

1 **LU AND CR ELIMINATION**

2 GILBERT STRANG AND CLEVE MOLER

3 **Abstract.** The reduced row echelon form $\mathbf{rref}(A)$ has traditionally been used for classroom
4 examples: small matrices A with integer entries and low rank r . This paper creates a column-row
5 rank-revealing factorization $A = CR$, with the first r independent columns of A in C and the r
6 nonzero rows of $\mathbf{rref}(A)$ in R . We want to reimagine the start of a linear algebra course, by helping
7 students to see the independent columns of A and the rank and the column space.

8 If B contains the first r independent rows of A , then those rows of $A = CR$ produce $B = WR$.
9 The r by r matrix W has full rank r , where B meets C . Then the triple factorization $A = CW^{-1}B$
10 treats columns and rows of A (C and B) in the same way.

11 **Key words.** elimination, factorization, row echelon form, matrix, rank

12 **AMS subject classifications.** 15A21, 15A23, 65F05, 65F55

13 **1. Introduction.** Matrix factorizations like $A = LU$ and $A = U\Sigma V^T$ have be-
14 come the organizing principles of linear algebra. This expository paper develops a
15 column-row factorization $A = CR = (m \times r)(r \times n)$ for any matrix of rank r . The
16 matrix C contains the *first r independent columns* of A : a basis for the column
17 space. The matrix R contains the *nonzero rows* of the reduced row echelon form
18 $\mathbf{rref}(A)$. We will put R in the form $R = [I \ F] P$, with an r by r identity ma-
19 trix that multiplies the r columns of C . Then $A = CR = [C \ CF] P$ expresses
20 the $n - r$ remaining (dependent) columns of A as combinations CF of the r in-
21 dependent columns in C . When those independent columns don't all come first in A ,
22 P permutes those columns of I and F into their correct positions.

23 The example in Section 3 shows how invertible row operations find the first r
24 independent columns of A . For a large matrix this row reduction is expensive and
25 numerically perilous. But Section 6 will explain the value of an approximate CR or
26 $CW^{-1}B$ factorization of A . This is achievable by randomized linear algebra.

27 The key point is: $A = CR$ is an “interpolative decomposition” that includes r
28 actual columns of A in C . A more symmetric two-sided factorization $A = CW^{-1}B$
29 also includes r actual rows of A in B . The r by r matrix W lies at the “intersection”
30 inside A of the columns of C with the rows of B . The mixing matrix W^{-1} removes
31 that repetition to produce $W^{-1}B = R$. If $P = I$ then we are seeing block elimination
32 with W as the block pivot:

33
$$A = \begin{bmatrix} W & H \\ J & K \end{bmatrix} = \begin{bmatrix} W \\ J \end{bmatrix} W^{-1} \begin{bmatrix} W & H \end{bmatrix} = CW^{-1}B$$

34 That matrix W is invertible, where a row basis B meets the column basis C . For
35 large matrices, a low rank version $A \approx CUB$ can give a high quality approximation.

36 **2. LU Elimination.** This is familiar to all numerical analysts. It applies best
 37 to an invertible n by n matrix A . A typical step subtracts a multiple ℓ_{ij} of row j from
 38 row i , to produce a zero in the ij position for each $i > j$. A column at a time, going
 39 left to right, all entries below the main diagonal become zero. LU elimination arrives
 40 at an upper triangular matrix U , with the n nonzero pivots on its main diagonal.

41 We could express that result as a matrix multiplication $EA = U$. The lower
 42 triangular matrix E is the product of all the single-step elimination matrices E_{ij} . A
 43 more successful idea—which reverses the order of the steps as it inverts them—is to
 44 consider the matrix $L = E^{-1}$ that brings U back to A . Then $EA = U$ becomes
 45 $A = LU$.

46 In this order, the lower triangular L contains all the multipliers ℓ_{ij} *exactly in their*
 47 *proper positions*. The pattern is only upset if any row exchanges become necessary
 48 to avoid a zero pivot or to obtain a larger pivot (and smaller multipliers). If all
 49 row exchanges are combined into a permutation matrix, elimination factors that row-
 50 permuted version of A into LU .

51 **3. CR Elimination.** Start now with an m by n matrix A of rank r . Elimination
 52 will again proceed left to right, a column at a time. The new goal is to produce an r
 53 by r identity matrix I . So each pivot row in turn is divided by its first nonzero entry,
 54 to produce the desired 1 in I . Then multiples of that pivot row are subtracted from
 55 the rows *above and below*, to achieve the zeros in that column of I .

56 If at any point a row is entirely zero, it moves to the bottom of the matrix. If a
 57 column is entirely zero except in the rows already occupied by earlier pivots, then no
 58 pivot is available in that column—and we move to the next column. The final result
 59 of this elimination is the m by n **reduced row echelon form** of A . We denote that
 60 form by R_0 :

$$61 \quad (3.1) \quad R_0 = \text{rref}(A) = \begin{bmatrix} I & F \\ \mathbf{0} & \mathbf{0} \end{bmatrix} P \quad \begin{array}{ll} r & \text{rows} \\ m-r & \text{rows} \end{array}$$

$$62 \quad \begin{array}{ll} r & n-r \end{array} \text{ columns}$$

63 The n by n permutation matrix P puts the columns of $I_{r \times r}$ into their correct positions,
 64 matching the positions of the first r independent columns of the original matrix A .
 65

$$66 \quad (3.2) \quad A = \begin{bmatrix} \mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{4} \\ \mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{5} \\ \mathbf{2} & \mathbf{4} & \mathbf{6} & \mathbf{9} \end{bmatrix} \rightarrow \begin{bmatrix} \mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{4} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{1} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \end{bmatrix} \rightarrow \begin{bmatrix} \mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{1} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \end{bmatrix} = R_0.$$

67 C contains columns 1 and 4 of A and R contains rows 1 and 2 of R_0 . Then CR con-
 68 tains the four columns of A : column 1, 2(column 1), 3(column 1), column 4. P^T puts I
 69 first:

$$70 \quad (3.3)$$

$$71 \quad R_0 P^T = \begin{bmatrix} \mathbf{1} & \mathbf{2} & \mathbf{3} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{1} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{1} & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{1} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{1} \\ \mathbf{0} & \mathbf{1} & \mathbf{0} & \mathbf{0} \end{bmatrix} = \begin{bmatrix} \mathbf{1} & \mathbf{0} & \mathbf{2} & \mathbf{3} \\ \mathbf{0} & \mathbf{1} & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \mathbf{0} \end{bmatrix} = \begin{bmatrix} I & F \\ \mathbf{0} & \mathbf{0} \end{bmatrix}.$$

72 All permutations have $P^T P = I$. So multiplying equation (3.3) by P and removing
 73 row 3 produces $R = \begin{bmatrix} I & F \end{bmatrix} P$.

74 Usually the description of elimination stops at R_0 . There is *no connection* to a
 75 matrix factorization $A = CR$. And the matrix F in the columns without pivots is
 76 given *no interpretation*. This misses an opportunity to understand more fully the
 77 **rref** algorithm and the structure of a general matrix A . We believe that $A = CR$ is
 78 a valuable idea in teaching linear algebra [11]: a good way to start.

79 **4. The Factorization $A = CR$: m by r times r by n .** The matrix C is
 80 easy to describe. It contains the *first r independent columns* of A . The positions
 81 of those independent columns are revealed by the identity matrix I in **rref**(A) and
 82 by P (the permutation). All other columns of A (with rank r) are combinations
 83 of these r columns. Those combinations come from the submatrix F . **The matrix**
 84 **R is the reduced row echelon form $R_0 = \mathbf{rref}(A)$ without its zero rows:**

$$85 \quad (4.1) \quad R = [I \quad F] P \quad \text{and} \quad A = CR = [C \quad CF] P$$

86 Again, the matrix F tells how to construct the $n - r$ dependent columns of A from
 87 the r independent columns in C . This interpretation is often missing from explana-
 88 tions of **rref**—it comes naturally when the process is expressed as a factorization. C
 89 *gives the independent columns of A and CF gives the dependent columns.* P orders
 90 those n columns correctly in A .

91 **Example 2** A is a 3 by 3 matrix of rank 2. Column 3 is $-$ column 1 + 2(column 2),
 92 so F contains -1 and 2. This example has $P = I$, and the zero row of R_0 is gone:

$$93 \quad A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix} = \begin{bmatrix} 1 & 2 \\ 4 & 5 \\ 7 & 8 \end{bmatrix} \begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & 2 \end{bmatrix} = CR = C [I \quad F]$$

94 The first two columns of A are a basis for its column space. Those columns are in C .
 95 **The two rows of R are a basis for the row space of A .** Those rows are
 96 independent (they contain I). They span the row space because $A = CR$ expresses
 97 every row of A as a combination of the rows of R .

98 Piziak and Odell [9] include $A = CR$ as the first of their “full rank factoriza-
 99 tions”: independent columns times independent rows. Then Gram-Schmidt ($A =$
 100 QR) achieves *orthogonal columns*, and the SVD also achieves *orthogonal rows*. Those
 101 great factorizations offer a numerical stability that elimination can never match.

102

103 Here are the steps to establish $A = CR$. We know that an invertible elimination
 104 matrix E (a product of simple steps) gives $EA = R_0 = \mathbf{rref}(A)$. Then $A = E^{-1}R_0 =$
 105 $(m \times m)(m \times n)$. Drop the $m - r$ zero rows of R_0 and the last $m - r$ columns of E^{-1} . This
 106 leaves $A =$
 107 $C [I \quad F] P$, where the identity matrix in R allows us to identify C in the columns of
 108 E^{-1} .

109

110 The factorization $A = CR$ reveals the first great theorem of linear algebra.

111 **The column rank r equals the row rank.**

112 **Proof** The r rows of R are independent (*from its submatrix I*). And all rows of A are
 113 combinations of the rows of R (*because $A = CR$*). Note that the rows of R belong to
 114 the

115 row space of A because $R = (C^T C)^{-1} C^T A$. So the row rank is r .

117 The $A = CR$ factorization reveals all these essential facts and more. The equation
 118 $A\mathbf{x} = \mathbf{0}$ becomes easy to solve. *Each dependent column of A is a combination of the r*
 119 *independent columns in C* . That gives $n - r$ “special solutions” to $A\mathbf{x} = \mathbf{0}$: a basis for
 120 the nullspace of A .

121 Put those solutions into the nullspace matrix N . Recalling that $PP^T = I$, here
 122 is $AN = \mathbf{zero\ matrix}$ (m by $n - r$):

123 (4.2) $A = [C \quad CF]P$ times $N = P^T \begin{bmatrix} -F \\ I_{n-r} \end{bmatrix}$ equals $-CF + CF = \mathbf{0}$

124 In Example 2 with $P = I$, N has one column $\mathbf{x} = [-F \quad 1]^T = [1 \quad -2 \quad 1]^T$.

125 The columns of N are “special” because of the submatrix I_{n-r} : One dependent
 126 column of A is a combination (given by F) of the independent columns in C .

128 This analysis of the $A = CR$ factorization was new to us. In a linear algebra
 129 course,

130 it allows early examples of the column space and the crucial ideas of *linear indepen-*
 131 *dence*

132 and *basis*. Those ideas appear first in examples like the matrix A above (before
 133 the student has a system for computing C). That system comes later—it is the elimi-
 134 nation that produces $R_0 = \mathbf{rref}(A)$ and then R , with their submatrix I that identifies
 135 the columns that go into C .

136 The goal is to start a linear algebra class with simple examples of independence
 137 and rank and matrix multiplication. Integer matrices A reveal the *idea* of indepen-
 138 dence before the definition. And the matrix $\mathbf{rref}(A)$ now has a meaning! The identity
 139 matrix locates the first r independent columns (a basis for the column space of A).
 140 And F (r by $n - r$) gives the combinations of those r columns in C that produce the
 141 $n - r$ dependent columns CF of A .

142 **5. The Magic Factorization $A = CW^{-1}B$.** The factorization $A = CR$ has
 143 useful properties, but symmetric treatment of columns and rows is not one of them.
 144 The rows of R were constructed by elimination. They identify the independent col-
 145 umns in C . There is a closely related factorization that takes r *independent rows* (*as*
 146 *well as r independent columns*) *directly from A* . Now an r by r submatrix W appears
 147 both in the column matrix C and the row matrix B : W is the “intersection” of r
 148 rows with r columns. So a factor W^{-1} must go in between C and B :

149 $A = CR = CW^{-1}B$ as in $A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix} = \begin{bmatrix} 1 & 2 \\ 4 & 5 \\ 7 & 8 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 4 & 5 \end{bmatrix}^{-1} \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix}$

150 The first r independent columns and independent rows of A meet in W . A key
 151 point is that W is invertible. By interposing W^{-1} between C and B , the factorization
 152 $A = CW^{-1}B$ succeeds with R replaced by $W^{-1}B$.

153 (5.1) Notice that $\mathbf{WR} = \begin{bmatrix} 1 & 2 \\ 4 & 5 \end{bmatrix} \begin{bmatrix} 1 & 0 & -1 \\ 0 & 1 & 2 \end{bmatrix} = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} = \mathbf{B}$.

154
 155 THEOREM 5.1. Suppose C contains the first r independent columns of A , with
 156 $r = \text{rank}(A)$. Suppose R contains the r nonzero rows of $\mathbf{rref}(A)$. Then $\mathbf{A} = \mathbf{CR}$.
 157 Now suppose the matrix B contains the first r independent rows of A . Then W
 158 is the r by r matrix where C meets B . If we look only at those r rows of $A = CR$, we
 159 see $\mathbf{B} = \mathbf{WR}$. Since B and R both have rank r , the square matrix W must have rank
 160 r and it is invertible.

162 In case A itself is square and invertible, we have $A = C = B = W$. Then the
 163 factorization $A = CW^{-1}B$ reduces to $A = WW^{-1}W$. For a rectangular matrix with
 164 independent columns and rows coming first in W , we see columns times W^{-1} times
 165 rows:

166 (5.2) $A = \begin{bmatrix} W & H \\ J & K \end{bmatrix} = \begin{bmatrix} W \\ J \end{bmatrix} W^{-1} \begin{bmatrix} W & H \end{bmatrix} = CW^{-1}B$

167 All this is the work of a “coordinated” person. If we ask an algebraist (as we did),
 168 the factors B and W^{-1} and C become coordinate-free linear maps that reproduce A :

169 (5.3) $A : \mathbf{R}^n \rightarrow \mathbf{R}^n / \text{kernel}(A) \rightarrow \text{image}(A) \rightarrow \mathbf{R}^m$

170 The map in the middle is an isomorphism between r -dimensional spaces. But the
 171 natural bases are different for those two spaces. The columns of C give the right basis
 172 for $\text{image}(A) =$ “column space”. The rows of B are natural for the dual space (row
 173 space) $\mathbf{R}^n / \text{kernel}(A)$. So there is an invertible r by r matrix—a change of basis—
 174 when we introduce coordinates. *That step is represented by W^{-1} .* The invertibility
 175 of W follows from the algebra.

176 When we verify equation (5.3) in matrix language (with coordinates), the in-
 177 vertibility of the column-row intersection W will be explicitly proved. Here is that
 178 additional proof.

| | | |
|--|--------------------------------------|--|
| <p>181 $\mathbf{A} =$ any matrix of rank r $m \times n$</p> <p>182 $\mathbf{C} =$ r independent columns of A $m \times r$</p> <p>183 $\mathbf{B} =$ r independent rows of A $r \times n$</p> <p>184 $\mathbf{W} =$ intersection of C and B $r \times r$</p> | <p>Let $\mathbf{A} =$</p> | $\begin{bmatrix} \mathbf{W} & \mathbf{H} \\ \mathbf{J} & \mathbf{K} \end{bmatrix} \begin{matrix} r \\ m-r \\ r & n-r \end{matrix}$ |
| | <p>$\mathbf{C} =$</p> | $\begin{bmatrix} \mathbf{W} \\ \mathbf{J} \end{bmatrix} \quad \mathbf{B} = \begin{bmatrix} \mathbf{W} & \mathbf{H} \end{bmatrix}$ |

185 THEOREM 5.2. The r by r matrix \mathbf{W} also has rank r . Then $\mathbf{A} = \mathbf{CW}^{-1}\mathbf{B}$.

186 1. Combinations V of the rows of B must produce the dependent rows in
 187 $\begin{bmatrix} \mathbf{J} & \mathbf{K} \end{bmatrix}$
 188

189 Then $\begin{bmatrix} \mathbf{J} & \mathbf{K} \end{bmatrix} = \mathbf{V}\mathbf{B} = \begin{bmatrix} \mathbf{V}\mathbf{W} & \mathbf{V}\mathbf{H} \end{bmatrix}$ for some matrix \mathbf{V} and $\mathbf{C} =$
 190 $\begin{bmatrix} \mathbf{I} \\ \mathbf{V} \end{bmatrix} \mathbf{W}$

191
 192 **2.** Combinations \mathbf{T} of the columns of \mathbf{C} must produce the dependent columns
 193 in $\begin{bmatrix} \mathbf{H} \\ \mathbf{K} \end{bmatrix}$

194
 195 Then $\begin{bmatrix} \mathbf{H} \\ \mathbf{K} \end{bmatrix} = \mathbf{C}\mathbf{T} = \begin{bmatrix} \mathbf{W}\mathbf{T} \\ \mathbf{J}\mathbf{T} \end{bmatrix}$ for some matrix \mathbf{T} and $\mathbf{B} = \mathbf{W} \begin{bmatrix} \mathbf{I} & \mathbf{T} \end{bmatrix}$

196
 197 **3.** $\mathbf{A} = \begin{bmatrix} \mathbf{W} & \mathbf{H} \\ \mathbf{V}\mathbf{W} & \mathbf{V}\mathbf{H} \end{bmatrix} = \begin{bmatrix} \mathbf{W} & \mathbf{W}\mathbf{T} \\ \mathbf{V}\mathbf{W} & \mathbf{V}\mathbf{W}\mathbf{T} \end{bmatrix} = \begin{bmatrix} \mathbf{I} \\ \mathbf{V} \end{bmatrix} \begin{bmatrix} \mathbf{W} \end{bmatrix} \begin{bmatrix} \mathbf{I} & \mathbf{T} \end{bmatrix} = \mathbf{C}\mathbf{W}^{-1}\mathbf{B}$

198
 199 Each step **1, 2, 3** proves again that \mathbf{W} is invertible. If its rank were less than r , then
 200 \mathbf{W} could not be a factor of \mathbf{C} or \mathbf{B} or \mathbf{A} . In case \mathbf{C} and \mathbf{B} are not in the first r
 201 columns and rows of \mathbf{A} , permutations P_R and P_C will give $P_R \mathbf{A} P_C = \begin{bmatrix} \mathbf{W} & \mathbf{H} \\ \mathbf{J} & \mathbf{K} \end{bmatrix}$
 202 and the proof goes through.

203

204 **Uniqueness** The columns of \mathbf{C} are a basis for the column space of \mathbf{A} . Then the
 205 columns of \mathbf{R} contain the unique coefficients that express each column of \mathbf{A} in that
 206 basis: $\mathbf{A} = \mathbf{C}\mathbf{R}$. But \mathbf{W} and \mathbf{B} are not unique. They can come from any r independent
 207 rows of \mathbf{A} .

208 This is consistent with (and proves!) the uniqueness of the reduced row echelon
 209 form R_0 , and the non-uniqueness of the steps from \mathbf{A} to R_0 .

210

211 $\mathbf{A} = \mathbf{C}\mathbf{R}$ and $\mathbf{A} = \mathbf{C}\mathbf{W}^{-1}\mathbf{B}$ both lead to explicit expressions for the *pseudoinverse*
 212 \mathbf{A}^+ . Those factors have full rank, so we know that $\mathbf{A}^+ = \mathbf{R}^+\mathbf{C}^+$ and $\mathbf{A}^+ = \mathbf{B}^+\mathbf{W}\mathbf{C}^+$.
 213 These factors are one-sided inverses as in $\mathbf{R}^+ = \mathbf{R}^T(\mathbf{R}\mathbf{R}^T)^{-1}$ and $\mathbf{C}^+ = (\mathbf{C}^T\mathbf{C})^{-1}\mathbf{C}^T$.

214 We don't know the full history of this column-row factorization of \mathbf{A} into
 215 $\mathbf{C}\mathbf{W}^{-1}\mathbf{B}$. But an excellent paper by Hamm and Huang [3] extends the theory to
 216 the general case when \mathbf{W}^{-1} becomes a pseudoinverse. They analyze *CUR* approxi-
 217 mations for large matrices (see below), and they have also provided a valuable set of
 218 references—including the 1956 paper [8] by Roger Penrose. That paper followed his
 219 1955 introduction of the Moore-Penrose pseudoinverse.

220 **6. Applications to Large Matrices.** Our starting point for $\mathbf{A} = \mathbf{C}\mathbf{R}$ has
 221 been its connection to the reduced row echelon form $R_0 = \mathbf{rref}(\mathbf{A})$. That is tradi-
 222 tionally a classroom construction, and classrooms are usually limited to small matri-
 223 ces. But the idea of using actual columns and rows of the matrix \mathbf{A} can be highly
 224 attractive. Those vectors have meaning. They are often *sparse and/or nonnega-*
 225 *tive*. They reflect useful properties that we wish to preserve in approximating \mathbf{A} .

226

227 A persuasive and widely read essay by Lee and Seung [5] made this point strongly.
 228 And there has been good progress on the key decision: *Which k columns and k rows*
 229 *to choose?* It is understood that for large matrices, k may be much less than the rank
 230 r of A .

231 An early start was the analysis of “pseudoskeleton approximations” to an exact
 232 skeleton $A = CW^{-1}B$. This can be accurate (provided there exists a rank k matrix
 233 close to A). The optimal choice of columns and rows for C and B is connected in [2]
 234 to finding a subdeterminant of maximum volume.

235 Interpolative decomposition has become part of *randomized* numerical linear alge-
 236 bra.

237 Exact factors C, W, B would be submatrices of A and require no new storage.
 238 But if A itself is very large, its size will have to be reduced. The computation often
 239 samples the columns of A by $Y = AG$, for an n by k Gaussian random matrix G .
 240 Then a crucial step is to produce an orthonormal basis Q for that k -dimensional ap-
 241 proximate column space (at reasonable cost). This is Gram-Schmidt with column
 242 pivoting, aiming to put the most important columns first: $AP = QR$.

243 The whole randomized algorithm is beautifully presented by Martinsson in [6]
 244 and in the extended survey [7] by Martinsson and Tropp. When the singular val-
 245 ues of A decay rapidly, these algorithms succeed with a moderate choice of k —the
 246 computational rank that replaces the actual rank of A .

247 **7. Implementation.** An `rref` function has always been part of MATLAB. It
 248 was intended for academic use in computer experiments supplementing a traditional
 249 linear algebra course, and has been largely ignored by most users for many years. We
 250 were pleasantly surprised when the function proved to be exactly what was needed to
 251 compute the $A = CR$ and ultimately the $A = CW^{-1}B$ factorizations.

252 **7.1. rref.** This description of `rref` was probably written in the 1980’s. It is
 253 still valid today.

254 `R = rref(A)` produces the reduced row echelon form of A .

255 `[R,jb] = rref(A)` also returns a vector, `jb`, so that:
 256 `r = length(jb)` is this algorithm’s idea of the rank of A ,
 257 `x(jb)` are the bound variables in a linear system, $Ax = b$,
 258 `A(:,jb)` is a basis for the range of A ,
 259 `R(1:r,jb)` is the r -by- r identity matrix.

260 An important output is the indices `jb` of the leading r independent columns of A . The
 261 statement

262 `[~,jb] = rref(A)`

263 ignores the echelon form and just retains the column indices.

264 **7.2. cr and cab.** We have developed two new functions, `cr` and `cab`. The
 265 function `cr` produces two matrices. The first is C , a basis for the column space. The
 266 second output is R , the `rref` of A with any zero rows removed.

267 `[C,R] = cr(A)`

268 `rref` is used once by `cr` to find both C and R . Here is the code for `cr`.

```
269 [R,jb] = rref(A);
270 r = length(jb);      % r = rank.
271 R = R(1:r,:);      % R(:,jb) == eye(r).
```

272 `C = A(:,jb)`

273 The function `cab` produces three matrices, `C`, `W` and `B`, with

274 `[C,W,B] = cab(A)` or `[C,W,B,cols,rows] = cab(A)`

275 `C = A(:,cols)` is a subset of the columns and forms a basis for the column space.

276 The same `C` is produced by `cr`. `B = A(rows,:)` is a subset of the rows and forms a

277 basis for the row space. And `W = A(rows,cols)` is the set intersection of `C` and `B`.

278 The original `A` can be reconstructed with

279 `A = C*inv(W)*B`

280 Moreover the rank of `A` is `r = rank(A) = length(cols) = length(rows)`.

281 Consequently `W` is square and its size is `r` -by- `r`

282 `rref` is used twice by `cab`. The first use finds a basis for the column space. The second

283 use is with the transpose of `A` and finds a basis for the row space. The echelon forms

284 themselves are discarded; only the pivot indices are retained. Here is the code for `cab`.

285 `[~,cols] = rref(A); % Column space`

286 `C = A(:,cols);`

287 `[~,rows] = rref(A'); % Row space`

288 `B = A(rows,:);`

289 `W = A(rows,cols); % Intersection`

291 The theorem that the column space and the row space have the same dimension implies

292 that `cols` and `rows` have the same length and ultimately that `W` is square and nonsingular.

293 lar.

294 Its size is the rank.

295 **7.3. Gauss-Jordan.** The algorithm used by `rref` is known in numerical analysis
 296 as Gauss-Jordan elimination. When Gauss-Jordan is applied to an augmented matrix
 297 $[A \ b]$, it solves the linear system $A*x = b$. The Gaussian elimination algorithm is
 298 less work. It does not zero elements above the pivot. It stops with an upper triangular
 299 factor U and then solves the modified augmented system by back substitution.

300 Most modern computers have a fused multiply-add instruction, `FMA`, that multiplies
 301 two floating point values and then adds a third value to the result in one operation.
 302 For an n -by- n system, Gauss-Jordan requires about $(1/2)*n^3$ FMAs, while Gaussian
 303 elimination requires less, only $(1/3)*n^3$ FMAs. This is one of the reasons why the
 304 numerical linear algebra community has been less interested in `rref`.

305 **7.4. Examples.** Magic squares provide good examples of our factorizations. An
 306 n -by- n magic square is a matrix whose elements are the integers from 1 to n^2 and
 307 whose row sums, column sums, and sums along both principal diagonals are all the
 308 same magic sum. Many magic squares are rank deficient.

309 A 4-by-4 magic square is one of the mathematical objects in Melencolia I, a 1514
 310 engraving by the German Renaissance artist Albrecht Dürer.

311 $A =$

| | | | | |
|-----|----|----|----|----|
| 312 | 16 | 3 | 2 | 13 |
| 313 | 5 | 10 | 11 | 8 |
| 314 | 9 | 6 | 7 | 12 |
| 315 | 4 | 15 | 14 | 1 |

316

317 This matrix is rank deficient, but the linear dependencies are not obvious. The rank
 318 is revealed to be three by the `cr` function.

319 $[C,R] = cr(A)$

320

321 $C =$

| | | | | |
|-----|----|----|----|--|
| 322 | | | | |
| 323 | 16 | 3 | 2 | |
| 324 | 5 | 10 | 11 | |
| 325 | 9 | 6 | 7 | |
| 326 | 4 | 15 | 14 | |

327

328 $R =$

| | | | | |
|-----|---|---|---|----|
| 329 | | | | |
| 330 | 1 | 0 | 0 | 1 |
| 331 | 0 | 1 | 0 | -3 |
| 332 | 0 | 0 | 1 | 3 |

333 The matrix R is the `rref` of A with a zero row removed.

334 Equal treatment of rows and columns is provided by the `cab` function.

335 $[C,W,B] = cab(A) = C*inv(W)*B.$

336 $C =$

| | | | |
|-----|----|----|----|
| 337 | 16 | 3 | 2 |
| 338 | 5 | 10 | 11 |
| 339 | 9 | 6 | 7 |
| 340 | 4 | 15 | 14 |

```

341
342 W =
343   16   3   2
344    5  10  11
345    9   6   7
346
347 B =
348   16   3   2  13
349    5  10  11   8
350    9   6   7  12

```

351 The Franklin *semimagic* square is attributed to Benjamin Franklin in c.1752 [1].
 352 Its rows and columns all have the required magic sum, but the diagonals do not, so it
 353 isn't fully magic. However, many other interesting submatrices are magic, including
 354 bent diagonals and any eight elements arranged symmetrically about the center.

```

355 A =
356   52   61   4   13   20   29   36   45
357   14   3   62   51   46   35   30   19
358   53   60   5   12   21   28   37   44
359   11   6   59   54   43   38   27   22
360   55   58   7   10   23   26   39   42
361    9   8   57   56   41   40   25   24
362   50   63   2   15   18   31   34   47
363   16   1   64   49   48   33   32   17

```

364 It is hard to think of Franklin's square as a linear transformation, but `cab` can still
 365 compute its rank.

```

366 [C,W,B] = cab(A)
367 C =
368   52   61   4
369   14   3   62
370   53   60   5
371   11   6   59
372   55   58   7
373    9   8   57
374   50   63   2
375   16   1   64
376 W =
377   52   61   4
378   14   3   62
379   53   60   5
380 B =
381   52   61   4   13   20   29   36   45
382   14   3   62   51   46   35   30   19
383   53   60   5   12   21   28   37   44

```

384 So the rank is three. And, `C*inv(W)*B` is within roundoff error of `A`.

```

385 test = C*inv(W)*B

```

```
386 test =
387 52.00 61.00 4.00 13.00 20.00 29.00 36.00 45.00
388 14.00 3.00 62.00 51.00 46.00 35.00 30.00 19.00
389 53.00 60.00 5.00 12.00 21.00 28.00 37.00 44.00
390 11.00 6.00 59.00 54.00 43.00 38.00 27.00 22.00
391 55.00 58.00 7.00 10.00 23.00 26.00 39.00 42.00
392 9.00 8.00 57.00 56.00 41.00 40.00 25.00 24.00
393 50.00 63.00 2.00 15.00 18.00 31.00 34.00 47.00
394 16.00 1.00 64.00 49.00 48.00 33.00 32.00 17.00
```

395 Rounding `test` to the nearest integers reproduces Franklin's magic square exactly.

396

REFERENCES

- 397 [1] Benjamin Franklin, *The Papers of Benjamin Franklin* 4: 1750-1753,
 398 To Peter Collinson (American Philosophical Society and Yale University) 392.
 399 <https://franklinpapers.org/framedVolumes.jsp?vol=4&page=392a>.
 400 Also printed in *Experiments and Observations on Electricity* (London, 1769), 350-354.
- 401 [2] S. A. Goreinov, E. E. Tyrtyshnikov, and N. L. Zamarashkin, *A theory of pseu-*
 402 *doskeleton approximation*, *Linear Algebra and Its Applications* **261** (1997) 1-21.
- 403 [3] K. Hamm and L. Huang, *Perspectives on CUR decompositions*, *Applied and Comp. Harmonic*
 404 *Analysis* **48** (2020) 1088-1099 and arXiv: 1907.12668.
- 405 [4] R. Kannan and S. Vempala, *Randomized algorithms in numerical linear algebra*, *Acta Nu-*
 406 *merica* **26** (2017) 95-135.
- 407 [5] D. D. Lee and H. S. Seung, *Learning the parts of objects by nonnegative matrix factorization*,
 408 *Nature* **401** (1999) 788-791.
- 409 [6] P. G. Martinsson, *Randomized methods for matrix computations*, in *The Mathematics of*
 410 *Data*, American Mathematical Society (2018).
- 411 [7] Per-Gunnar Martinsson and Joel Tropp, *Randomized numerical linear algebra: Foundations*
 412 *and Algorithms*, *Acta Numerica* (2020) and arXiv: 2002.01387.
- 413 [8] R. Penrose, *On best approximate solutions of linear matrix equations*, *Math. Proc. Cambridge*
 414 *Phil. Soc.* **52** (1956) 17-19.
- 415 [9] R. Piziak and P. L. Odell, *Full rank factorization of matrices*, *Mathematics Magazine* **72**
 416 (1999) 193-201.
- 417 [10] D. C. Sorensen and Mark Embree, *A DEIM-induced CUR factorization*, *SIAM J. Scientific*
 418 *Computing* **38** (2016) 1454-1482 and arXiv: 1407.5516.
- 419 [11] Gilbert Strang, *Linear Algebra for Everyone*, Wellesley-Cambridge Press (2020).
- 420 [12] Gilbert Strang and Steven Lee, *Row reduction of a matrix and $A = CaB$* , *American Mathe-*
 421 *matical Monthly* **107** (2000) 681-688.
- 422 [13] S. Voronin and P. G. Martinsson, *Efficient algorithms for CUR and interpolative*
 423 *matrix decompositions*, *Advances in Computational Mathematics* **43** (2017) 495-516.
- 424

425 **Acknowledgments.** We thank David Vogan for translating $CW^{-1}B$ into alge-
 426 bra in equation (5.3), and Dan Drucker, Steven Lee, and Alexander Lin for a stream
 427 of ideas and references about CR and $CW^{-1}B$. A joint paper on their construction
 428 and applications is in preparation.