

# The fine art of point counting

Andrew V. Sutherland

Massachusetts Institute of Technology

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Simons Collaboration in Arithmetic Geometry, Number Theory, and Computation

# What is arithmetic geometry?

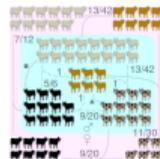
Arithmetic geometers study solutions to polynomial equations like



Babylonians  
circa 1800 BCE

$$x^2 + y^2 = z^2,$$

$$x^2 - 4729494y^2 = 1,$$

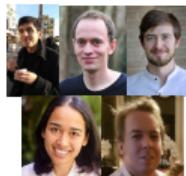


Archimedes  
251 BCE

$$x^3 + y^3 + z^3 = 3,$$

and even “cursed” equations like

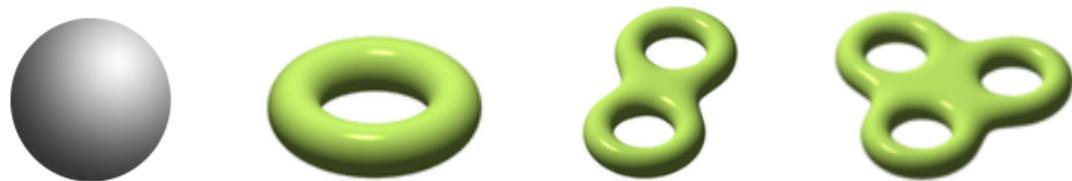
$$y^4 + 5x^4 - 6x^2y^2 + 6x^3z + 26x^2yz + 10xy^2z - 10y^3z \\ - 32x^2z^2 - 40xyz^2 + 24y^2z^2 + 32xz^3 - 16yz^3 = 0.$$



Balakrishnan et al.  
2017

## Which solutions?

There is a robust theory (classical algebraic geometry) to address this problem over  $\mathbb{C}$ .



But number theorists are typically interested in solutions over rings like  $\mathbb{Z}$  or  $\mathbb{Q}$ . These can be very hard to find. Indeed, it took sixty-five years to find

$$569936821221962380720^3 + (-569936821113563493509)^3 + (-472715493453327032)^3 = 3.$$

This problem is **undecidable**, in general (this is Hilbert's 10th problem).

We can simplify the problem by looking for solutions modulo primes.

In other words, we can **count points** over finite fields  $\mathbb{F}_p$  (or  $\mathbb{F}_q$ ).

These point counts can be used to define **zeta functions** and **L-functions** that encode the underlying structure of an equation (or system of equations) in a canonical way.

## Counting points modulo $p$

Let's count points  $(x, y)$  on the curve  $C: x^2 + y^2 = 1$  modulo primes  $p$ :

|       |   |   |   |   |    |    |    |    |    |    |           |
|-------|---|---|---|---|----|----|----|----|----|----|-----------|
| $p$   | 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | ...       |
| <hr/> |   |   |   |   |    |    |    |    |    |    |           |
|       | 2 | 4 | 4 | 8 | 12 | 12 | 16 | 20 | 24 | 28 | $p \pm 1$ |

Better, count points  $(x, y, z) \sim (cx, cy, cz)$  on  $x^2 + y^2 = z^2 \pmod{p}$ :

|       |   |   |   |   |    |    |    |    |    |    |         |
|-------|---|---|---|---|----|----|----|----|----|----|---------|
| $p$   | 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | ...     |
| <hr/> |   |   |   |   |    |    |    |    |    |    |         |
|       | 3 | 4 | 6 | 8 | 12 | 14 | 18 | 20 | 24 | 30 | $p + 1$ |

We always get  $p + 1$ . The  $L$ -function of  $C$  is

$$L(C, s) = \prod (1 - p^{-s})^{-1} = \sum n^{-s} = \zeta(s),$$

and we will get the same  $L$ -function whenever  $C$  has [genus 0](#).

# Elliptic curves

Let  $E$  be an elliptic curve over  $\mathbb{Q}$ , which we can write as

$$E: y^2 = x^3 + ax + b.$$

Every curve of **genus 1** with rational points has this form.  
You use elliptic curves every day! (on your phone/tablet/laptop)



The number of points on  $E$  modulo  $p$  can be written as

$$\#E_p(\mathbb{F}_p) = p + 1 - a_p,$$

where the **trace of Frobenius**  $a_p$  satisfies  $|a_p| \leq 2\sqrt{p}$ .



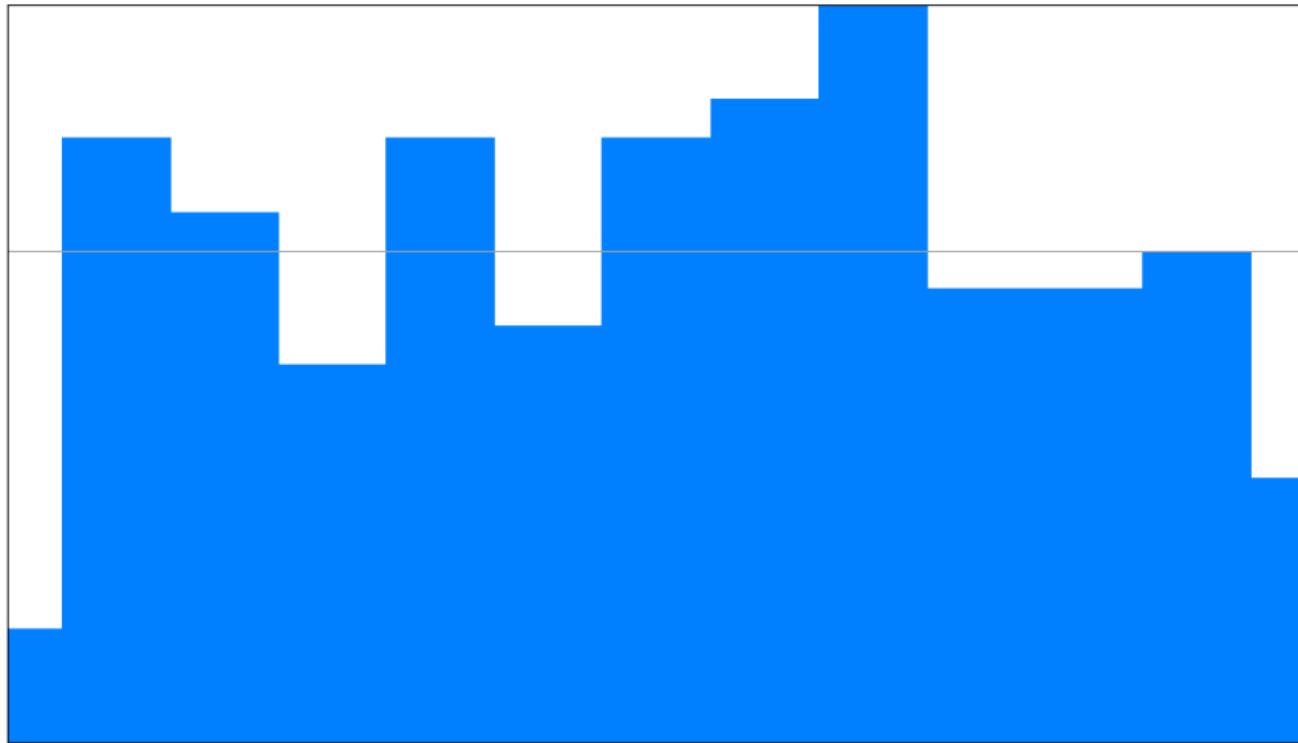
Frobenius

Let us now consider the sequence of real numbers  $x_p := -t_p/\sqrt{p} \in [-2, 2]$ .

Example:  $y^2 = x^3 + x + 1$

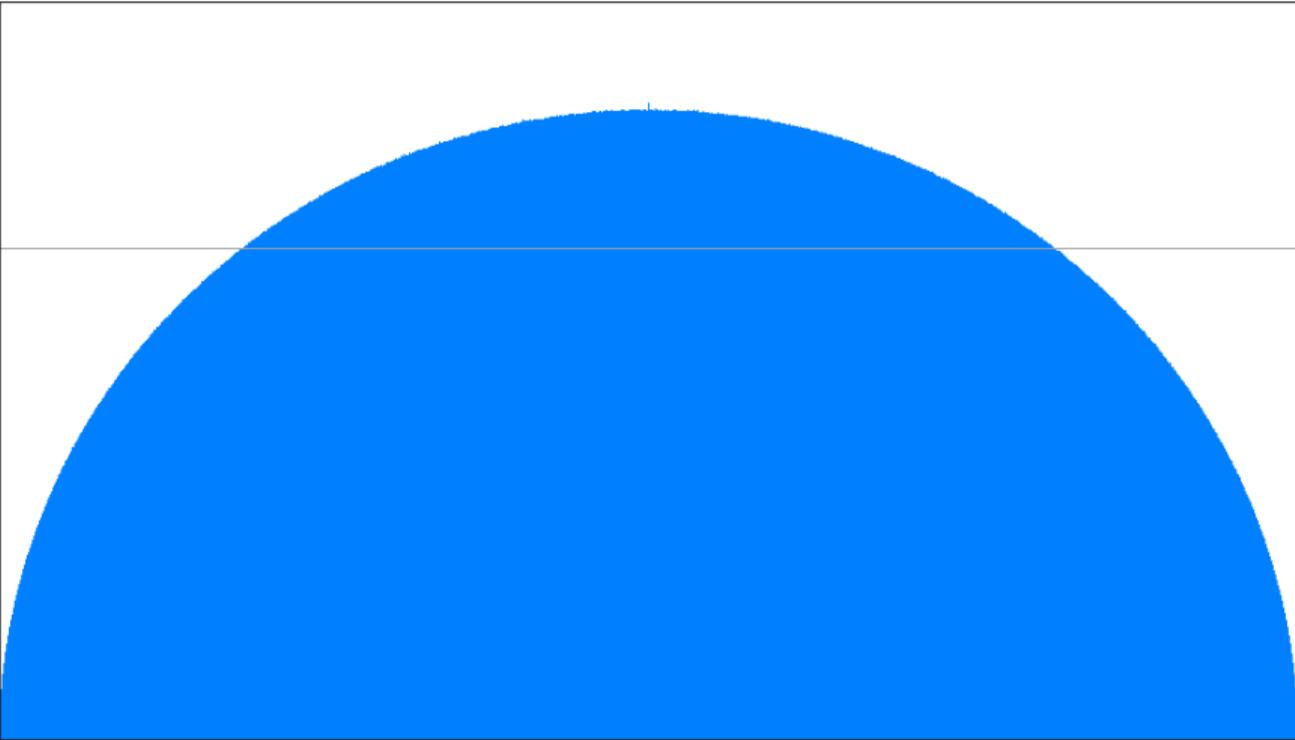
| $p$ | $t_p$ | $x_p$            | $p$ | $t_p$ | $x_p$            | $p$ | $t_p$ | $x_p$            |
|-----|-------|------------------|-----|-------|------------------|-----|-------|------------------|
| 3   | 0     | <b>0.000000</b>  | 71  | 13    | <b>-1.542816</b> | 157 | -13   | <b>1.037513</b>  |
| 5   | -3    | <b>1.341641</b>  | 73  | 2     | <b>-0.234082</b> | 163 | -25   | <b>1.958151</b>  |
| 7   | 3     | <b>-1.133893</b> | 79  | -6    | <b>0.675053</b>  | 167 | 24    | <b>-1.857176</b> |
| 11  | -2    | <b>0.603023</b>  | 83  | -6    | <b>0.658586</b>  | 173 | 2     | <b>-0.152057</b> |
| 13  | -4    | <b>1.109400</b>  | 89  | -10   | <b>1.059998</b>  | 179 | 0     | <b>0.000000</b>  |
| 17  | 0     | <b>0.000000</b>  | 97  | 1     | <b>-0.101535</b> | 181 | -8    | <b>0.594635</b>  |
| 19  | -1    | <b>0.229416</b>  | 101 | -3    | <b>0.298511</b>  | 191 | -25   | <b>1.808937</b>  |
| 23  | -4    | <b>0.834058</b>  | 103 | 17    | <b>-1.675060</b> | 193 | -7    | <b>0.503871</b>  |
| 29  | -6    | <b>1.114172</b>  | 107 | 3     | <b>-0.290021</b> | 197 | -24   | <b>1.709929</b>  |
| 37  | -10   | <b>1.643990</b>  | 109 | -13   | <b>1.245174</b>  | 199 | -18   | <b>1.275986</b>  |
| 41  | 7     | <b>-1.093216</b> | 113 | -11   | <b>1.034793</b>  | 211 | -11   | <b>0.757271</b>  |
| 43  | 10    | <b>-1.524986</b> | 127 | 2     | <b>-0.177471</b> | 223 | -20   | <b>1.339299</b>  |
| 47  | -12   | <b>1.750380</b>  | 131 | 4     | <b>-0.349482</b> | 227 | 0     | <b>0.000000</b>  |
| 53  | -4    | <b>0.549442</b>  | 137 | 12    | <b>-1.025229</b> | 229 | -2    | <b>0.132164</b>  |
| 59  | -3    | <b>0.390567</b>  | 139 | 14    | <b>-1.187465</b> | 233 | -3    | <b>0.196537</b>  |
| 61  | 12    | <b>-1.536443</b> | 149 | 14    | <b>-1.146925</b> | 239 | -22   | <b>1.423062</b>  |
| 67  | 12    | <b>-1.466033</b> | 151 | -2    | <b>0.162758</b>  | 241 | 22    | <b>-1.417145</b> |

a1 histogram of  $y^2 = x^3 + x + 1$  for  $p \leq 2^{10}$   
170 data points in 13 buckets,  $z_1 = 0.029$ , out of range data has area 0.018



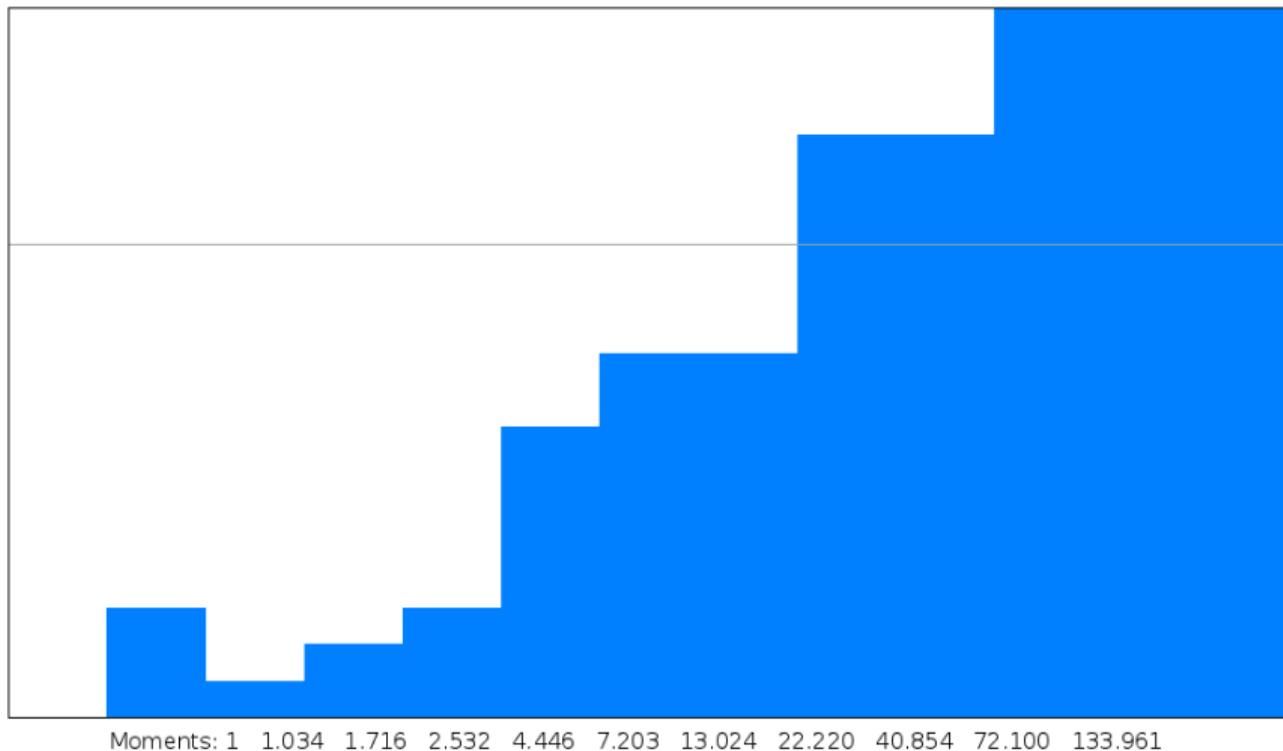
Moments: 1 0.051 1.039 0.081 2.060 0.294 4.971 1.134 13.278 4.308 37.954

a1 histogram of  $y^2 = x^3 + x + 1$  for  $p \leq 2^{40}$   
41203088794 data points in 202985 buckets

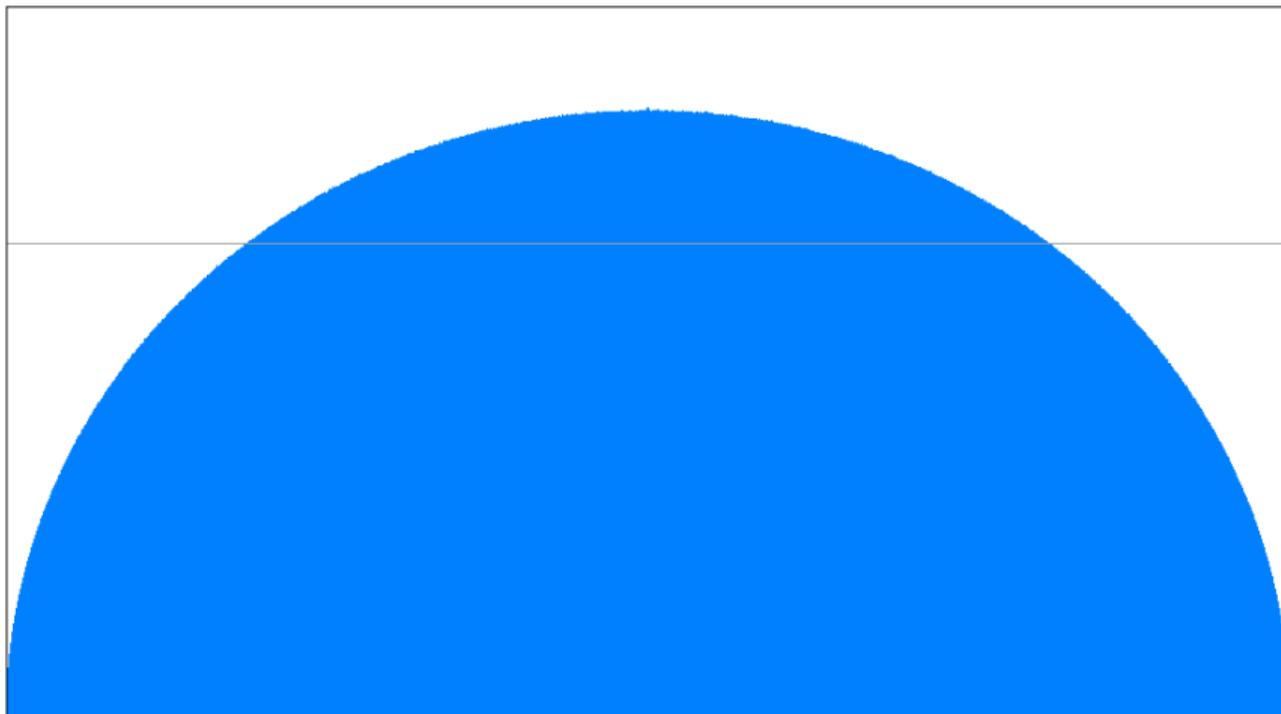


Moments: 1 0.000 1.000 0.000 2.000 0.000 5.000 0.000 14.000 0.000 41.999

a1 histogram of  $y^2 + xy + y = x^3 - x^2 - 20067762415575526585033208209338542750930230312178956502x + 34481611795030556467032985690390720374855944359319180361266008296291939448732243429$  for  $p \leq 2^{10}$   
172 data points in 13 buckets,  $z1 = 0.023$ , out of range data has area 0.250



a1 histogram of  $y^2 + xy + y = x^3 - x^2 - 20067762415575526585033208209338542750930230312178956502x$   
+ 34481611795030556467032985690390720374855944359319180361266008296291939448732243429 for  $p \leq 2^{40}$   
41203088796 data points in 202985 buckets



Moments: 1 0.000 1.000 0.000 2.000 0.000 5.000 0.001 14.000 0.003 42.000

# The Sato-Tate conjecture

The **Sato-Tate conjecture** states that, except for certain families of well understood exceptions, we will always get the same limiting distribution as  $p \rightarrow \infty$ .



Mikio Sato



John Tate

## Theorem (Taylor et al. 2008)

*Let  $E/\mathbb{Q}$  be an elliptic curve without any extra endomorphisms.  
The sequence  $x_p$  converges to the semi-circular distribution.*



Richard Taylor

Richard Taylor received the 2014 Breakthrough Prize in Mathematics for this work.

# Ranks of elliptic curves

Elliptic curves are not just curves, they are also **abelian varieties**. For  $E/\mathbb{Q}$  the rational points are finitely generated:  $E(\mathbb{Q}) \simeq T \oplus \mathbb{Z}^r$ , where  $T$  is finite and  $r$  is the **rank**.

There are many things we do not know about  $r$ :

- ▶ Is there an algorithm to compute  $r$ ?
- ▶ Which values of  $r$  can occur? Is there an upper limit?
- ▶ How often does each possible value of  $r$  occur, on average?

Theorem (Elkies 1990)

*The value of  $r$  can be as large as 28.*



Theorem (Bhargava-Shankar 2012)

*The average value of  $r$  lies strictly between 0 and 1.*



## Conjectures of Birch and Swinnerton-Dyer

Based on extensive computer experiments (in the early 1960s!), Birch and Swinnerton-Dyer conjectured that

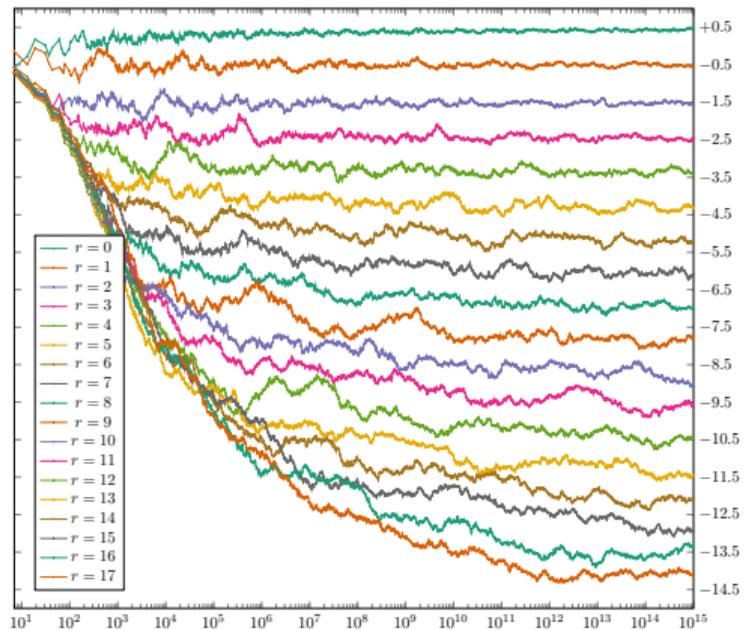
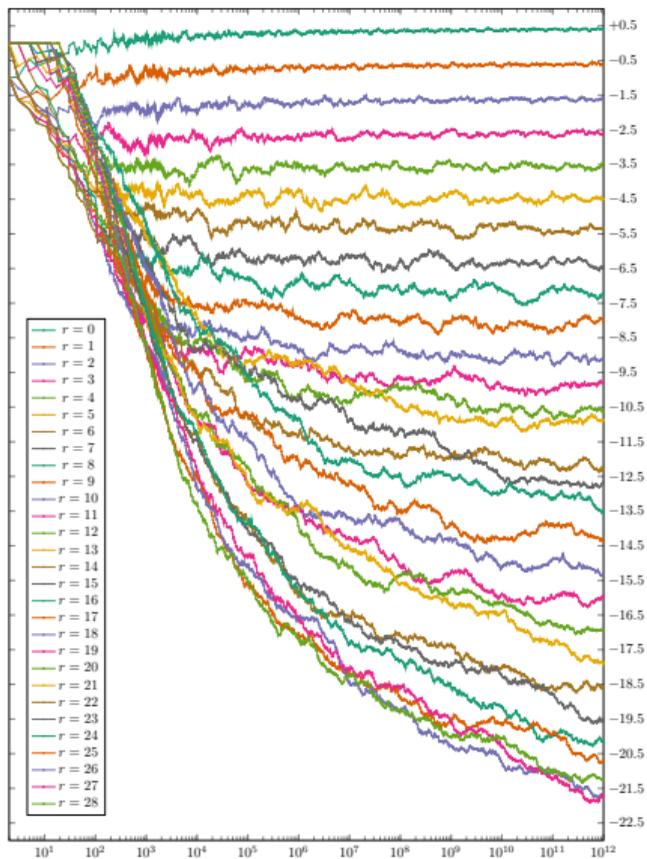
$$\lim_{x \rightarrow \infty} \prod_{p \leq x} \frac{\#E(\mathbb{F}_p)}{p} = c_E (\log x)^r$$

for some constant  $c_E$ , where  $r$  is the rank. This implies

$$\lim_{x \rightarrow \infty} \frac{1}{\log x} \sum_{p \leq x} \frac{a_p \log p}{p} = -r + \frac{1}{2},$$

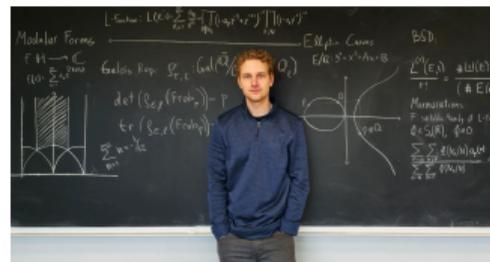
and suggests that one can compute (or at least predict)  $r$  by counting points.

Birch and Swinnerton-Dyer eventually formulated a weaker conjecture relating the order of vanishing of  $L(E, s)$  at  $s = 1$  to  $r$  (part of the [BSD conjecture](#)).



# Murmurations of elliptic curves

In 2022, He, Lee, Oliver, and Pozdnyakov ran a series of machine learning experiments in an attempt to predict ranks of elliptic curves over  $\mathbb{Q}$  using Frobenius traces.



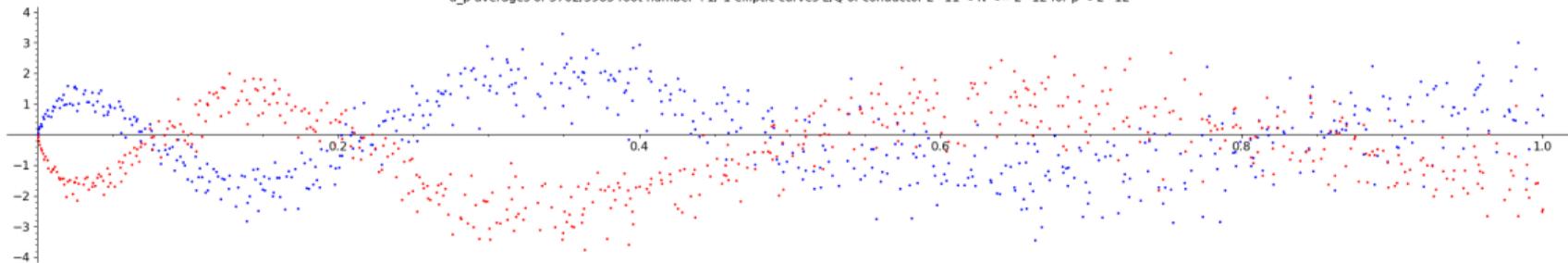
Their efforts to predict ranks were not particularly successful, but in the process they discovered a new and as yet unexplained oscillation in average Frobenius traces in families of elliptic curves ordered by conductor and separated by rank.

You can read more about their discovery in this recent [Quanta article](#).

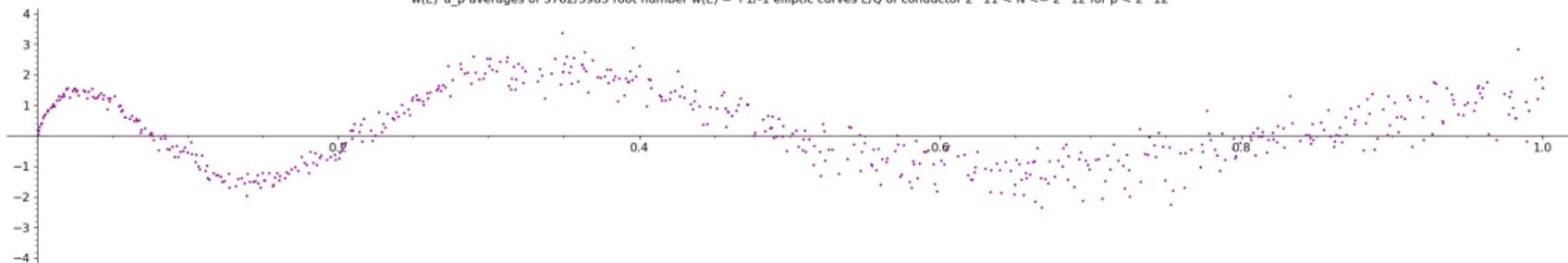
# Murmurations of elliptic curves

Elliptic curves of conductor  $N \in (2^n, 2^{n+1}]$  for  $11 \leq n \leq 18$ . Blue/red/purple dots at  $(p, \bar{a}_p$  or  $\bar{m}_p)$  are averages of  $a_p$  or  $m_p := (-1)^r a_p(E)$  over even/odd/all  $E/\mathbb{Q}$ .

$a_p$  averages of 3762/3985 root number +1/-1 elliptic curves  $E/\mathbb{Q}$  of conductor  $2^{11} < N \leq 2^{12}$  for  $p < 2^{12}$

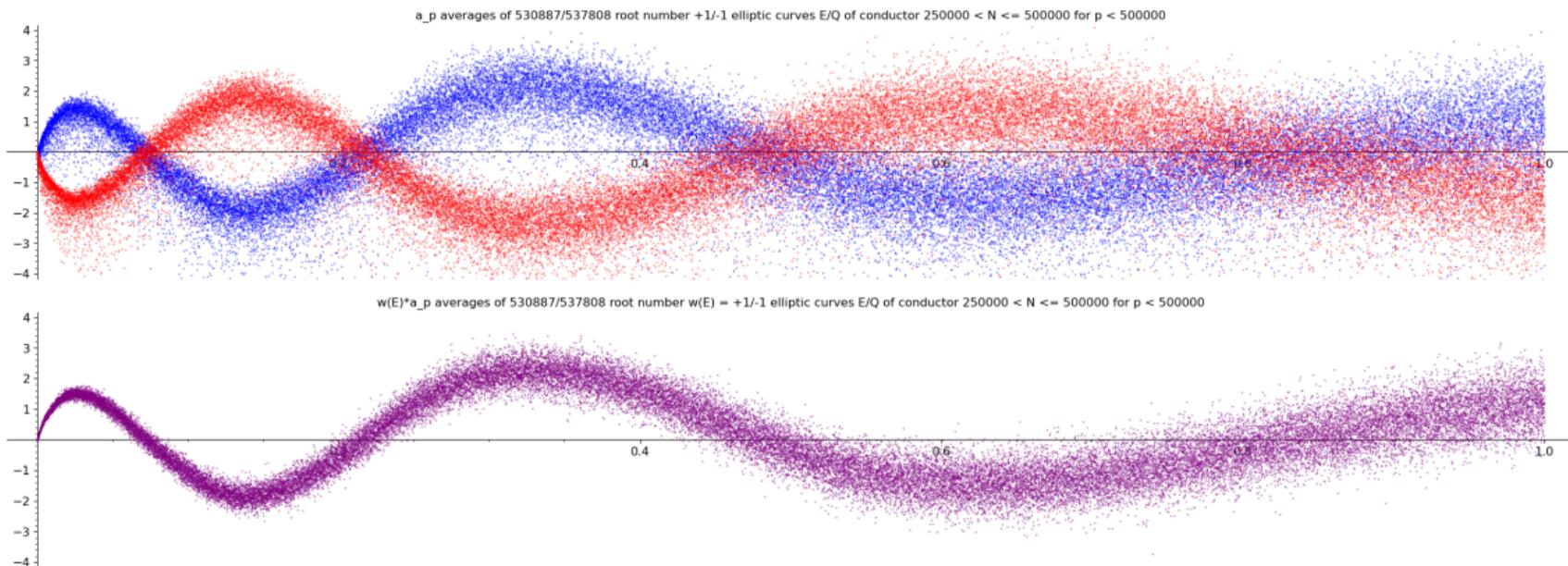


$w(E) \cdot a_p$  averages of 3762/3985 root number  $w(E) = +1/-1$  elliptic curves  $E/\mathbb{Q}$  of conductor  $2^{11} < N \leq 2^{12}$  for  $p < 2^{12}$



# Murmurations of elliptic curves

Elliptic curves of conductor  $N \in (2^n, 2^{n+1}]$  for  $11 \leq n \leq 18$ . Blue/red/purple dots at  $(p, \bar{a}_p$  or  $\bar{m}_p)$  are averages of  $a_p$  or  $m_p := (-1)^r a_p(E)$  over even/odd/all  $E/\mathbb{Q}$ .



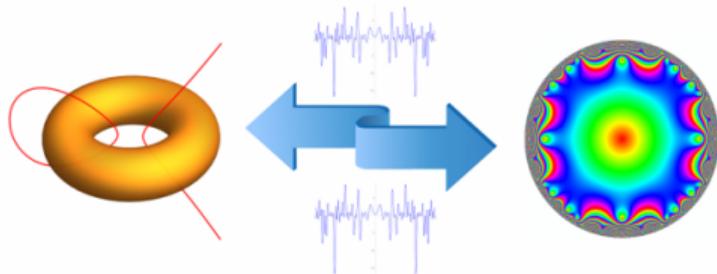
# Modularity

The proof of the Sato-Tate conjecture is built on the [Modularity Theorem](#).

[Theorem \(Taylor-Wiles 1995, Breuil-Conrad-Diamond-Taylor 2001\)](#)

*For every elliptic curve  $E/\mathbb{Q}$  there is a modular form  $f_E$  for which  $L(E, s) = L(f_E, s)$ .*

*The  $q$ -expansion  $f_E(q) = \sum a_n q^n$  of  $f_E$  is determined by the Frobenius traces  $a_p$  of  $E$ .*



[Corollary \(Wiles 1995\)](#)

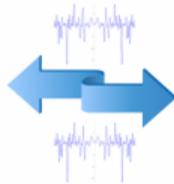
*The equation  $x^n + y^n = z^n$  has no nontrivial integer solutions for  $n > 2$ .*

# The $L$ -functions and Modular Forms Database (LMFDB)

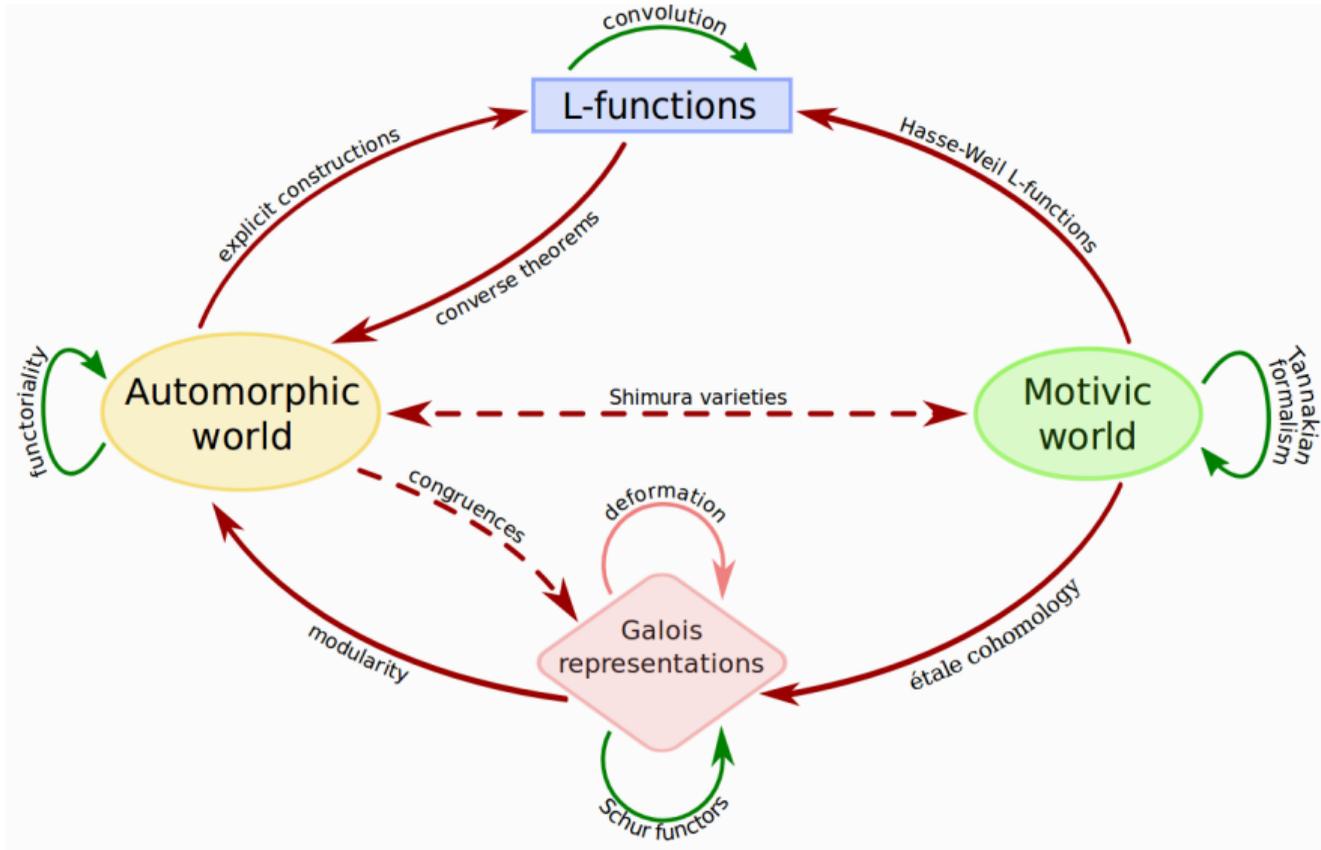
The relationship between elliptic curves and modular forms established by the Modularity Theorem was conjectured fifty years earlier.

Compelling evidence for this conjecture was obtained over many decades by tabulating elliptic curves and modular forms and computing their  $L$ -functions.

Extensive tables of these (and many other mathematical objects) are now available in the  $L$ -functions and Modular Forms Database.



# The Langlands Program



## L-functions and zeta functions

The *L-function* of a (nice) curve  $X/\mathbb{Q}$  can be written as

$$L(X, s) := \prod L_p(p^{-s})^{-1} = \sum a_n n^{-s}.$$

The *completed L-function*  $\Lambda(s) := \Gamma_{\mathbb{C}}(s)L(X, s)$  is expected to satisfy

$$\Lambda(s) = \pm N^{1-s} \Lambda(2-s).$$

For good primes  $p$  the *L-polynomial*  $L_p(T)$  appears in the *zeta function*

$$Z(X_p; T) := \exp \left( \sum_{r \geq 1} \#X_p(\mathbb{F}_{p^r}) \frac{T^r}{r} \right) = \frac{L_p(T)}{(1-T)(1-pT)}.$$

The Langlands conjectures imply an effective *multiplicity one* result for  $L(X, s)$ , which is determined by  $\#X(\mathbb{F}_p)$  good primes  $p \leq B$ , where  $B$  depends on the *conductor*  $N$ .

## Algorithms to compute zeta functions

Given  $X/\mathbb{Q}$  of genus  $g$ , we want to compute  $L_p(T)$  for all good  $p \leq B$ .

| algorithm            | complexity per prime<br>(ignoring factors of $O(\log \log p)$ ) |                      |                      |
|----------------------|---|----------------------|----------------------|
|                      | $g = 1$   | $g = 2$              | $g = 3$              |
| point enumeration    | $p \log p$  | $p^2 \log p$         | $p^3 (\log p)^2$     |
| group computation    | $p^{1/4} \log p$  | $p^{3/4} \log p$     | $p (\log p)^2$       |
| $p$ -adic cohomology | $p^{1/2} (\log p)^2$  | $p^{1/2} (\log p)^2$ | $p^{1/2} (\log p)^2$ |
| CRT (Schoof-Pila)    | $(\log p)^5$  | $(\log p)^8$         | $(\log p)^{12}$      |
| average poly-time    | $(\log p)^4$  | $(\log p)^4$         | $(\log p)^4$         |

The bottom row is due to a 2014 breakthrough by David Harvey, and many later refinements [HS 2016, HS 2018, S 2020, CHS 2023].



## Timings for genus 3 curves

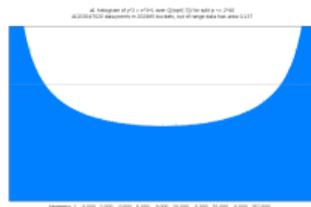
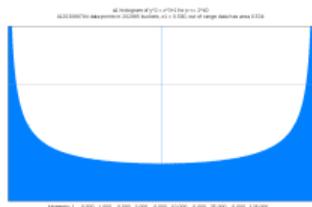
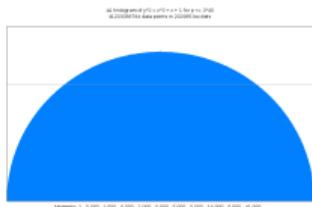
Time to compute  $L_p(T) \bmod p$  for all good  $p \leq B$ .

| $B$      | plane quartic (old) | plane quartic (new) | hyperelliptic (old) | hyperelliptic (new) |
|----------|---------------------|---------------------|---------------------|---------------------|
| $2^{12}$ | 18                  | 1.4                 | 1.3                 | 0.1                 |
| $2^{13}$ | 49                  | 2.4                 | 2.6                 | 0.2                 |
| $2^{14}$ | 142                 | 4.6                 | 5.4                 | 0.5                 |
| $2^{15}$ | 475                 | 8.4                 | 12                  | 1.0                 |
| $2^{16}$ | 1,670               | 18                  | 29                  | 2.1                 |
| $2^{17}$ | 5,880               | 34                  | 74                  | 5.3                 |
| $2^{18}$ | 22,300              | 71                  | 192                 | 14                  |
| $2^{19}$ | 78,100              | 159                 | 532                 | 37                  |
| $2^{20}$ | 297,000             | 421                 | 1,480               | 97                  |
| $2^{21}$ | 1,130,000           | 996                 | 4,170               | 244                 |
| $2^{22}$ | 4,280,000           | 2,430               | 12,200              | 617                 |
| $2^{23}$ | 16,800,000          | 5,930               | 36,800              | 1,500               |
| $2^{24}$ | 66,800,000          | 13,700              | 113,000             | 3,520               |
| $2^{25}$ | 244,000,000         | 31,100              | 395,000             | 8,220               |
| $2^{26}$ | 972,000,000         | 73,400              | 1,060,000           | 19,700              |

(Intel Xeon E7-8867v3 3.3 GHz CPU seconds).

# Sato-Tate groups and their distributions

There are two Sato-Tate distributions that arise for elliptic curves  $E/\mathbb{Q}$ , depending on whether  $E$  is exceptional or not, but over general number fields there are three:

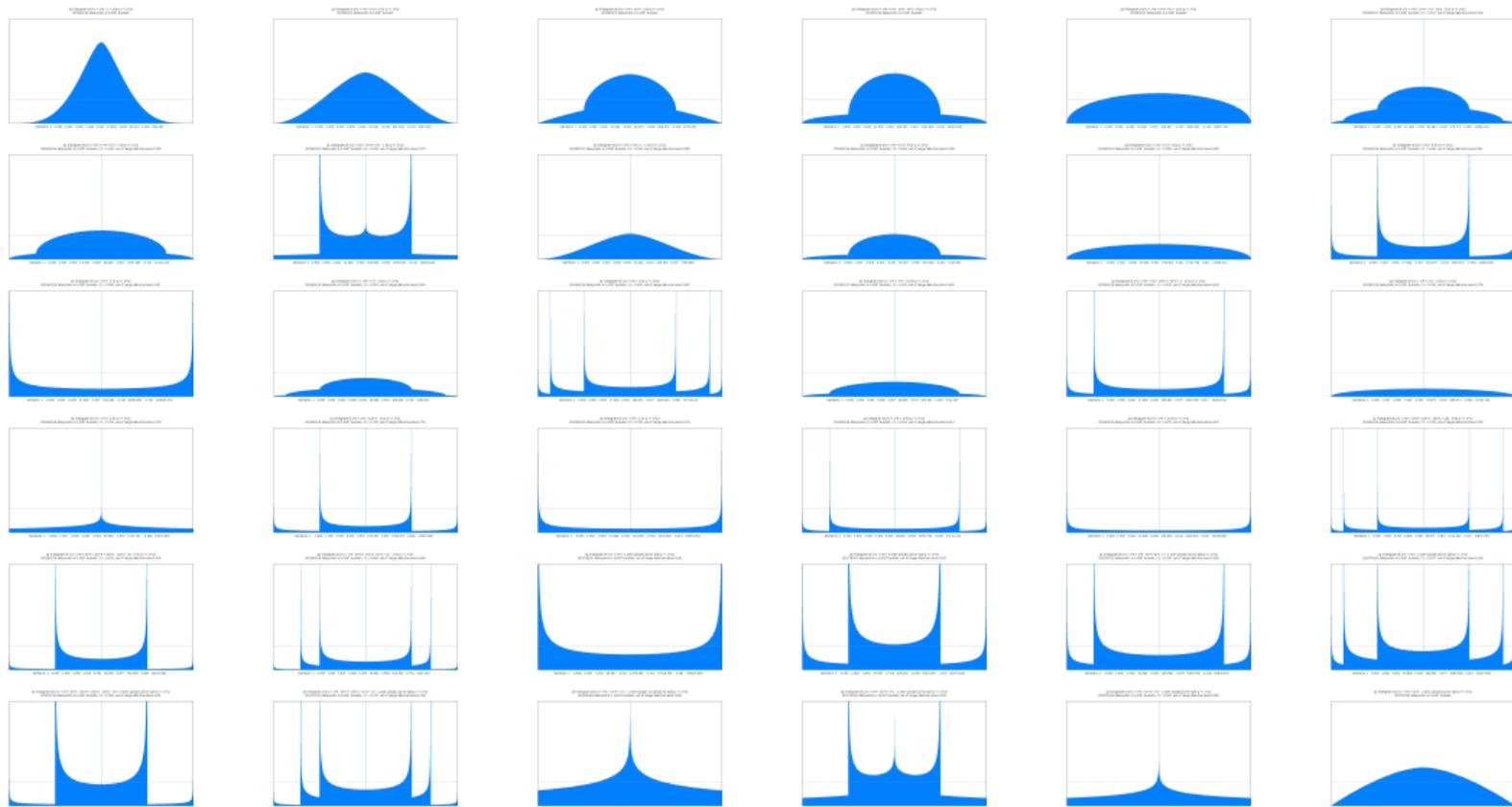


Each corresponds to the distributions of traces in a compact subgroup of  $SU(2)$ , the **Sato-Tate group** of  $E$ . Katz-Sarnak and Serre extended this random matrix model to curves of genus  $g$  and abelian varieties of dimension  $g$  using subgroups of  $USp(2g)$ .

The Sato-Tate conjecture is open for genus  $g > 1$ , but Sato-Tate groups have been completely classified for  $g \leq 3$ . There are 52 in genus 2 and 410 in genus 3 [Fité-Kedlaya-Rotger-S 2012, Fité-Kedlaya-S 2019].

These classifications involved more than a thousand CPU-years of computation.

# Sato-Tate trace distributions of genus 2 curves:



## Building a database of low genus curves

To make it feasible to compute  $L$ -functions, to make the Langlands correspondence explicit and to investigate murmurations, a we want to tabulate curves by **conductor**.

**No one knows how to do this for curves of genus  $g > 1$ , not even in principle!**

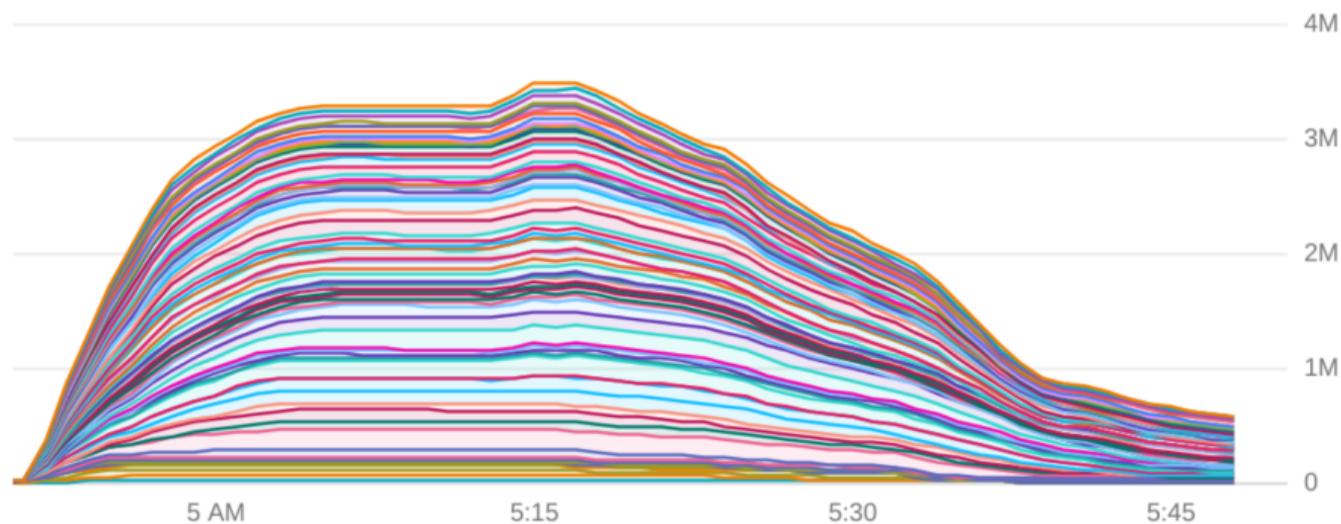
We instead “sieve the sky”. We enumerate vast numbers of curves with small coefficients along with their discriminants (which bound the conductor).

In our genus 2 and genus 3 searches we enumerated a total of more than  $10^{18}$  curves, keeping only about  $10^6$  curves of interest in each case.

To make such a computation feasible requires:

- efficient enumeration algorithms (equations/discriminants in parallel);
- code optimization down to the metal (about 10 nanoseconds per curve);
- massive parallelism (thousands or even millions of cores).

## Computation of Jan 9, 2022



We used 126080 32-vCPU instances in 73 data centers in 23 different locations, including Taiwan, Hong Kong, Tokyo, Osaka, Mumbai, Singapore, Sydney, Finland, Belgium, London, Frankfurt, the Netherlands, Zurich, Montréal, São Paulo, Iowa, South Carolina, Virginia, Oregon, Los Angeles, Salt Lake City, and Las Vegas.

## How much carbon does a 300 vCPU-year computation emit?

This is a question <http://www.green-algorithms.org/> can help answer.

300 vCPU-years is about 1 314 900 core-hours (2 vCPUs per core).

| CPU                 | cores | platform | location      | energy    | carbon    |
|---------------------|-------|----------|---------------|-----------|-----------|
| i9-9900K (64GB)     | 1     | desktop  | Massachusetts | 46.99 MWh | 19 750 Kg |
| i9-9900K (64GB)     | 16    | desktop  | Massachusetts | 17,61 MWh | 7 400 Kg  |
| Ryzen 3990X (256GB) | 64    | desktop  | Massachusetts | 7.44 MWh  | 3 260 Kg  |
| Ryzen 3990X (256GB) | 64    | cloud    | Virginia      | 8.60 MWh  | 2 650 Kg  |
| Ryzen 3990X (256GB) | 64    | cloud    | Montreal      | 8.60 MWh  | 13 Kg     |



## Some preliminary results

We found more than 500 genus 2 curves with conductors  $N \leq 1000$ , including the curve

$$C_{903} : y^2 + (x^2 + 1)y = x^5 + 3x^4 - 13x^3 - 25x^2 + 61x - 28$$

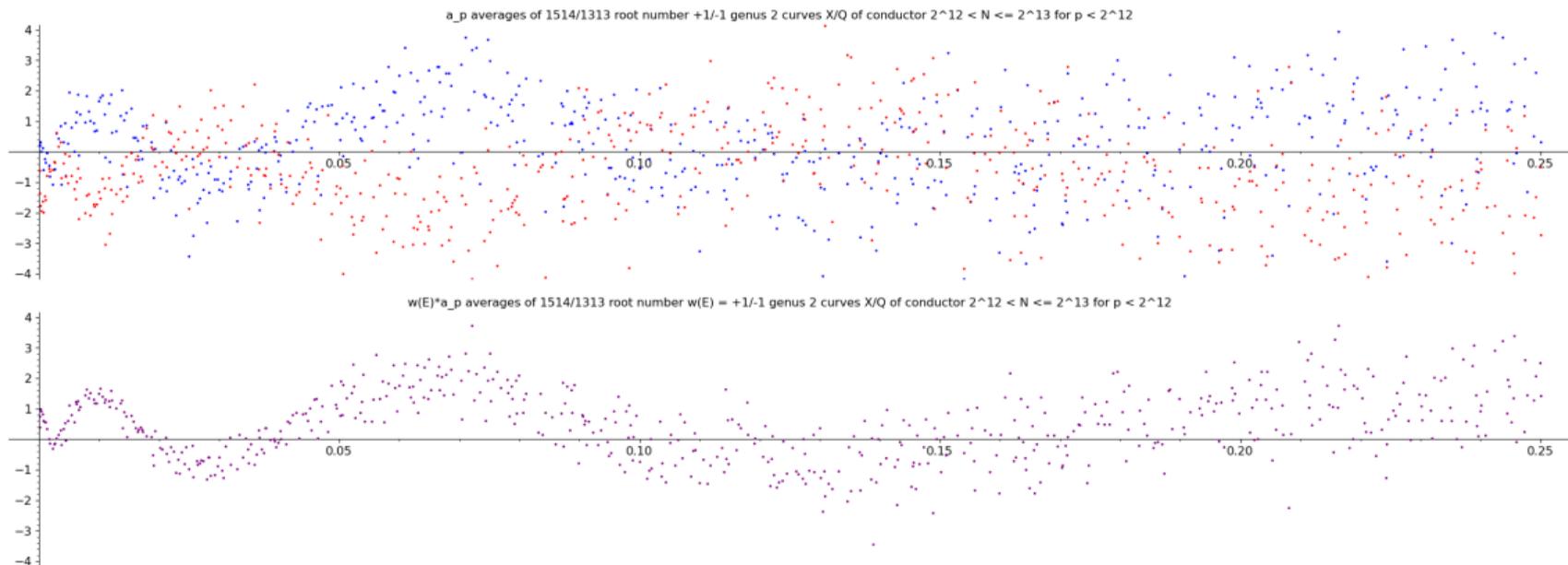
of conductor 903 and whose  $L$ -function coefficients match the  $a_n$  of the paramodular form of conductor 903 computed by Poor-Yuen and extended by Mellit to  $n \leq 100$ .

We also found curves of conductor 657, 760, 775, 924 not previously known to occur for Jacobians, examples of arithmetic phenomena that do not arise among the curves in the LMFDB, and many more curves of small conductor:

| conductor bound      | 1000 | 10000 | 100000 | 1000000 |
|----------------------|------|-------|--------|---------|
| curves in LMFDB      | 159  | 3069  | 20265  | 66158   |
| curves found         | 807  | 25438 | 426396 | 5059724 |
| L-functions in LMFDB | 109  | 2807  | 19775  | 65534   |
| L-functions found    | 200  | 9409  | 186811 | 2315328 |

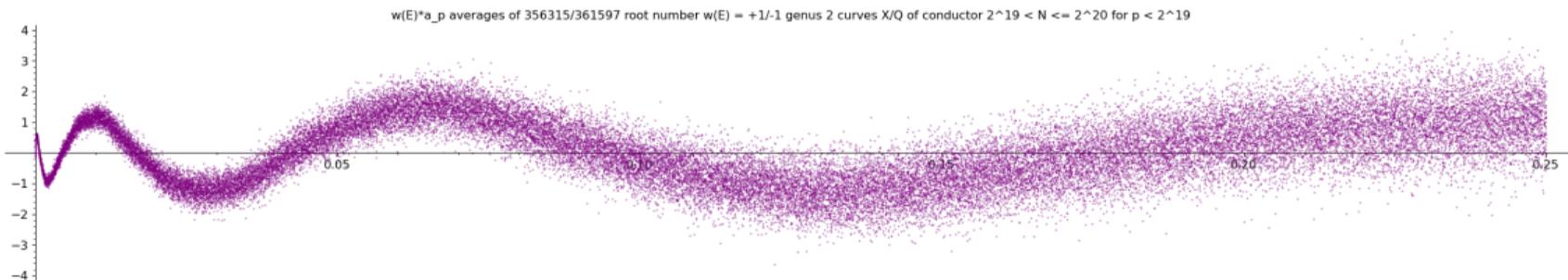
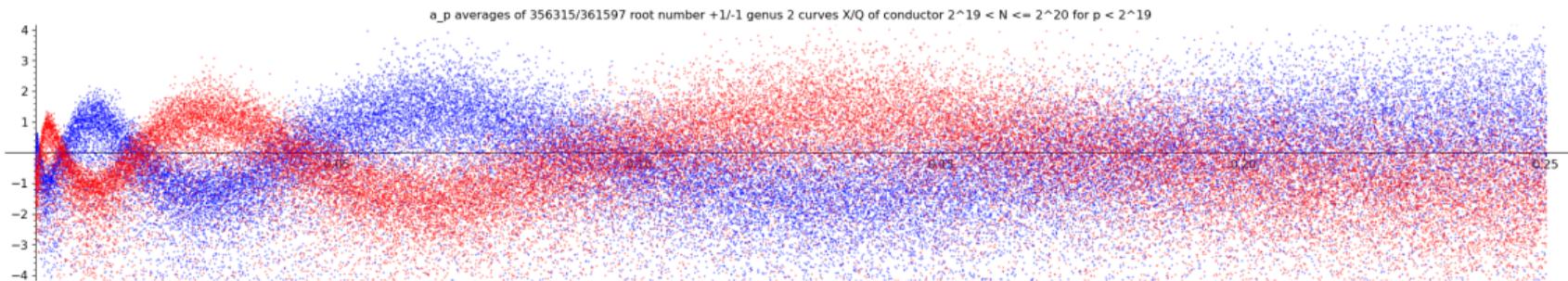
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Genus 2 curves  $X/\mathbb{Q}$  of conductor  $N \in (2^n, 2^{n+1}]$  for  $12 \leq n \leq 19$ . Blue/red/purple dots at  $(p, \bar{a}_p$  or  $\bar{m}_p)$  are averages of  $a_p$  or  $m_p := (-1)^r a_p(X)$  over even/odd/all  $X/\mathbb{Q}$ .



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