MIT 18.211: COMBINATORIAL ANALYSIS

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Lecture 23: Intro to Trees

In this lecture, we introduce trees and discuss some basic related properties.

Definition 1. A simple connected graph is called a *tree* if it does not contain any cycle.

A connected simple graph G is said to be *minimally connected* if any graph obtained from G by deleting an edge is disconnected. As the following proposition indicates, trees are precisely those graphs that are minimally connected.

Proposition 2. A graph is a tree if and only if it is minimally connected.

Proof. Let G be a graph.

To prove the direct implication, assume that G is a tree and suppose, towards a contradiction, that G is not minimally connected. Then G contains an edge e = vwthat we can delete to obtain a connected graph G'. Since G' is connected, there must be a path $v_1v_2...v_\ell$ in G' connecting v to w (in particular, $v_1 = v$ and $v_\ell = w$). Since $vw \notin V(G')$, we see that $v_1v_2 \ldots v_\ell v_1$ is a cycle in G, which yields the desired contradiction.

To establish the converse, assume that G is minimally connected and suppose, by way of contradiction, that G is not a tree. Then G contains a cycle $v_1v_2 \ldots v_\ell v_1$. Let G' be the graph that we obtain from G by removing the edge v_1v_2 . We claim that G' is a connected graph. To prove this, let v and w be two distinct vertices of G', and let Pbe a path in G from v to w, which must exist because G is connected. If v_1v_2 is not an edge of P, then P is a path from v to w in G'. Otherwise, we can replace the edge v_1v_2 in P by the path $v_1v_\ell v_{\ell-1}\ldots v_2$ to turn P into a walk from v to w in G'. Among all such walks, pick one P_0 of minimum length. It is clear that P_0 is a path from v to w in G'. Hence we conclude that G' is connected, which contradicts that G is minimally connected.

As we show in the next proposition a tree with n vertices must contain precisely n-1 edges.

Proposition 3. If G is a tree, then |E(G)| = |V(G)| - 1.

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Proof. We will argue that for every tree *G*, the equality |E(G)| = |V(G)| - 1 by using induction on |V(G)|. If |V(G)| = 1, then *G* is the graph consisting of one vertex and no edges, and so the equality |E(G)| = |V(G)| - 1 holds immediately. Assume that every tree with $n \in \mathbb{N}$ vertices has precisely n - 1 edges, and let *G* be a tree with |V(G)| = n + 1. Let *P* be a path in *G* with maximum number of edges. Since *G* is connected and $|V(G)| \ge 2$, it follows that deg $v \ge 1$ for every $v \in V(G)$. Thus, *P* has at least two vertices, and the maximality of *P* implies that the first and last vertices of *P* have degree 1. Take $v \in V(G)$ such that deg v = 1. Observe that the graph *G'* we obtain from *G* after removing *v* (and the only edge in *G* that is adjacent to *v*) is a tree with |V(G')| = n. By our induction hypothesis, |E(G')| = n - 1 and, therefore |E(G)| = |E(G')| + 1 = n = |V(G)| - 1. □

If we drop connectedness from the definition of a tree, we obtain a class of graphs that are also quite relevant. They are called forests.

Definition 4. A *forest* is a graph whose connected components are trees.

We can generalize Proposition 3 to forests as follows.

Proposition 5. Let G be a forest with exactly k connected components. Then |E(G)| = |V(G)| - k.

Proof. ¹ Let T_1, \ldots, T_k be the connected components of G. Now construct a graph G' from G be adding k-1 edges e_1, \ldots, e_{k-1} such that the edge e_i connect a vertex of T_i with a vertex of T_{i+1} . It is clear that G' is a connected graph with no cycles. Hence G' is a tree, and it follows from Proposition 3 that |E(G')| = |V(G')| - 1. This implies that |E(G)| = |E(G')| - (k-1) = |V(G')| - k = |V(G)| - k.

PRACTICE EXERCISES

Exercise 1. [1, Exercise 10.29] Prove that in any tree, any two longest paths cross each other.

Exercise 2. [1, Exercise 10.31] Let T be a tree, and let P(T) be the induced subgraph of T whose vertices of T in the intersection of all longest paths of T. Describe the graph P(T).

Exercise 3. Let G be a connected graph satisfying |E(G)| = |V(G)| - 1. Prove that G is a tree.

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¹Observe that this prove also follows mimicking the inductive proof given for trees.

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References

[1] M. Bóna: A Walk Through Combinatorics: An Introduction to Enumeration and Graph Theory (Fourth Edition), World Scientific, New Jersey, 2017.

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