### Intro to Generating Functions

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### Outline

**Ordinary Generating Functions** 

Exponential Generating Functions (EGF)

Rational Generating Functions

The Exponential Formula

# Ordinary Generating Functions (OGF)

#### Definition

The formal series

$$F(x) := \sum_{n \ge 0} f(n) x^n \in \mathbb{C}[[x]]$$

associated to the counting map  $f: \mathbb{N}_0 \to \mathbb{C}$  is called *ordinary* generating function. We also use the notation  $[x^n]F(x) := f(n)$ .

### Sum and Product of OGF

We recall that  $\mathbb{C}[[x]]$  is an integral domain (actually a PID). We can perform sums and multiplications of OGF according to the way we sum and multiply elemnents in  $\mathbb{C}[[x]]$ .

### Definition

Let  $F(x) = \sum_{n \ge 0} f(n)x^n$  and  $G(n) = \sum_{n \ge 0} g(n)x^n$  be two OGF. Then their *sum* is the OGF

$$F(x) + G(x) := \sum_{n \ge 0} (f(n) + g(n))x^n,$$

and their product, also called their convolution is the OGF

$$F(x)G(x) := \sum_{n>0} \left( \sum_{k=0}^{n} f(k)g(n-k) \right) x^{n}.$$

**Example:** The OGF of the Fibonacci sequence is  $F(x) = \frac{1}{1-x-x^2}$ .

### The Inverse of an OGF

#### Theorem

An OGF F(x) is invertible (meaning that there exists an OGF G(x) such that F(x)G(x) = 1) iff is  $F(0) \neq 0$ .

**Proof:** If G(x) is the inverse of F(x) then F(0)G(0)=1 and so  $F(0)\neq 0$ . If  $F(0)\neq 0$  then  $G(0)=F(0)^{-1}$ , and we can recurrently define the remaining coefficients of G(0) by using the multiplication formula for OGF.

**Example:** The OGF F(x) = 1 - x satisfies that  $F(0) = 1 \neq 0$ . Therefore, it is invertible. The inverse of F(x) is  $G(x) = \sum_{n \geq 0} x^n$ .

## Convergence of OGF

#### Definition

The *degree* of an OGF  $F(x) = \sum_{n \geq 0} f(n)x^n$ , denoted by deg F(x) is the smallest n such that  $f(n) \neq 0$ . A sequence of OGF  $\{F_i(x)\}_{i \in \mathbb{N}}$  *converges* to the OGF F(x) if

$$\lim_{i\to\infty}\deg(F(x)-F_i(x))=\infty.$$

We say that the infinite sum  $\sum_{i\geq 0} F_i$  coverges to the OGF F(x) if the sequence of partial sum converges to F(x).

#### Theorem

The infinite series  $\sum_{i>0} F_i(x)$  converges iff  $\lim_{i\to\infty} \deg F_i(x) = \infty$ .

**Proof:** It follows directly from the definition of convergence.

# Composition of OGF

#### Definition

Let  $F(x) = \sum_{n \geq 0} f(n)x^n$  and  $G(n) = \sum_{n \geq 0} g(n)x^n$  be two OGF such that G(0) = 0. Then the composition of F(x) and G(x) is the OGF

$$F(G(x)) := \sum_{n>0} f(n)G(x)^n.$$

#### Remarks:

- Notice that the condition G(0)=0 guarantees that  $\sum_{n\geq 0}f(n)G(x)^n$  converges. This is because  $\deg G(x)^n\geq n\deg G(x)$ , and so  $\lim_{i\to\infty}\deg(f(i)G(x)^i)=\infty$ .
- ▶ The expression  $e^{1+x} = \sum_{n\geq 0} (x+1)^n/n!$  is not a valid OGF because the series does not converge in the sense we defined above.

# A Few Popular OGFs

### $\mathsf{Theorem}$

The following are popular and useful OGFs:

$$1. \sum_{n\geq 0} x^n = \frac{1}{1-x},$$

2. 
$$\sum_{n\geq 0} (-1)^n x^n = \frac{1}{1+x}$$

3. 
$$\sum_{n\geq 0} x^{2n} = \frac{1}{1-x^2}$$
,

4. 
$$\sum_{n\geq 0} {m \choose n} x^n = (1+x)^m$$

5. 
$$\sum_{n\geq 0} {n+m \choose n} x^n = \frac{1}{(1-x)^{m+1}}$$
,

6. 
$$\sum_{n\geq 0} \binom{n}{m} x^n = \frac{x^m}{(1-x)^{m+1}}$$
.

Sketch of Proof: Pending...

# **Exponential Generating Functions**

#### **Definition**

The formal series

$$F(x) := \sum_{n \ge 0} f(n) \frac{x^n}{n!}$$

associated to the counting map  $f: \mathbb{N}_0 \to \mathbb{C}$  is called the exponential generating function (EGF) of f. We also use the notation  $[x^n/n!]F(x) := f(n)$ .

## Sum, Product, and Composition of EGFs

#### Definition

Let  $F(x) = \sum_{n \geq 0} f(n) x^n / n!$  and  $G(n) = \sum_{n \geq 0} g(n) x^n / n!$  be two EGF. Then their *sum* is the EGF

$$F(x) + G(x) := \sum_{n \ge 0} (f(n) + g(n))x^n/n!,$$

and their product, also called their convolution is the OGF

$$F(x)G(x) := \sum_{n>0} \Big( \sum_{k=0}^{n} \binom{n}{k} f(k)g(n-k) \Big) x^{n} / n!.$$

If, in addition, G(0) = 0, we composition of the EGFs F and G is given by

$$F(G(x)) := \sum_{n>0} f(n)G(x)^n/n!.$$

### The formal derivative of a GF

### Definition

- ► The *derivative* of the OGF  $F(x) = \sum_{n\geq 0} f(n)x^n$  is  $F'(x) := \sum_{n\geq 0} nf(n)x^{n-1}$ .
- ► The *derivative* of the EGF  $F(x) = \sum_{n\geq 0} f(n)x^n/n!$  is  $F'(x) := \sum_{n\geq 0} f(n)x^{n-1}$ .

### Theorem

Let F(x) and G(x) be two OGF (EGF) then the following hold:

- (F(x) + G(x))' = F'(x) + G'(x),
- ► F(x)G(x) = F'(x)G(x) + F(x)G'(x),
- (F(G(x))' = F'(G(x))G'(x).

### The formal derivative of a GF

#### **Theorem**

Let F(x) and G(x) be two OGFs such that F(0) = 1 and G(0) = 0. If G'(x) = F'(x)/F(x) then  $F(x) = \exp(G(x))$ , where  $\exp(G(x)) = \sum_{n \geq 0} \frac{G(x)^n}{n!}$ .

Sketch of Proof: Pending...

**Example 1:** The EGF of the function  $f: \mathbb{N} \to \mathbb{C}$  given by f(0) = 1 and f(n+1) = f(n) + nf(n-1) if  $n \ge 0$ 

### General Newton Coefficients

### Definition

- ▶ For  $\lambda \in \mathbb{C}$  and  $k \in \mathbb{N}_0$ , set  $\binom{\lambda}{k} := \lambda(\lambda 1) \dots (\lambda k + 1)$ .
- ▶ For an OGF F(x) such that F(0) = 0, we define:

$$(1+F(x))^{\lambda}:=\sum_{n\geq 0}\binom{\lambda}{n}F(x)^n.$$

**Example 2:** We want to find all  $f: \mathbb{N}_0 \to \mathbb{R}$  satisfying

$$\sum_{k=0}^{n} f(k)f(n-k) = 1.$$

# A Practice Example

**Example:** Suppose that the function  $f: \mathbb{N} \to \mathbb{C}$  has EGF  $F(x) = e^{x+x^2/2}$ .

- Find a recurrence formula for f.
- ► Find an explicit formula for f.

# Rational Generating Functions

#### Theorem

Let  $f: \mathbb{N}_0 \to \mathbb{C}$  and  $Q(x) = 1 + c_1 x + \dots + c_d x^d = \prod_{i=1}^k (1 - \alpha_i x)^{d_i}$ , where  $c_1, \dots, c_d \in \mathbb{C}$   $(c_d \neq 0)$ . Then TFAE:

- 1.  $f(n+d) + c_1 f(n+d-1) + \cdots + c_d f(n) = 0$  for every  $n \in \mathbb{N}_0$ ;
- 2.  $F(x) = \sum_{n>0} f(b)x^n = P(x)/Q(x)$ , where deg P(x) < d;
- 3.  $F(x) = \sum_{n>0} f(n)x^n = \sum_{i=1}^k g_i(x)/(1-\alpha_i x)^{d_i};$
- 4.  $f(n) = \sum_{i=1}^{k} p_i(n)\alpha_i^n$ , where  $p_i(n)$  is a polynomial in n such that  $\deg p_i < d_i$  for each  $i \in \{1, \ldots, k\}$ .

**Sketch of Proof:** For each  $i \in \{1, 2, 3, 4\}$ , define the complex space

$$V_i := \{ f : \mathbb{N}_0 \to \mathbb{C} \mid f \text{ satisfies } (i) \}.$$

Check that dim  $V_i = d$  for each i. Use this to check that  $V_1 = V_2$  and  $V_3 = V_4$ . Finally, show that  $V_3 \subseteq V_2$ .

# Rational Generating Functions (continuation)

#### Definition

A generating function  $F(x) = \sum_{n \geq 0} f(n)x^n$  satisfying any of the four condition in the previous theorem is called a (proper) *rational* generating function.

### Examples

**Example 1:** Let f(n) be the number of paths with n non-intersecting steps starting from (0,0) with directions east, north, or west.

- 1. Find the generating function of F of f.
- 2. Find a close formula for f.

**Hint:** Count the paths of length n ending in EE, WW, and NE.

**Example 2:** Write  $(\sqrt{2}+\sqrt{3})^{1980}$  in decimal form. What is the last digit before and the first digit after the decimal point? **Hint:** Compute the generating function of  $(\sqrt{2}+\sqrt{3})^{2n}$ .

## Putting Structures on Finite Sets

#### **Theorem**

Let  $f_1, \ldots, f_n \colon \mathbb{N}_0 \to \mathbb{C}$ , and denote by  $E_{f_i}(x)$  the EGF of  $f_i$ . For every finite set S, let

$$h(|S|) = \sum_{(T_1,...,T_n)} f_1(|T_1|) ... f_n(|T_n|),$$

where the sum runs over every ordered n-partition of S. Then the EGF  $E_h(x)$  of h satisfies that  $E_h(x) = E_{f_1}(x) \dots E_{f_n}(x)$ .

**Sketch of Proof:** Suppose first that n=2. If |S|=s, the fact that there are  $\binom{s}{k}$  ordered partitions  $(T_1, T_2)$  such that  $|T_1|=k$  of S implies that

$$h(s) = \sum_{k=0}^{s} {s \choose k} f_1(k) f_2(n-k).$$

Then  $E_h(x) = E_{f_1}(x)E_{f_2}(x)$ . Now extend to n by induction.

## The Compositional Formula

#### **Theorem**

Given  $f: \mathbb{N} \to \mathbb{C}$  and  $g: \mathbb{N}_0 \to \mathbb{C}$  with g(0) = 1, and for every finite set S let

$$h(0) = h([n]) = \sum_{\{T_1, \dots, T_k\} \in \pi([n])} f(|T_1|) \dots f(|T_k|) g(k)$$

if |S| > 0 and h(0) = 1. Then  $E_h(x) = E_g(E_f(x))$ .

**Sketch of Proof:** Defining, for every  $k \in \{1, ..., n\}$ 

$$h_k(n) = \frac{1}{k!} \sum_{(T_1,...,T_k)} f(|T_1|) ... f(|T_k|) g(k),$$

we have  $h(n) = \sum h_k(n)$ . By the previous theorem,  $E_{h_k}(x) = g(k)/k!E_f(x)^k$ . Hence

$$E_h(x) = \sum_{k>1} g(k) \frac{E_f(x)^k}{k!} = E_g(E_f(x)).$$

# The Compositional Formula: An Example

**Example:** In how many ways h(n) we can form n people into nonempty lines, and then arrange these lines in a circular order?

**Explanation:** Let f(n) and g(n) the number of ways to form n people in a line and in a circle, respectively. Then f(n) = n! and g(n) = (n-1)!. So the EGFs of f and g are

$$E_f(x) = \sum_{n \ge 1} x^n = \frac{x}{1-x}$$
 and  $E_g(x) = \sum_{n \ge 1} \frac{x^n}{n} = \ln(1-x)^{-1}$ .

Hence, using the previous theorem,

$$E_h(x) = E_g(E_f(x)) = \ln(\frac{1-x}{1-2x}) = \sum_{n>1} (2^n-1)(n-1)! \frac{x^n}{n!}.$$

Thus 
$$h(n) = (2^n - 1)(n - 1)!$$
.

### References

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