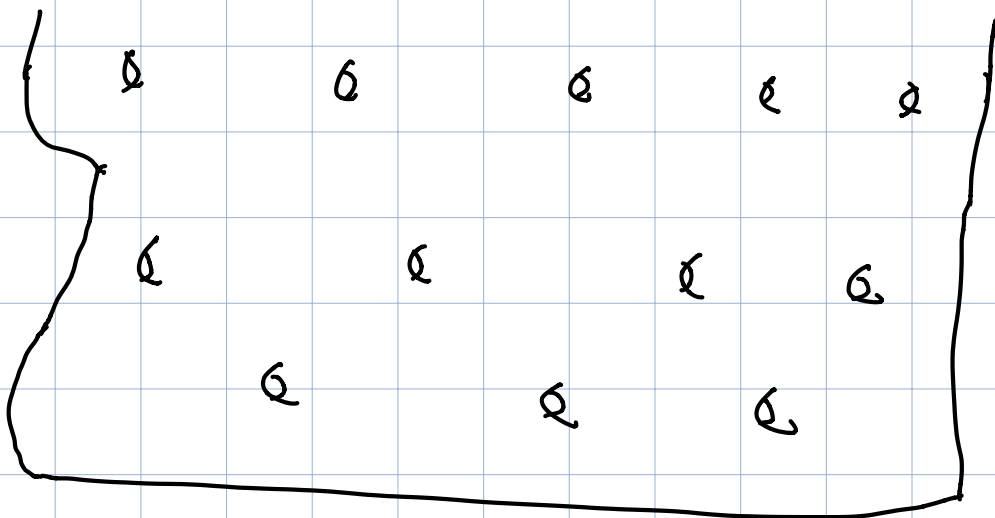
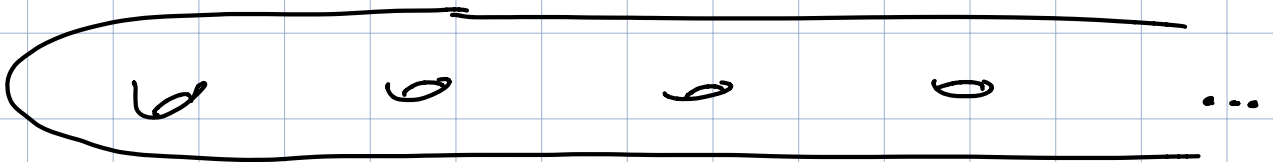
$MV = 2$ -sided curve

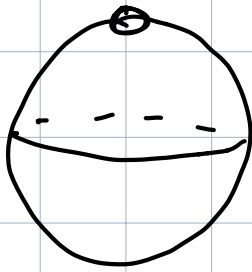
$MV = 1$ -sided curve

Def: A surf is orientable if it has no 1-sided curves.

Def: A surf is compact if a poly cpx associated to it has a finite # of polygons

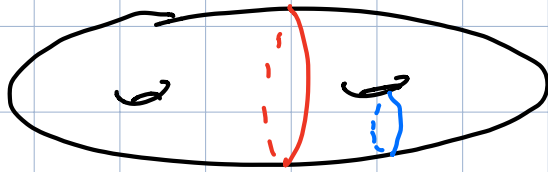
Ex:



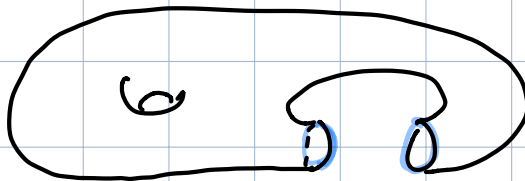


Rmk: If γ is a 2-sided curve in X that does not sep. X , then $X = T^2 \# Y$, where Y is some other surf.

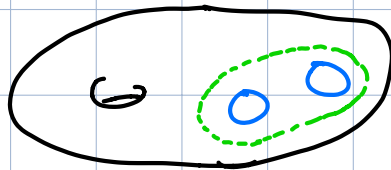
cut along it and X w/ the cut is still conn.



cut \rightsquigarrow



deform \rightarrow



reglue \rightsquigarrow



green w/ blue glued is a T^2 w/ a boundary comp

$\Rightarrow X$ is realized as a connect sum w/ T^2 along the green curve

Rmk: $X \# S^2 = X$

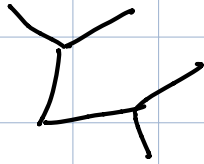
Lemma: $\Gamma = \text{conn graph}$, then

$$\chi(\Gamma) \leq 1$$

w/ $\chi(\Gamma) = 1$ if and only if $\Gamma = \text{tree}$.

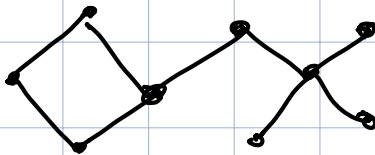
Exs

i)



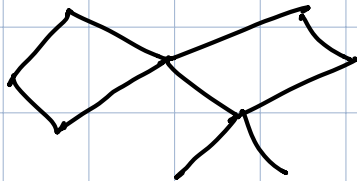
$$7 - 6 = 1 \quad \checkmark$$

ii)



$$9 - 9 = 0$$

iii)



$$9 - 10 = -1$$

Remark: Each loop contributes -1 to the χ .

Proof: Build Γ up (in a conn manner) in a seq.

$\Gamma_1, \Gamma_2, \Gamma_3, \dots, \Gamma_N$.

Compute $\chi(\Gamma_i)$ inductively

$$\Gamma_1 = \text{---} \Rightarrow \chi(\Gamma_1) = 1$$

$$\Gamma_2 = \left\{ \begin{array}{l} \text{---} \\ \text{---} \end{array} \right. \Rightarrow \chi(\Gamma_2) = 3 - 2 = 1$$

$$\Rightarrow \chi(\Gamma_2) = 2 - 2 = 0 \leq 1$$

Γ_{i+1} from Γ_i , we ^① either add 1 new edge and 1 new ver. or ^② just add 1 new edge.

$$\begin{aligned}
 \textcircled{1} \Rightarrow \chi(\Gamma_{i+1}) &= V(\Gamma_{i+1}) - E(\Gamma_{i+1}) \\
 &= V(\Gamma_i) + 1 - (E(\Gamma_i) + 1) \\
 &= V(\Gamma_i) - E(\Gamma_i) \\
 &= 1
 \end{aligned}$$

$$\textcircled{2} \Rightarrow \chi(\Gamma_{i+1}) = \chi(\Gamma_i) - 1$$

Note if we just do $\textcircled{1}$, then we create no loops
 $\Rightarrow \Gamma$ is a tree but we don't drop $\chi(\Gamma)$
 from 1 w/ this set-up.

Thm: (2-dim'l Poincaré Conjecture) If $X = \text{surf}$ w/
 $\chi(X) = 2$, then $X = S^2$.

Cor: X surf, then $\chi(X) \leq 2$

Proof: Fix some poly cpx for X

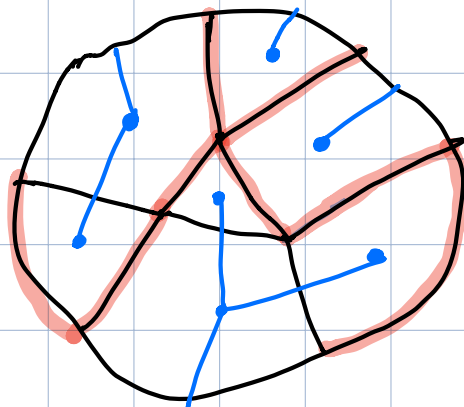
Note edges form a graph in X , say Γ

Pick spanning tree for Γ , say T .

Defn another graph Γ' on X as follows

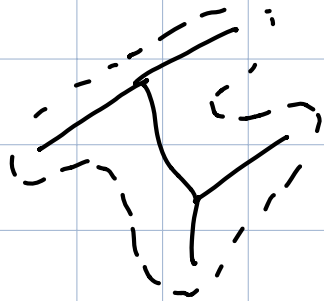
i) at each face we drop a vertex

ii) if two faces are conn via edge \Rightarrow conn vertices.



Claim Γ' is connected

$\Leftrightarrow T \subset X$ thickened up is a disk.



$\Rightarrow X$ -disk is connected but Γ' lives on
 X -disk $\Rightarrow \Gamma'$ is connected itself.

$$\begin{aligned}\chi(X) &= V - E + F \\ &= \underline{V(T)} - (\underline{E(T)} + \underline{E(\Gamma')}) + \underline{V(\Gamma')} \\ &= \chi(T) + \chi(\Gamma')\end{aligned}$$

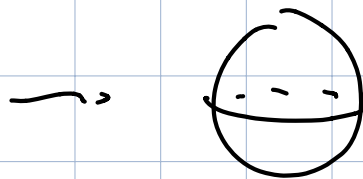
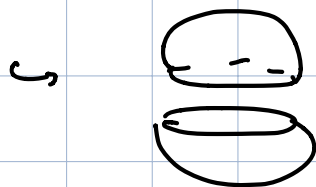
$$2 = \chi(X) = \chi(T) + \chi(\Gamma') = 1 + \chi(\Gamma')$$

$\Rightarrow \Gamma' = \text{tree}$.

Slowly thicken T and Γ' to fill up X

But thick (T) , thick (Γ') are both homeo
to disks.

\Rightarrow fill up realizes X as two disks glued along
their boundaries.



\square

Proof: Let T, Γ' be as before.

$$\Rightarrow \chi(X) = \chi(T) + \chi(\Gamma') = 1 + \chi(\Gamma').$$

If $\chi(X) = 2 \Rightarrow$ done $\Rightarrow X = S^2 \quad \ddot{\smile}$

$$\text{If } \chi(X) < 2 \Rightarrow 2 > 1 + \chi(\Gamma')$$

$$\Rightarrow 1 > \chi(\Gamma')$$

$\Rightarrow \Gamma'$ is not a tree

$\Rightarrow \Gamma'$ has a loop.

Call this loop γ .

X is orientable $\Rightarrow \gamma = 2$ -sided curve!

\Leftrightarrow using γ we realize X as a connect sum

$$X = T^2 \# X'$$

Note,

$$\begin{aligned} \chi(X) &= \chi(T^2) + \chi(X') - 2 \\ &= 0 + \chi(X') - 2 \end{aligned}$$

$$\Rightarrow \chi(X') > \chi(X).$$

$$X = T^2 \# X' = T^2 \# T^2 \# X''$$

Eventually X'', X''' will have $\chi \geq 2$

\Rightarrow eventually we get

$$\begin{aligned} X &= T^2 \# \dots \# T^2 \# S^2 \\ &= T^2 \# \dots \# T^2. \end{aligned}$$

