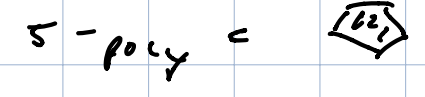
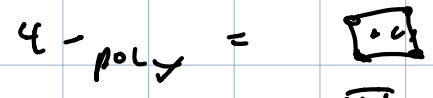
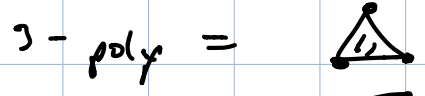
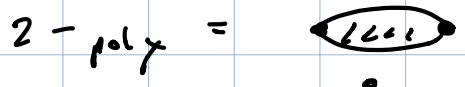


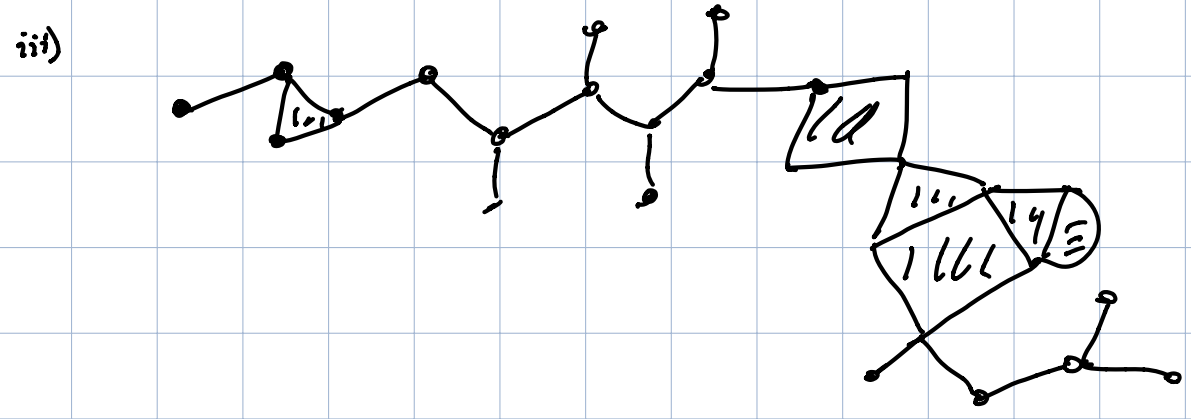
Lecture # 2

Defn: A polygonal cpx is a gluing of vertices, edges, and n -polygons where glue means glue edges to edges **along boundaries of polygons.**

\hookrightarrow n -poly = disk w/ n -sides



Ex: i)



Def: A graph is a polygonal cpx composed entirely of edges. A graph is a tree if each pair of vert may be joined via a unique path of edges

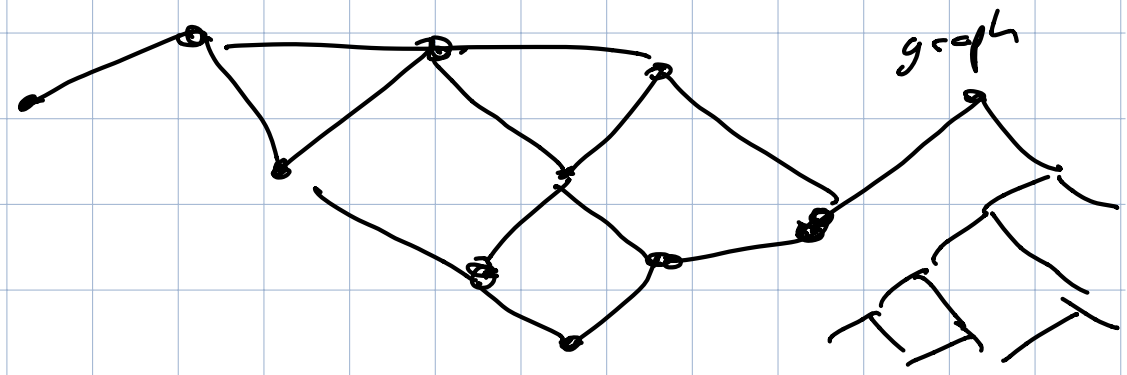
$E \subseteq$

i)

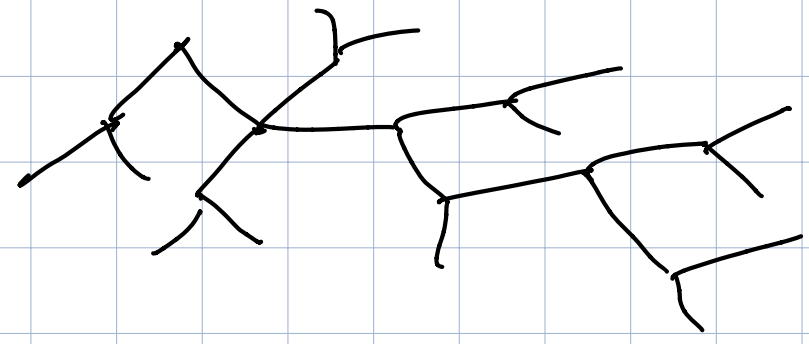


graph

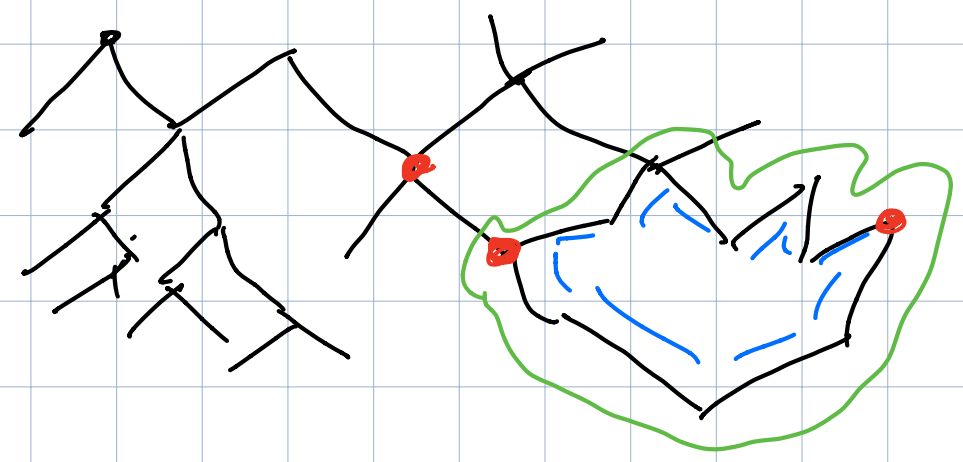
ii)



iii)



iv)



v)



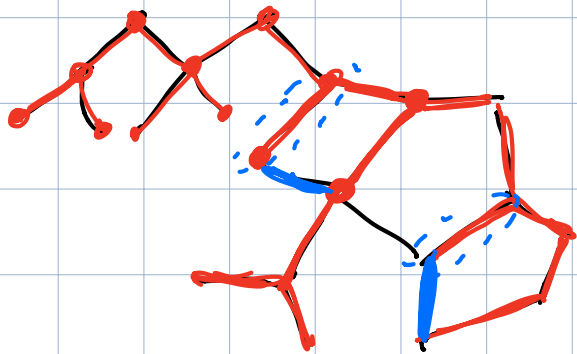
Remark:

Tree is a graph w/ no loops.

Def:

A spanning tree of a ^{conn.} graph is a subcollection of edges that form a conn tree and contain all the vertices of the graph.

Ex: i)



ii)



Remark: There can exist mult diff. spanning associated to any graph.

Prop: Every conn graph admits a spanning tree.

Proof: Build up the graph, Γ , in a seq of edge gluing. Denote this seq of cons. graphs by

$$\Gamma_0, \Gamma_1, \Gamma_2, \dots, \Gamma_n = \Gamma$$

\hookrightarrow

$$\Gamma_0 = \text{---}$$

$$\Gamma_1 = \text{---}$$

$$\Gamma_2 = \text{---}$$

$$\Gamma_3 = \text{---}$$

⋮

$$\Gamma_n$$

$$\Gamma_0 = \text{---}$$

$$\Gamma_1 = \text{---}$$

$$\Gamma_2 = \text{---}$$

$$\Gamma_3 = \text{---}$$

\hookrightarrow We can assume each Γ_i is connected as well.

For each Γ_i we construct a spanning tree T_i .

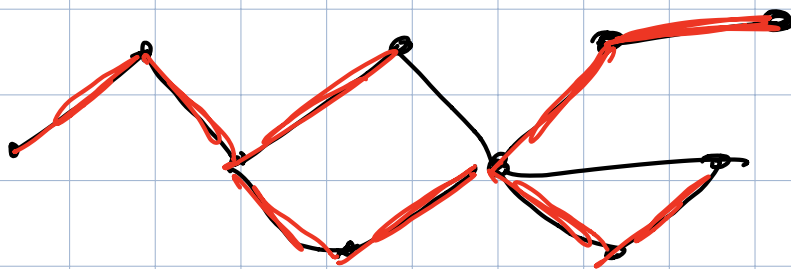
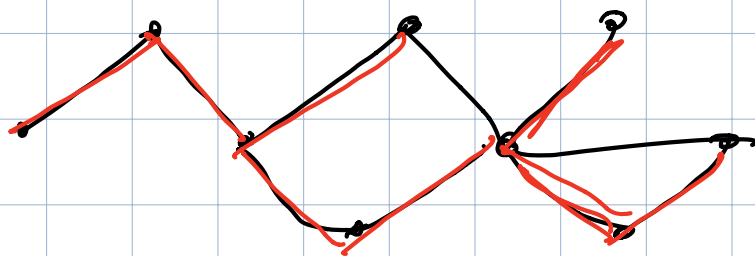
$$\begin{cases} T_0 = \Gamma_0 = \Gamma_0 \\ T_i = T_{i-1} & \text{if we don't add new vert to } \Gamma_i \\ T_i = \Gamma_i & \text{if we add a new vert. to } \Gamma_i \end{cases}$$

Spse that we know what the spanning tree is for Γ_i .

Say T_i .

Want to construct T_{i+1}

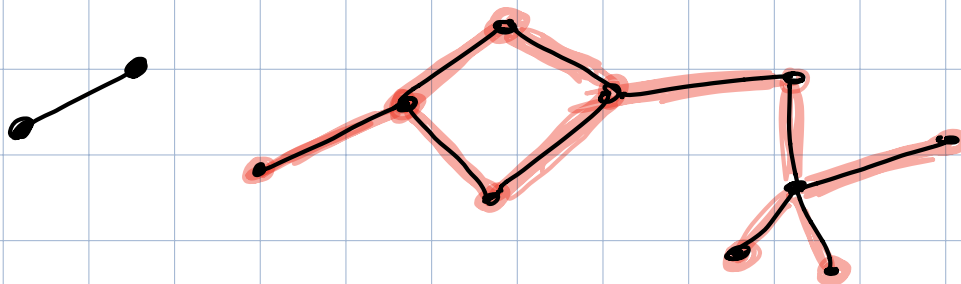
$$\begin{cases} T_{i+1} = T_i & \text{if we don't add new vert to } \Gamma_{i+1} \\ T_{i+1} = T_i \cup \text{newly added edges} & \text{if we add a new vert. to } \Gamma_{i+1}. \end{cases}$$



Repeat and repeat until $i = n$ and we are done!

Q.E.D. \square

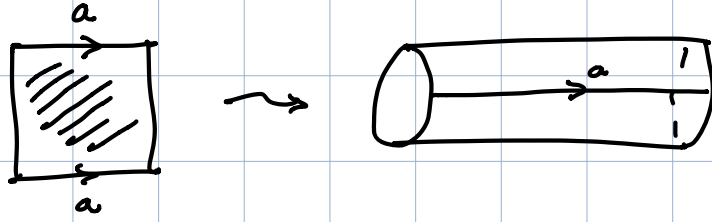
Link:



Defn 8 A planar dgm is a polygonal cpx obtained from gluing pairs of edges of a single $2n$ -polygon.

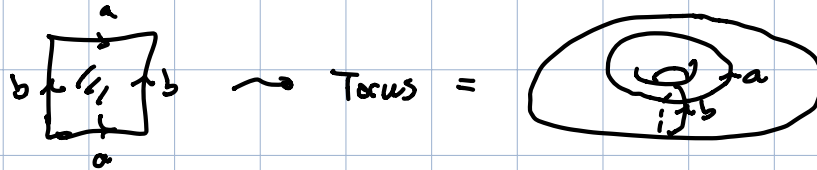
Ex:

i)



\rightarrow 2-sides

ii)



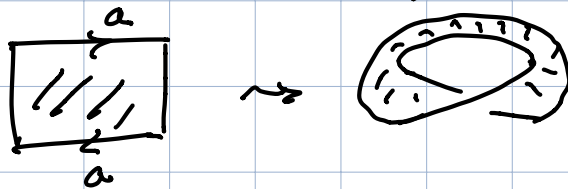
\equiv
 \uparrow
 set to be equal
 $\equiv T^2 = 2\text{-torus}$

iii)



($S^1 = \text{circle}$,
 $T^1 = 1\text{-torus}$)

iv)



\rightarrow 1-side

= Möbius band

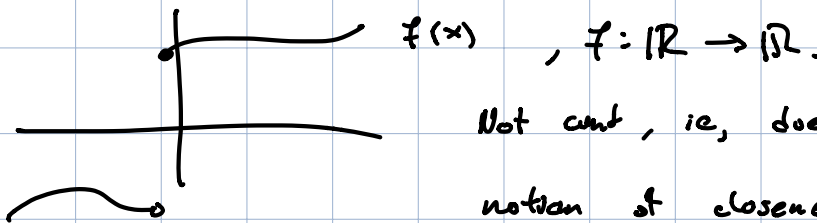
Defn:

A continuous map between top. spaces,

$f: X \rightarrow Y$, is an assignment of pts in X to points in Y st f takes pts infinitesimally close together to pts inf close together

\hookrightarrow cont. map preserves the notion of closeness.

\hookrightarrow



Not cont, ie, doesn't pres. notion of closeness.

\hookrightarrow

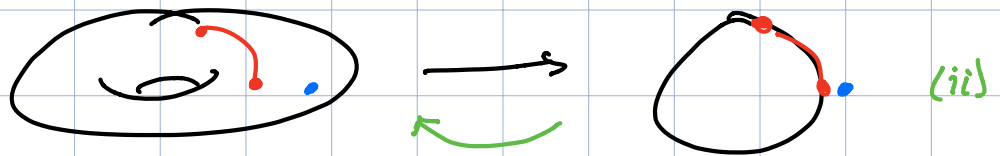


$f(\text{Northern hemi}) = 1$

$f(\text{southern hemi}) = -1$

$f: S^1 \rightarrow S^1$

$f(x) = \text{angle of } x \text{ about the rotating direction.}$



f^{-1} = Doesn't exist since it would not be a fun.

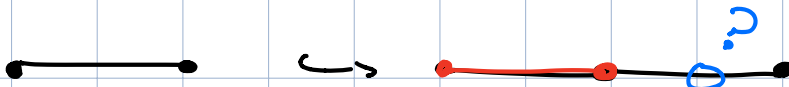
Yes it is cont.



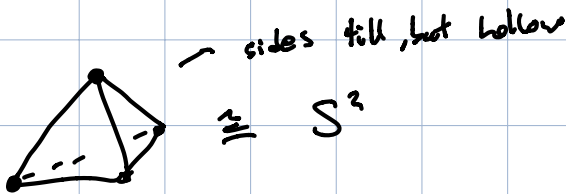
Defn: A homeomorphism is a continuous map $f: X \rightarrow Y$ that is a bijection.

\hookrightarrow if it is invertible (go backwards, $f^{-1}: Y \rightarrow X$ st $f \circ f^{-1} = \text{Identity}$ $f^{-1} \circ f = \text{Id.}$)

Remark:



Ex:



Defn: A surface is a top space that "locally" is homeo to an "open" disk in \mathbb{R}^2 \hookrightarrow about each point.

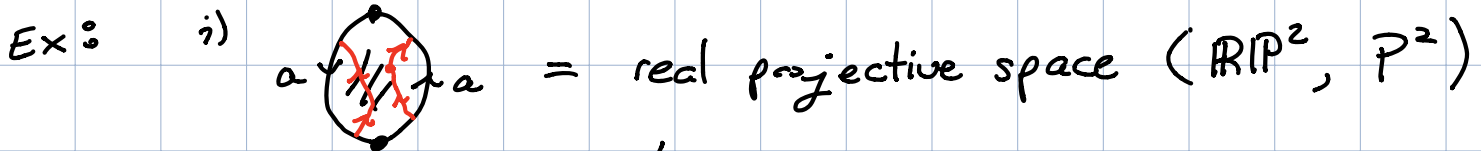
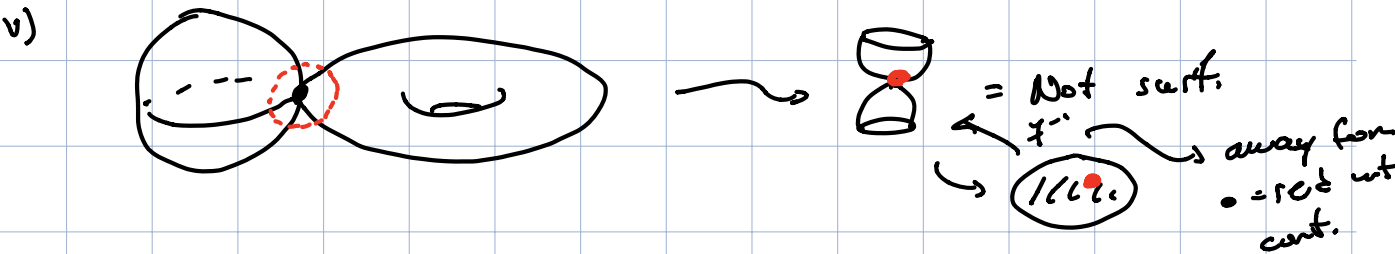
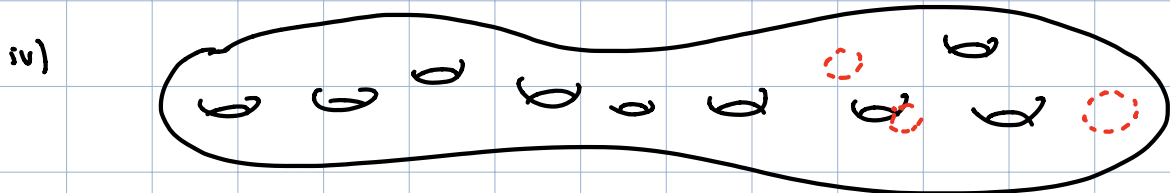
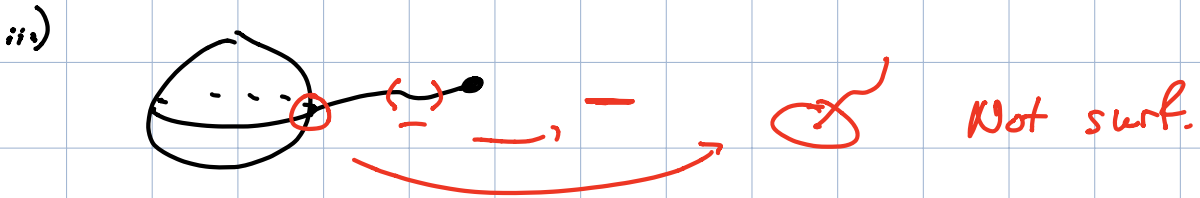
$\hookrightarrow \{(x, y) \mid x^2 + y^2 < 1\} =$

If you zoom in on the space it just looks like a

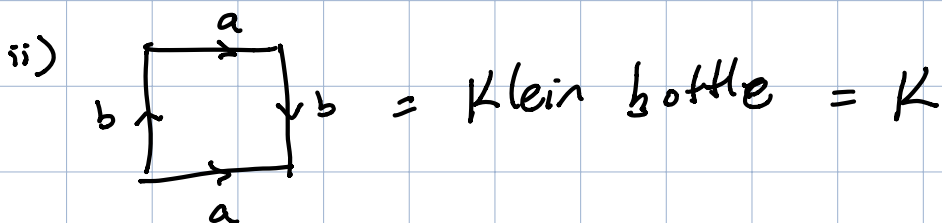
piece of paper.



Ex:



↳ can't be embedded in 3-dim space



Prop: A planar dgm w/ all edges glued (not like cycl) is a surface.

Proof:

