

A stochastic approach for noise stability on the hypercube

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Joint work with Ronen Eldan and Prasad Raghavendra

1. **Noise stability**
2. “Majority is Stablest”
3. A stochastic approach for noise stability via a re-normalized Brownian motion
4. Interlude - from a toy example to the Courtade-Kumar conjecture
5. Back to Majority is Stablest

Noise Operators

Consider the discrete hypercube $\mathcal{C}_n = \{-1, 1\}^n$ with its uniform probability measure μ .

For $\rho \in (0, 1)$ define the noise operator T_ρ , by

$$T_\rho f(x) = \mathbb{E}_{y \sim \rho \text{ correlated with } x} [f(y)].$$

We say that y is ρ correlated with x if $\mathbb{E}[y_i x_i] = \rho$. In other words, the law of y is the unique product measure with $\mathbb{E}[y] = \rho x$.

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For a Boolean function $f : \mathcal{C}_n \rightarrow \{-1, 1\}$, define its noise stability by,

$$\text{Stab}_\rho(f) := \mathbb{E}_\mu [fT_\rho f].$$

Important concept in social choice theory and Boolean analysis.

Example:

Theorem (Kalai 02')

If $f : \mathcal{C}_n \rightarrow \{-1, 1\}$ is used to rank three candidates,

$$\mathbb{P}_\mu (f \text{ gives a rational outcome}) = \frac{3}{4}(1 + \text{Stab}_{\frac{1}{3}}(f)).$$

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Question

Among all Boolean functions, which one maximizes the noise stability?

Easy answer: among all Boolean functions the dictator $f(x) := x_1$ has the largest noise stability.

Not a very useful fact in social choice theory.

Define the maximal influence of a Boolean function by:

$$\inf = \max_{i \in [n]} \mathbb{E}_{\mu} [(\partial_i f)^2].$$

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Among all Boolean functions **with small maximal influence**, which one maximizes the noise stability?

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Theorem (Mossel-O'Donnell-Oleszkiewicz 05')

Let f be a balanced Boolean function and suppose $\inf(f) \leq \kappa$, then,

$$\text{Stab}_\rho(f) \leq \frac{2}{\pi} \arcsin(\rho) + O\left(\frac{\log \log(\frac{1}{\kappa})}{\log(\frac{1}{\kappa})}\right).$$

Define the majority function $\text{Maj}_n(x) = \text{sgn}\left(\frac{1}{\sqrt{n}} \sum_{i=1}^n x_i\right)$.

- Computation: $\inf(\text{Maj}_n) \leq \frac{1}{\sqrt{n}}$.
- CLT: $\text{Stab}_\rho(\text{Maj}_n) \xrightarrow{n \rightarrow \infty} \frac{2}{\pi} \arcsin(\rho)$.

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Majority is Stablest - Proof Sketch

1. Prove analogous result in Gaussian space:
 - Noise semi-group is replaced by Ornstein-Uhlenbeck semi-group.
 - Majority is replaced by indicator of halfspace.

Result follows from the isoperimetric inequality.

2. Prove invariance principle for low-influence polynomials:

$$|\mathbb{E}_\mu[p] - \mathbb{E}_\gamma[p]| \leq O(2^{\text{degree}(p)} \cdot \text{inf}(p)).$$

3. Replace f by $T_\epsilon f$, essentially a log-degree polynomial.

Turns out that $\epsilon = \Theta\left(\frac{\log \log(\frac{1}{\kappa})}{\log(\frac{1}{\kappa})}\right)$ works.

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Quantitative Majority is Stablest

We prove a quantitative version of the Majority theorem.

Theorem

Let f be a balanced Boolean function and suppose $\inf(f) \leq \kappa$, then,

$$\text{Stab}_\rho(f) \leq \frac{2}{\pi} \arcsin(\rho) + \text{poly}(\kappa).$$

- The main idea is to realize $(\text{Stab}_\rho(f))_{\rho \geq 0}$ as a measurement of some stochastic process.
- Allows using stochastic analysis to bypass the invariance principle.
- For the proof we introduce a new martingale embedding of μ as a re-normalized Brownian motion.

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Noise Stability - an Observation

If $f : \mathcal{C}_n \rightarrow \mathbb{R}$, we extend it harmonically to $f : [-1, 1]^n \rightarrow \mathbb{R}$. In particular, $T_\rho f(x) = f(\rho x)$. So, if $\mu_\rho = \text{Uniform}(\{-\sqrt{\rho}, \sqrt{\rho}\}^n)$,

$$\text{Stab}_\rho(f) = \mathbb{E}_\mu[f(x) \cdot f(\rho x)] = \mathbb{E}_\mu[f(\sqrt{\rho}x) \cdot f(\sqrt{\rho}x)] = \mathbb{E}_{\mu_\rho}[f^2].$$

Now, if ν is any measure on $[-1, 1]$, an orthogonal decomposition of $L^2(\mu)$ can be used to show

$$\text{Stab}_\nu(f) := \mathbb{E}_{\nu^{\otimes n}}[f^2] = \text{Stab}_{\text{Var}(\nu)}(f).$$

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A Re-normalized Brownian Motion

Consider the following martingale,

$$dX(t) = \sigma_t dB(t) \text{ with } \sigma_t = \text{diag}(\sqrt{(1 - X_i(t))(1 + X_i(t))}),$$

and define $\nu_t = \text{Law}(X_1(t))$.

Lemma

$$\text{Var}(\nu_t) = 1 - e^{-t}.$$

Proof.

$X_1(t)^2 = \text{martingale} + (1 - X_1(t)^2)dt$. So,

$$\frac{d}{dt} \mathbb{E} [X_1(t)^2] = 1 - \mathbb{E} [X_1(t)^2]. \text{ Now solve an ODE.} \quad \square$$

If $Y(t) \sim X(\infty)|X(t)$ then $\mathbb{E}[Y(t)] = X(t)$ and $\text{Cov}(Y(t)) = \sigma_t$.

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General Proof Strategy

Let $f : \mathcal{C}_n \rightarrow \{-1, 1\}$ and define the martingale $N_t = f(X(t))$.

Observe,

$$\begin{aligned}\mathbb{E} [[M]_t] &= \mathbb{E}[N_t^2] = \mathbb{E}_{\nu_t^{\otimes n}} [f^2] \\ &= \text{Stab}_{\text{Var}(\nu_t)}(f) = \text{Stab}_{1-e^{-t}}(f).\end{aligned}$$

The proof goes by finding a “model process” M_t to represent $\text{Stab}_\rho(\text{Maj})$ and a coupling which affords an almost-sure path-wise inequality,

$$[N]_t \leq [M]_t.$$

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Theorem

Among all Boolean functions, the dictator maximizes noise stability.

- Let $f : \mathcal{C}_n \rightarrow \{-1, 1\}$ and let $g : \mathcal{C}_n \rightarrow \{-1, 1\}$, $g(x) = x_1$.
- Define the martingales $N_t = f(X(t))$, $M_t = g(X(t)) = X_1(t)$.
- The theorem will follow, if we can find a coupling of N_t and M_t , such that $[N]_t \leq [M]_t$ almost surely.

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Interlude - Quadratic Variation

By Itô's formula

$$dM_t = \nabla g(X(t))\sigma_t dB_t = \sqrt{(1 - X_1(t))(1 + X_1(t))} dB_t.$$

Hence,

$$\frac{d}{dt}[M]_t = (1 - X_1(t))(1 + X_1(t)) = (1 - M_t^2).$$

In a similar way,

$$\frac{d}{dt}[N]_t = \|\nabla f(X(t))\sigma_t\|_2^2 = \sum_i (1 - X_i(t))(1 + X_i(t))\partial_i f(X(t)).$$

An application of Parseval's inequality gives,

$$\frac{d}{dt}[N]_t \leq (1 - f(X(t)))(1 + f(X(t))) = (1 - N_t^2).$$

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Interlude - a Coupling

By the Dambis-Dubins-Schwartz theorem, there exists a Brownian motion W_t , such that,

$$W_{[N]_t} = N_t \text{ and } W_{[M]_t} = M_t.$$

Reversing roles, for $\tau \geq 0$, write,

$$W_\tau = N_{T_1(\tau)} = M_{T_2(\tau)}.$$

So, keeping in mind that T_1 is the inverse function of $t \rightarrow [N]_t$

$$T_2'(\tau) = \frac{1}{1 - M_{T_2(\tau)}^2} = \frac{1}{1 - W_\tau^2} = \frac{1}{1 - N_{T_1(\tau)}^2} \leq T_1'(\tau).$$

Hence, almost surely, $T_2(\tau) \leq T_1(\tau) \implies [M]_\tau \geq [N]_\tau$.

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Interlude - Beyond the Toy Example

Let $\varphi : \mathbb{R} \rightarrow \mathbb{R}$ be any convex function and fix $t \geq 0$.

$$\begin{aligned}\mathbb{E}[\varphi(M_t)] &= \mathbb{E}[\varphi(W_{[M]_t})] = \mathbb{E}[\mathbb{E}[\varphi(W_{[M]_t}) | W_{[N]_t}]] \\ &\geq \mathbb{E}[\varphi(\mathbb{E}[W_{[M]_t} | W_{[N]_t}])] = \mathbb{E}[\varphi(W_{[N]_t})] \\ &= \mathbb{E}[\varphi(N_t)].\end{aligned}$$

Choose $\varphi(x) = x \log(x) + (1-x) \log(1-x)$, to get,

$$\mathbb{E}[\varphi(N_t)] = \mathbb{E}[\varphi(\mathbb{E}[f(X(\infty)) | X(t)])] = -\text{Ent}(f(X(\infty)) | X(t)).$$

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Interlude - Beyond the Toy Example

Define the mutual information $I(X; Y) := \text{Ent}(X) - \text{Ent}(X|Y)$.

Theorem (Most informative $X(t)$ bit)

Among all Boolean functions, the dictator maximizes the mutual information,

$$I(f(X(\infty)); X(t)).$$

Compare this with the 'most informative bit' conjecture of Courtade and Kumar.

Conjecture

Among all Boolean functions, the dictator maximizes the mutual information,

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where X and Y are ρ -correlated copies of uniform vectors on C_n .

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Most informative bit theorem

- Note that while $X(\infty)$ and $X(t)$ are correlated vectors, in general

$$(X(\infty), X(t)) \neq (X, Y),$$

for a ρ -correlated pair (X, Y) .

- Thus while the theorem is in the spirit of the Courtade-Kumar conjecture, it proves it with respect to a different noise model.
- Interestingly, the analog of the relation $\mathbb{E}[\varphi(M_t)] \geq \mathbb{E}[\varphi(N_t)]$, for general convex φ is known to be under for the 'usual' noise semi-group.

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Main ingredients:

- A martingale $N_t := f(X_t)$.
- A martingale M_t to represent noise stability of majority, or $\frac{2}{\pi} \arcsin(\rho)$.
- A differential *equality* for $[M]_t$.
- A differential *inequality* for $[N]_t$.

Constructing the Martingale M_t

There are infinitely many martingales M_t , which satisfy

$$\mathbb{E} [M_t] = \mathbb{E} [M_t^2] = \frac{2}{\pi} \arcsin(1 - e^{-t}) = \frac{2}{\pi} \arcsin(\rho).$$

We require one whose paths interact well with the paths of $f(X(t))$ when f has low influence.

One possibility is to take $M_t = f(\text{Maj}_n)$. However, that depends on the dimension.

Instead, we take a limiting object of $\text{Maj}_n(x) = \text{sign} \left(\frac{1}{\sqrt{n}} \sum x_i \right)$ in Gaussian space.

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Constructing the Martingale M_t

Let Φ stand for the Gaussian CDF and define the *Gaussian isoperimetric profile*:

$$I(x) := \Phi' \circ \Phi^{-1}(x).$$

Now, define M_t by, $dM_t = I(M_t)dB_t$.

It can be shown that $\mathbb{E}[[M]_t]$ encodes the limit of $\text{Stab}_\rho(\text{Maj}_n)$.
Evidently, we have the differential equality

$$\frac{d}{dt}[M]_t = I(M_t)^2$$

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For the proof, we use the representation,

$$\frac{d}{dt}[N]_t = \|\nabla f(X(t))\sigma_t\|_2^2 = \int_{t \geq \alpha} t d\nu(t),$$

where ν is a marginal of $\frac{X(\infty)|X(t)-X(t)}{\sigma_t}$ in direction ∇f .

Level 1 Inequality

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Thank You