

**INTRODUCTION  
TO  
LINEAR  
ALGEBRA  
Fifth Edition**

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**MANUAL FOR INSTRUCTORS**

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**Wellesley - Cambridge Press**

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## Problem Set 11.1, page 516

- 1 Without exchange, pivots .001 and 1000; with exchange, 1 and  $-1$ . When the pivot is

larger than the entries below it, all  $|\ell_{ij}| = \frac{|\text{entry}|}{|\text{pivot}|} \leq 1$ .  $A = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & -1 \\ -1 & 1 & 1 \end{bmatrix}$ .

2 The exact inverse of  $\text{hilb}(3)$  is  $A^{-1} = \begin{bmatrix} 9 & -36 & 30 \\ -36 & 192 & -180 \\ 30 & -180 & 180 \end{bmatrix}$ .

3  $A \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 11/6 \\ 13/12 \\ 47/60 \end{bmatrix} = \begin{bmatrix} 1.833 \\ 1.083 \\ 0.783 \end{bmatrix}$  compares with  $A \begin{bmatrix} 0 \\ 6 \\ -3.6 \end{bmatrix} = \begin{bmatrix} 1.80 \\ 1.10 \\ 0.78 \end{bmatrix}$ .  $\|\Delta \mathbf{b}\| < .04$  but  $\|\Delta \mathbf{x}\| > 6$ .

The difference  $(1, 1, 1) - (0, 6, -3.6)$  is in a direction  $\Delta \mathbf{x}$  that has  $A\Delta \mathbf{x}$  near zero.

- 4 The largest  $\|\mathbf{x}\| = \|A^{-1}\mathbf{b}\|$  is  $\|A^{-1}\| = 1/\lambda_{\min}$  since  $A^T = A$ ; largest error  $10^{-16}/\lambda_{\min}$ .

- 5 Each row of  $U$  has at most  $w$  entries. Use  $w$  multiplications to substitute components of  $\mathbf{x}$  (already known from below) and divide by the pivot. Total for  $n$  rows  $< wn$ .

- 6 The triangular  $L^{-1}$ ,  $U^{-1}$ ,  $R^{-1}$  need  $\frac{1}{2}n^2$  multiplications.  $Q$  needs  $n^2$  to multiply the right side by  $Q^{-1} = Q^T$ . So  $QR\mathbf{x} = \mathbf{b}$  takes 1.5 times longer than  $LU\mathbf{x} = \mathbf{b}$ .

- 7  $UU^{-1} = I$ : Back substitution needs  $\frac{1}{2}j^2$  multiplications on column  $j$ , using the  $j$  by  $j$  upper left block. Then  $\frac{1}{2}(1^2 + 2^2 + \dots + n^2) \approx \frac{1}{2}(\frac{1}{3}n^3) =$  total to find  $U^{-1}$ .

8  $\begin{bmatrix} 1 & 0 \\ 2 & 2 \end{bmatrix} \rightarrow \begin{bmatrix} 2 & 2 \\ 1 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 2 & 2 \\ 0 & -1 \end{bmatrix} = U$  with  $P = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$  and  $L = \begin{bmatrix} 1 & 0 \\ .5 & 1 \end{bmatrix}$ ;

$A \rightarrow \begin{bmatrix} 2 & 2 & 0 \\ 1 & 0 & 1 \\ 0 & 2 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 2 & 2 & 0 \\ 0 & -1 & 1 \\ 0 & 2 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 2 & 2 & 0 \\ 0 & 2 & 0 \\ 0 & -1 & 1 \end{bmatrix} \rightarrow \begin{bmatrix} 2 & 2 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 1 \end{bmatrix} = U$  with

$P = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$  and  $L = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ .5 & -.5 & 1 \end{bmatrix}$ .

9  $A = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$  has cofactors  $C_{13} = C_{31} = C_{24} = C_{42} = 1$  and  $C_{14} = C_{41} = -1$ .  $A^{-1}$  is a full matrix!

10 With 16-digit floating point arithmetic the errors  $\|\mathbf{x} - \mathbf{x}_{\text{computed}}\|$  for  $\varepsilon = 10^{-3}, 10^{-6}, 10^{-9}, 10^{-12}, 10^{-15}$  are of order  $10^{-16}, 10^{-11}, 10^{-7}, 10^{-4}, 10^{-3}$ .

11 (a)  $\cos \theta = 1/\sqrt{10}$ ,  $\sin \theta = -3/\sqrt{10}$ ,  $R = \frac{1}{\sqrt{10}} \begin{bmatrix} 1 & 3 \\ -3 & 1 \end{bmatrix} \begin{bmatrix} 1 & -1 \\ 3 & 5 \end{bmatrix} = \frac{1}{\sqrt{10}} \begin{bmatrix} 10 & 14 \\ 0 & 8 \end{bmatrix}$ .

(b)  $A$  has eigenvalues 4 and 2. Put one of the unit eigenvectors in row 1 of  $Q$ : either

$$Q = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & 1 \end{bmatrix} \text{ and } QAQ^{-1} = \begin{bmatrix} 2 & -4 \\ 0 & 4 \end{bmatrix} \text{ or}$$

$$Q = \frac{1}{\sqrt{10}} \begin{bmatrix} 1 & -3 \\ 3 & 1 \end{bmatrix} \text{ and } QAQ^{-1} = \begin{bmatrix} 4 & -4 \\ 0 & 2 \end{bmatrix}.$$

12 When  $A$  is multiplied by a plane rotation  $Q_{ij}$ , this changes the  $2n$  (not  $n^2$ ) entries in rows  $i$  and  $j$ . Then multiplying on the right by  $(Q_{ij})^{-1} = (Q_{ij})^T$  changes the  $2n$  entries in columns  $i$  and  $j$ .

13  $Q_{ij}A$  uses  $4n$  multiplications (2 for each entry in rows  $i$  and  $j$ ). By factoring out  $\cos \theta$ , the entries 1 and  $\pm \tan \theta$  need only  $2n$  multiplications, which leads to  $\frac{2}{3}n^3$  for  $QR$ .

14 The  $(2, 1)$  entry of  $Q_{21}A$  is  $\frac{1}{3}(-\sin \theta + 2 \cos \theta)$ . This is zero if  $\sin \theta = 2 \cos \theta$  or  $\tan \theta = 2$ . Then the  $2, 1, \sqrt{5}$  right triangle has  $\sin \theta = 2/\sqrt{5}$  and  $\cos \theta = 1/\sqrt{5}$ .

Every 3 by 3 rotation with  $\det Q = +1$  is the product of 3 plane rotations.

15 This problem shows how elimination is more expensive (the nonzero multipliers in  $L$  and  $LL$  are counted by  $\mathbf{nnz}(L)$  and  $\mathbf{nnz}(LL)$ ) when we spoil the tridiagonal  $K$  by a random permutation.

If on the other hand we start with a poorly ordered matrix  $K$ , an improved ordering is found by the code **symamd** discussed in this section.

- 16** The “red-black ordering” puts rows and columns 1 to 10 in the odd-even order 1, 3, 5, 7, 9, 2, 4, 6, 8, 10. When  $K$  is the  $-1, 2, -1$  tridiagonal matrix, odd points are connected only to even points (and 2 stays on the diagonal, connecting every point to itself):

$$K = \begin{bmatrix} 2 & -1 & & & & & & & & \\ -1 & 2 & -1 & & & & & & & \\ & & \cdot & \cdot & \cdot & & & & & \\ & & & & & & & & & \\ & & & & -1 & 2 & & & & \\ & & & & & & & & & \end{bmatrix} \quad \text{and } PKP^T = \begin{bmatrix} 2I & D \\ D^T & 2I \end{bmatrix} \quad \text{with}$$

$$D = \begin{bmatrix} -1 & & & & & & & & & \\ -1 & -1 & & & & & & & & \\ 0 & -1 & -1 & & & & & & & \\ & & & -1 & -1 & & & & & \\ & & & & -1 & -1 & & & & \end{bmatrix} \begin{array}{l} 1 \text{ to } 2 \\ 3 \text{ to } 2, 4 \\ 5 \text{ to } 4, 6 \\ 7 \text{ to } 6, 8 \\ 9 \text{ to } 8, 10 \end{array}$$

- 17** Jeff Stuart’s **Shake a Stick** activity has long sticks representing the graphs of two linear equations in the  $x$ - $y$  plane. The matrix is nearly singular and Section 9.2 shows how to compute its condition number  $c = \|A\|\|A^{-1}\| = \sigma_{\max}/\sigma_{\min} \approx 80,000$ :

$$A = \begin{bmatrix} 1 & 1.0001 \\ 1 & 1.0000 \end{bmatrix} \quad \|A\| \approx 2 \quad A^{-1} = 10000 \begin{bmatrix} -1 & 1.0001 \\ 1 & -1 \end{bmatrix} \quad \begin{array}{l} \|A^{-1}\| \approx 20000 \\ c \approx 40000. \end{array}$$

### Problem Set 11.2, page 522

- 1**  $\|A\| = 2$ ,  $\|A^{-1}\| = 2$ ,  $c = 4$ ;  $\|A\| = 3$ ,  $\|A^{-1}\| = 1$ ,  $c = 3$ ;  $\|A\| = 2 + \sqrt{2} = \lambda_{\max}$  for positive definite  $A$ ,  $\|A^{-1}\| = 1/\lambda_{\min}$ ,  $\text{comd} = (2 + \sqrt{2})/(2 - \sqrt{2}) = 5.83$ .
- 2**  $\|A\| = 2$ ,  $c = 1$ ;  $\|A\| = \sqrt{2}$ ,  $c = \infty$  (singular matrix);  $A^T A = 2I$ ,  $\|A\| = \sqrt{2}$ ,  $c = 1$ .
- 3** For the first inequality replace  $\mathbf{x}$  by  $B\mathbf{x}$  in  $\|A\mathbf{x}\| \leq \|A\|\|\mathbf{x}\|$ ; the second inequality is just  $\|B\mathbf{x}\| \leq \|B\|\|\mathbf{x}\|$ . Then  $\|AB\| = \max(\|AB\mathbf{x}\|/\|\mathbf{x}\|) \leq \|A\|\|B\|$ .
- 4**  $1 = \|I\| = \|AA^{-1}\| \leq \|A\|\|A^{-1}\| = c(A)$ .

- 5** If  $\Lambda_{\max} = \Lambda_{\min} = 1$  then all  $\Lambda_i = 1$  and  $A = SIS^{-1} = I$ . The only matrices with  $\|A\| = \|A^{-1}\| = 1$  are *orthogonal matrices*.
- 6** All orthogonal matrices have norm 1, so  $\|A\| \leq \|Q\|\|R\| = \|R\|$  and in reverse  $\|R\| \leq \|Q^{-1}\|\|A\| = \|A\|$ . Then  $\|A\| = \|R\|$ . Inequality is usual in  $\|A\| < \|L\|\|U\|$  when  $A^T A \neq AA^T$ . Use **norm** on a random  $A$ .

**7** The triangle inequality gives  $\|A\mathbf{x} + B\mathbf{x}\| \leq \|A\mathbf{x}\| + \|B\mathbf{x}\|$ . Divide by  $\|\mathbf{x}\|$  and take the maximum over all nonzero vectors to find  $\|A + B\| \leq \|A\| + \|B\|$ .

**8** If  $A\mathbf{x} = \lambda\mathbf{x}$  then  $\|A\mathbf{x}\|/\|\mathbf{x}\| = |\lambda|$  for that particular vector  $\mathbf{x}$ . When we maximize the ratio  $\|A\mathbf{x}\|/\|\mathbf{x}\|$  over all vectors we get  $\|A\| \geq |\lambda|$ .

**9**  $A + B = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$  has  $\rho(A) = 0$  and  $\rho(B) = 0$  but  $\rho(A + B) = 1$ .

The triangle inequality  $\|A + B\| \leq \|A\| + \|B\|$  fails for  $\rho(A)$ .  $AB = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$  has  $\rho(AB) > \rho(A)\rho(B)$ ; thus  $\rho(A) = \max |\lambda(A)| =$  spectral radius is not a norm.

**10** (a) The condition number of  $A^{-1}$  is  $\|A^{-1}\|\|(A^{-1})^{-1}\|$  which is  $\|A^{-1}\|\|A\| = c(A)$ .

(b) Since  $A^T A$  and  $AA^T$  have the same nonzero eigenvalues,  $A^T$  has the same norm as  $A$ .

**11** Use the quadratic formula for  $\lambda_{\max}/\lambda_{\min}$ , which is  $c = \sigma_{\max}/\sigma_{\min}$  since this  $A = A^T$  is positive definite:

$$c(A) = \left(1.00005 + \sqrt{(1.00005)^2 - .0001}\right) / \left(1.00005 - \sqrt{\quad}\right) \approx 40,000.$$

**12**  $\det(2A)$  is not  $2 \det A$ ;  $\det(A + B)$  is not always less than  $\det A + \det B$ ; taking  $|\det A|$  does not help. The only reasonable property is  $\det AB = (\det A)(\det B)$ . The condition number should not change when  $A$  is multiplied by 10.

**13** The residual  $\mathbf{b} - A\mathbf{y} = (10^{-7}, 0)$  is much smaller than  $\mathbf{b} - A\mathbf{z} = (.0013, .0016)$ . But  $\mathbf{z}$  is much closer to the solution than  $\mathbf{y}$ .

**14**  $\det A = 10^{-6}$  so  $A^{-1} = 10^3 \begin{bmatrix} 659 & -563 \\ -913 & 780 \end{bmatrix}$ :  $\|A\| > 1$ ,  $\|A^{-1}\| > 10^6$ , then  $c > 10^6$ .

- 15**  $\mathbf{x} = (1, 1, 1, 1, 1)$  has  $\|\mathbf{x}\| = \sqrt{5}$ ,  $\|\mathbf{x}\|_1 = 5$ ,  $\|\mathbf{x}\|_\infty = 1$ .  $\mathbf{x} = (.1, .7, .3, .4, .5)$  has  $\|\mathbf{x}\| = 1$ ,  $\|\mathbf{x}\|_1 = 2$  (sum),  $\|\mathbf{x}\|_\infty = .7$  (largest).
- 16**  $x_1^2 + \cdots + x_n^2$  is not smaller than  $\max(x_i^2)$  and not larger than  $(|x_1| + \cdots + |x_n|)^2 = \|\mathbf{x}\|_1^2$ .  $x_1^2 + \cdots + x_n^2 \leq n \max(x_i^2)$  so  $\|\mathbf{x}\| \leq \sqrt{n}\|\mathbf{x}\|_\infty$ . Choose  $y_i = \text{sign } x_i = \pm 1$  to get  $\|\mathbf{x}\|_1 = \mathbf{x} \cdot \mathbf{y} \leq \|\mathbf{x}\|\|\mathbf{y}\| = \sqrt{n}\|\mathbf{x}\|$ . The vector  $\mathbf{x} = (1, \dots, 1)$  has  $\|\mathbf{x}\|_1 = \sqrt{n}\|\mathbf{x}\|$ .
- 17** For the  $\ell^\infty$  norm, the largest component of  $\mathbf{x}$  plus the largest component of  $\mathbf{y}$  is not less than  $\|\mathbf{x} + \mathbf{y}\|_\infty =$  largest component of  $\mathbf{x} + \mathbf{y}$ .
- For the  $\ell^1$  norm, each component has  $|x_i + y_i| \leq |x_i| + |y_i|$ . Sum on  $i = 1$  to  $n$ :  $\|\mathbf{x} + \mathbf{y}\|_1 \leq \|\mathbf{x}\|_1 + \|\mathbf{y}\|_1$ .
- 18**  $|x_1| + 2|x_2|$  is a norm but  $\min(|x_1|, |x_2|)$  is not a norm.  $\|\mathbf{x}\| + \|\mathbf{x}\|_\infty$  is a norm;  $\|A\mathbf{x}\|$  is a norm provided  $A$  is invertible (otherwise a nonzero vector has norm zero; for rectangular  $A$  we require independent columns to avoid  $\|A\mathbf{x}\| = 0$ ).
- 19**  $\mathbf{x}^T \mathbf{y} = x_1 y_1 + x_2 y_2 + \cdots \leq (\max |y_i|)(|x_1| + |x_2| + \cdots) = \|\mathbf{x}\|_1 \|\mathbf{y}\|_\infty$ .
- 20** With  $\lambda_j = 2 - 2 \cos(j\pi/n+1)$ , the largest eigenvalue is  $\lambda_n \approx 2 + 2 = 4$ . The smallest is  $\lambda_1 = 2 - 2 \cos(\pi/n+1) \approx \left(\frac{\pi}{n+1}\right)^2$ , using  $2 \cos \theta \approx 2 - \theta^2$ . So the condition number is  $c = \lambda_{\max}/\lambda_{\min} \approx (4/\pi^2) n^2$ , growing with  $n$ .

### Problem Set 11.3, page 531

- 1** The iteration  $\mathbf{x}_{k+1} = (I - A)\mathbf{x}_k + \mathbf{b}$  has  $S = I$  and  $T = I - A$  and  $S^{-1}T = I - A$ .
- 2** If  $A\mathbf{x} = \lambda\mathbf{x}$  then  $(I - A)\mathbf{x} = (1 - \lambda)\mathbf{x}$ . Real eigenvalues of  $B = I - A$  have  $|1 - \lambda| < 1$  provided  $\lambda$  is between 0 and 2.
- 3** This matrix  $A$  has  $I - A = \begin{bmatrix} -1 & 1 \\ 1 & -1 \end{bmatrix}$  which has  $|\lambda| = 2$ . The iteration diverges.
- 4** Always  $\|AB\| \leq \|A\|\|B\|$ . Choose  $A = B$  to find  $\|B^2\| \leq \|B\|^2$ . Then choose  $A = B^2$  to find  $\|B^3\| \leq \|B^2\|\|B\| \leq \|B\|^3$ . Continue (or use induction) to find  $\|B^k\| \leq \|B\|^k$ . Since  $\|B\| \geq \max |\lambda(B)|$  it is no surprise that  $\|B\| < 1$  gives convergence.

**5**  $A\mathbf{x} = \mathbf{0}$  gives  $(S - T)\mathbf{x} = \mathbf{0}$ . Then  $S\mathbf{x} = T\mathbf{x}$  and  $S^{-1}T\mathbf{x} = \mathbf{x}$ . Then  $\lambda = 1$  means that the errors do not approach zero. We can't expect convergence when  $A$  is singular and  $A\mathbf{x} = \mathbf{b}$  is unsolvable!

**6** Jacobi has  $S^{-1}T = \frac{1}{3} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$  with  $|\lambda|_{\max} = \frac{1}{3}$ . Small problem, fast convergence.

**7** Gauss-Seidel has  $S^{-1}T = \begin{bmatrix} 0 & \frac{1}{3} \\ 0 & \frac{1}{9} \end{bmatrix}$  with  $|\lambda|_{\max} = \frac{1}{9}$  which is  $(|\lambda|_{\max}$  for Jacobi)<sup>2</sup>.

**8** Jacobi has  $S^{-1}T = \begin{bmatrix} a & \\ & d \end{bmatrix}^{-1} \begin{bmatrix} 0 & -b \\ -c & 0 \end{bmatrix} = \begin{bmatrix} 0 & -b/a \\ -c/d & 0 \end{bmatrix}$  with  $|\lambda| = |bc/ad|^{1/2}$ .

Gauss-Seidel has  $S^{-1}T = \begin{bmatrix} a & 0 \\ c & d \end{bmatrix}^{-1} \begin{bmatrix} 0 & -b \\ 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & -b/a \\ 0 & -bc/ad \end{bmatrix}$  with  $|\lambda| = |bc/ad|$ .

So Gauss-Seidel is twice as fast to converge if  $|\lambda| < 1$  (or to explode if  $|bc| > |ad|$ ).

**9** Gauss-Seidel will converge for the  $-1, 2, -1$  matrix.  $|\lambda|_{\max} = \cos^2\left(\frac{\pi}{n+1}\right)$  is given on page 527, together with the improvement from successive overrelaxation.

**10** If the iteration gives all  $x_i^{\text{new}} = x_i^{\text{old}}$  then the quantity in parentheses is zero, which means  $A\mathbf{x} = \mathbf{b}$ . For Jacobi change  $\mathbf{x}^{\text{new}}$  on the right side to  $\mathbf{x}^{\text{old}}$ .

**11**  $\mathbf{u}_k/\lambda_1^k = c_1\mathbf{x}_1 + c_2\mathbf{x}_2(\lambda_2/\lambda_1)^k + \cdots + c_n\mathbf{x}_n(\lambda_n/\lambda_1)^k \rightarrow c_1\mathbf{x}_1$  if all ratios  $|\lambda_i/\lambda_1| <$

1. The largest ratio controls the rate of convergence (when  $k$  is large).  $A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$

has  $|\lambda_2| = |\lambda_1|$  and no convergence.

**12** The eigenvectors of  $A$  and also  $A^{-1}$  are  $\mathbf{x}_1 = (.75, .25)$  and  $\mathbf{x}_2 = (1, -1)$ . The inverse power method converges to a multiple of  $\mathbf{x}_2$ , since  $|1/\lambda_2| > |1/\lambda_1|$ .

**13** In the  $j$ th component of  $A\mathbf{x}_1$ ,  $\lambda_1 \sin \frac{j\pi}{n+1} = 2 \sin \frac{j\pi}{n+1} - \sin \frac{(j-1)\pi}{n+1} - \sin \frac{(j+1)\pi}{n+1}$ .

The last two terms combine into  $-2 \sin \frac{j\pi}{n+1} \cos \frac{\pi}{n+1}$ . Then  $\lambda_1 = 2 - 2 \cos \frac{\pi}{n+1}$ .

**14**  $A = \begin{bmatrix} 2 & -1 \\ -1 & 2 \end{bmatrix}$  produces  $\mathbf{u}_0 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$ ,  $\mathbf{u}_1 = \begin{bmatrix} 2 \\ -1 \end{bmatrix}$ ,  $\mathbf{u}_2 = \begin{bmatrix} 5 \\ -4 \end{bmatrix}$ ,  $\mathbf{u}_3 = \begin{bmatrix} 14 \\ -13 \end{bmatrix}$ .

This is converging to the eigenvector direction  $\begin{bmatrix} 1 \\ -1 \end{bmatrix}$  with largest eigenvalue  $\lambda = 3$ .

Divide  $\mathbf{u}_k$  by  $\|\mathbf{u}_k\|$  to keep unit vectors.

**15**  $A^{-1} = \frac{1}{3} \begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix}$  gives  $\mathbf{u}_1 = \frac{1}{3} \begin{bmatrix} 2 \\ 1 \end{bmatrix}$ ,  $\mathbf{u}_2 = \frac{1}{9} \begin{bmatrix} 5 \\ 4 \end{bmatrix}$ ,  $\mathbf{u}_3 = \frac{1}{27} \begin{bmatrix} 14 \\ 13 \end{bmatrix} \rightarrow \mathbf{u}_\infty = \begin{bmatrix} 1/2 \\ 1/2 \end{bmatrix}$ .

**16**  $R = Q^T A = \begin{bmatrix} 1 & \cos \theta \sin \theta \\ 0 & -\sin^2 \theta \end{bmatrix}$  and  $A_1 = RQ = \begin{bmatrix} \cos \theta(1 + \sin^2 \theta) & -\sin^3 \theta \\ -\sin^3 \theta & -\cos \theta \sin^2 \theta \end{bmatrix}$ .

**17** If  $A$  is orthogonal then  $Q = A$  and  $R = I$ . Therefore  $A_1 = RQ = A$  again, and the “QR method” doesn’t move from  $A$ . But shift  $A$  slightly and the method goes quickly to  $\Lambda$ .

**18** If  $A - cI = QR$  then  $A_1 = RQ + cI = Q^{-1}(QR + cI)Q = Q^{-1}AQ$ . No change in eigenvalues from the shift and shift back, because  $A_1$  is similar to  $A$ .

**19** Multiply  $A\mathbf{q}_j = b_{j-1}\mathbf{q}_{j-1} + a_j\mathbf{q}_j + b_j\mathbf{q}_{j+1}$  by  $\mathbf{q}_j^T$  to find  $\mathbf{q}_j^T A\mathbf{q}_j = a_j$  (because the  $\mathbf{q}$ ’s are orthonormal). The matrix form (multiplying by columns) is  $AQ = QT$  where  $T$  is *tridiagonal*. The entries down the diagonals of  $T$  are the  $a$ ’s and  $b$ ’s.

**20** Theoretically the  $\mathbf{q}$ ’s are orthonormal. In reality this important algorithm is not very stable. We must stop every few steps to reorthogonalize—or find another more stable way to orthogonalize the sequence  $\mathbf{q}, A\mathbf{q}, A^2\mathbf{q}, \dots$

**21** If  $A$  is symmetric then  $A_1 = Q^{-1}AQ = Q^T AQ$  is also symmetric.  $A_1 = RQ = R(QR)R^{-1} = RAR^{-1}$  has  $R$  and  $R^{-1}$  upper triangular, so  $A_1$  cannot have nonzeros on a lower diagonal than  $A$ . If  $A$  is tridiagonal and symmetric then (by using symmetry for the upper part of  $A_1$ ) the matrix  $A_1 = RAR^{-1}$  is also tridiagonal.

**22** From the last line of code,  $\mathbf{q}_2$  is in the direction of  $\mathbf{v} = A\mathbf{q}_1 - h_{11}\mathbf{q}_1 = A\mathbf{q}_1 - (\mathbf{q}_1^T A\mathbf{q}_1)\mathbf{q}_1$ . The dot product with  $\mathbf{q}_1$  is zero. This is Gram-Schmidt with  $A\mathbf{q}_1$  as the second input vector; we subtract from  $A\mathbf{q}_1$  its projection onto the first vector  $\mathbf{q}_1$ .



*Note* The three lines after the short “pseudocodes” describe two key properties of conjugate gradients—the residuals  $\mathbf{r}_k = \mathbf{b} - A\mathbf{x}_k$  are orthogonal and the search directions are  $A$ -orthogonal ( $\mathbf{d}_i^T A \mathbf{d}_k = 0$ ). Then each new approximation  $\mathbf{x}_{k+1}$  is the **closest vector to  $\mathbf{x}$**  among all combinations of  $\mathbf{b}, A\mathbf{b}, \dots, A^k \mathbf{b}$ . Ordinary iteration  $S\mathbf{x}_{k+1} = T\mathbf{x}_k + \mathbf{b}$  does *not* find this best possible combination  $\mathbf{x}_{k+1}$ .

**23** The solution is straightforward and important. Since  $H = Q^{-1}AQ = Q^T A Q$  is *symmetric* if  $A = A^T$ , and since  $H$  has only one lower diagonal by construction, then  $H$  has only *one upper diagonal*:  $H$  is tridiagonal and all the recursions in Arnoldi’s method have only 3 terms.

**24**  $H = Q^{-1}AQ$  is *similar* to  $A$ , so  $H$  has the same eigenvalues as  $A$  (at the end of Arnoldi). When Arnoldi is stopped sooner because the matrix size is large, the eigenvalues of  $H_k$  (called *Ritz values*) are close to eigenvalues of  $A$ . This is an important way to compute approximations to  $\lambda$  for large matrices.

**25** In principle the conjugate gradient method converges in 100 (or 99) steps to the exact solution  $\mathbf{x}$ . But it is slower than elimination and its all-important property is to give good approximations to  $\mathbf{x}$  much sooner. (Stopping elimination part way leaves you nothing.) The problem asks how close  $\mathbf{x}_{10}$  and  $\mathbf{x}_{20}$  are to  $\mathbf{x}_{100}$ , which equals  $\mathbf{x}$  except for roundoff errors.

**26**  $A = \begin{bmatrix} 1 & 1 \\ 0 & 1.1 \end{bmatrix}$  has  $A^n = \begin{bmatrix} 1 & q \\ 0 & (1.1)^n \end{bmatrix}$  with  $q = 1 + 1.1 + \dots + (1.1)^{n-1} =$

$(1.1^n - 1)/(1.1 - 1) \approx 10 (1.1)^n$ . So the growing part of  $A^n$  is  $(1.1)^n \begin{bmatrix} 0 & 10 \\ 0 & 1 \end{bmatrix}$

with  $\|A^n\| \approx \sqrt{101}$  times  $1.1^n$  for larger  $n$ .