

## Lecture 5 – Mon Feb 27

- Recall last time:
- $B_n = \pi_0 \text{Homeo}_+^+(\mathbb{R}^2, Q_n)$
  - this induces a right action of  $B_n$  on free group  $F_n$ ,
  - faithful:  $B_n \hookrightarrow \text{Aut}(F_n)$ .

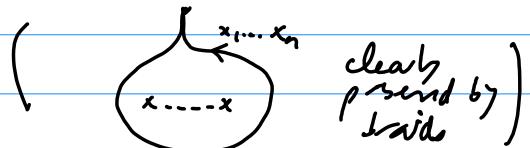
Now characterize the image of this map:

Thm (Arkin 1925):  $\beta \in \text{Aut}(F_n)$  is induced by an elt of  $B_n$  iff it satisfies:

- (1)  $\exists$  permutation  $\tau \in S_n$  &  $A_i \in F_n$  s.t.  $(x_i)\beta = A_i x_{\tau(i)} A_i^{-1} \quad \forall i=1..n.$
- (2)  $(x_1 \dots x_n)\beta = x_1 \dots x_n.$

PF:

- these 2 conditions are clearly necessary



&  $x_i$  is mapped to conjugates of each other

- conversely, we'll show that any autom. of  $F_n$  s.t. (1)+(2) can be expressed as product of  $(\sigma_i)_*$ .
- assume each word  $A_i x_{\tau(i)} A_i^{-1}$  is freely reduced,
- & consider the identity  $A_1 x_{\tau(1)} A_1^{-1} A_2 x_{\tau(2)} A_2^{-1} \dots A_n x_{\tau(n)} A_n^{-1} = x_1 \dots x_n \quad (*)$

Lemma 1: for  $(*)$  to hold, either  $\beta = \text{Id}$ , or  $\exists v \in \{1, \dots, n-1\}$  s.t. either

- $x_{\tau(v)} A_v^{-1}$  is absorbed by  $A_{v+1}$  (i.e.  $A_v x_{\tau(v)}^{-1}$  is a prefix in  $A_{v+1}$ )
- $A_v^{-1}$  absorbs  $A_{v+1} x_{\tau(v+1)}$  (i.e.  $A_{v+1} x_{\tau(v+1)}$  is a prefix in  $A_v$ ).

Assume lemma 1 holds: then reduce  $\beta$  to  $\text{Id}$  by induction on its "length"

Def.: length of  $\beta =$  sum of word lengths of  $(x_i)\beta = A_i x_{\tau(i)} A_i^{-1}, \quad i=1..n.$

Lemma 2: if (a) holds,  $\sigma_v \beta$  has length shorter than  $\beta$

if (b) holds,  $\bar{\sigma}_v \beta$  has length shorter than  $\beta$

[right action, so this means: first act by  $(\sigma_v^{\pm 1})_*$ , then by  $\beta$ ]

Then, lemma 1 + lemma 2  $\Rightarrow$  can reduce  $\beta$  to  $\text{Id}$  by composing repeatedly with action of  $\sigma_v$  or  $\bar{\sigma}_v$

$\Rightarrow \beta$  is indeed in the image of  $B_n \hookrightarrow \text{Aut}(F_n)$ .  $\square$

PF Lemma 2: assume (a) holds, then

$$(x_v)\beta = A_v x_{\tau(v)} A_v^{-1}$$

$$(x_{v+1})\beta = A_v x_{\tau(v)}^{-1} \tilde{A}_{v+1} x_{\tau(v+1)} \tilde{A}_{v+1}^{-1} x_{\tau(v)} A_v^{-1}$$

[this is a reduced word!]

so recalling that  $\sigma_v: x_v \mapsto x_v x_{v+1} x_v^{-1}$ , we have  
 $x_{v+1} \mapsto x_v$

$$(x_v)\sigma_v \beta = A_v \tilde{A}_{v+1} x_{\tau(v+1)} \tilde{A}_{v+1}^{-1} A_v^{-1} \rightarrow (\text{this might not be reduced, but is shorter than above anyway}).$$

→ total length decreased (since  $\sigma_v$  doesn't affect the other  $x_i$ ,  $i \neq v, v+1$ )!

$$\text{Similarly if (b) holds, } (x_v)\beta = A_{v+1} x_{\tau(v+1)} \tilde{A}_v x_{\tau(v)} \tilde{A}_v^{-1} x_{\tau(v+1)}^{-1} A_{v+1}^{-1}$$

$$(x_{v+1})\beta = A_{v+1} x_{\tau(v+1)} A_{v+1}^{-1}$$

[Reduced]

$$\rightarrow (x_v)\sigma_v^{-1}\beta = (x_{v+1})\beta = A_{v+1} x_{\tau(v+1)} A_{v+1}^{-1}$$

$$(x_{v+1})\sigma_v^{-1}\beta = (x_{v+1}^{-1} x_v x_{v+1})\beta = A_{v+1} \tilde{A}_v x_{\tau(v)} \tilde{A}_v^{-1} A_{v+1}^{-1} \leftarrow [\text{shorter}]$$

If Lemma 1: examine how LHS of (a) reduces to RHS:

$$A_1 x_{\tau(1)} A_1^{-1} A_2 x_{\tau(2)} A_2^{-1} \cdots A_n x_{\tau(n)} A_n^{-1} = x_1 \cdots x_n \quad (*)$$

first assume  $\exists v$  s.t. letter  $x_{\tau(v)}$  cancels in the reduction of LHS to RHS

& look at how it gets absorbed:

CNB: sequence of cancellations LHS  $\rightarrow$  RHS need not be unique, choose one & stick to it]

- can't be absorbed by a letter in  $A_v$  or  $A_v^{-1}$  since by assumption  $A_v x_{\tau(v)} A_v^{-1}$  is reduced.

- if absorbed by a letter in  $A_{v+1}$ , then (a) holds ( $x_{\tau(v)} A_v^{-1}$  absorbed by  $A_{v+1}$ )

- —————  $A_{v+1}^{-1}$ , then (b) holds ( $A_{v+1}^{-1}$  absorbs  $A_v x_{\tau(v)}$ )

- otherwise, it's absorbed by a letter to the left of  $x_{\tau(n-1)}$  or to the right of  $x_{\tau(n+1)}$

but then,  $x_{\tau(n-1)}$  or  $x_{\tau(n+1)}$  is absorbed by a letter that's closer

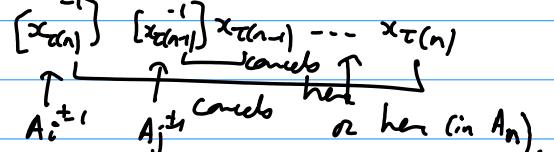
to it: i.e. if  $x_{\tau(n)}$  cancels w/ a letter

from  $A_i x_{\tau(i)} A_i^{-1}$  then  $x_{\tau(n+1)}$  cancels

w/ a letter from  $A_j x_{\tau(j)} A_j^{-1}$ ,

& either  $j=n$  ( $\Rightarrow$  done), or  $|j-(n+1)| < |i-n|$

so by induction on  $|i-n|$  reduce to previous cases & (a) or (b) holds!



Otherwise:  $\forall v$ ,  $x_{\tau(v)}$  survives the cancellations

$\Rightarrow$  then  $\tau(v)=v \quad \forall v$ , and the other letters all cancel!

however this implies  $A_{v-1}^{-1} A_v \mapsto 1 \quad \forall v$  ie  $A_{v-1} = A_v$

and  $A_1 \mapsto 1, A_n \mapsto 1$  ie  $A_1, A_n$  trivial  $\Rightarrow \text{Id}$

## Braid group of the sphere..

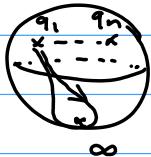
Theorem (Fadell - Van Buskirk 1962)

$B_n(S^2) = \pi_1 C_n(S^2)$  admits a presentation w/ generators  $\sigma_1, \dots, \sigma_{n-1}$ , & relations  $\sigma_i \sigma_j = \sigma_j \sigma_i$  if  $|i-j| \geq 2$   
 $\sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}$   
 $\sigma_1 \sigma_2 \dots \sigma_{n-1} \sigma_{n-1} \dots \sigma_2 \sigma_1 = 1$

Sketch PF:  $S^2 = \mathbb{R}^2 \cup \{\infty\}$ , so letting  $D_\infty = \{\{q_1 \dots q_n\} \in C_n(S^2), \{q_1 \dots q_n\} \ni \infty\}$   
 $D_\infty$  is a smooth submanif of  $C_n(S^2)$ , connected, of codim. 2.  
 $(^2 \{\infty\} \subset C_{n-1}(\mathbb{R}^2))$  and  $C_n(\mathbb{R}^2) = C_n(S^2) - D_\infty$ .

$\leftarrow$  Van Kampen's thm.  $\Rightarrow \pi_1 C_n(S^2) \cong \pi_1 C_n(\mathbb{R}^2) / \langle m \rangle$ , normal subgroup gen!  
 will see more carefully later  
 $(C_n(S^2) = C_n(\mathbb{R}^2) \cup V(D_\infty)$ , intersection =  $S^1$ -bundle/ $D_\infty$ ) by  $m$  = meridian b.p.

Meridian b.p. = keep  $q_2 \dots q_n$  fixed, move  $q_1$  around  $\infty$



=



$\equiv \sigma_1 \dots \sigma_{n-1} \sigma_{n-1} \dots \sigma_1$



(can also prove directly using a method similar to what we used for  $B_n(\mathbb{R}^2)$ : first understand pure braids  $P_n(S^2)$ ; for that use fibration  $\widetilde{C}_n(S^2) \rightarrow \widetilde{C}_{n-1}(S^2)$ , but beware the b.p.s. reduces to  $\pi_1$  only when  $n \geq 4$ .  $(\pi_2 \widetilde{C}_{n-1}(S^2) = 0$  only for  $n \geq 4$ )

## Braids and links:

- Work with PL links, i.e. union of disjoint polygonal simple closed curves in  $\mathbb{R}^3$ .
- Combinatorial equivalence of links ( $\leftrightarrow$  link isotopy): gen<sup>1</sup> by move =

replace an edge by edges provided that

- { - b, a, c not aligned in that order (nor a, c, b)
- { the convex hull  $\Delta(a, b, c)$  does not intersect the link anywhere outside the segment [a, c].

Call this move  $E_{a,c}^b$

(easy: this commutes w/ usual link isotopy)

- fix a line  $l \subset \mathbb{R}^3$  - call it axis.

Say the link  $L$  is in general position if none of its edges is coplanar w/ l.

(can always put a link in general position: if  $[ac]$  coplanar w/  $\ell$ , insert a vertex  $b$  very close to  $[ac]$  but not in the plane  $([ac], \ell)$ ).

• fix an orientation of  $L$ , & assume  $L$  is in general position:

view  $L$  in terms of its projection onto a plane  $\mathbb{R}^2 \perp \ell$ .

fix  $\textcirclearrowleft$  positive dir. of rotation around  $\ell$ .

Def: an oriented edge  $ab$  of  $L$  is positive if

its proj. to  $\mathbb{R}^2$  rotates around origin in + direction,

negative if rotates in - direction

$L$  is a closed braid if all its edges are positive

$h(L)$  ("height" of  $L$ ) := # neg edges



Ex. of closed braid ( $\Leftarrow$  no foil)

• Fact: any geometric braid  $\beta$  can be used to construct a closed braid  $\tilde{\beta}$

(identify initial & end pts of the strings) [& polygonal approx.]



Then (Alexander 1923):

$$\sigma_1^3 \in B_2$$

Every link is isotopic to a closed braid

If: get rid of negative edges by inserting sawtooths:

given a negative edge  $a_0 a$ :

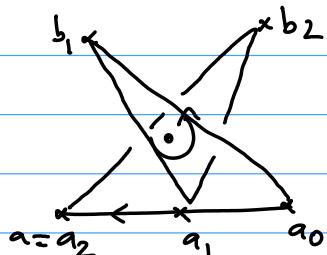
{ 1) subdivide it: take  $a_0, a_1, \dots, a_n = a$  aligned in that order

2) replace each negative edge  $a_{i-1} a_i$  by 2 positive edges  $a_{i-1} b_i, b_i a_i$

(st. this is an isotopy move  $\Sigma_{a_{i-1} a_i}^{b_i}$ , i.e. the triangle  $\Delta(a_{i-1}, a_i; b_i) \cap \bar{L} = [a_{i-1}, a_i]$ )

link obtained so far.

Ex:

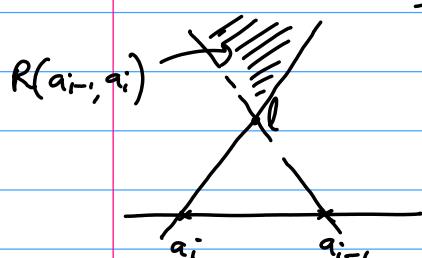


Lemma:  $[a_0, a]$  negative edge of  $L$  = link in general pos<sup>2</sup>.

$\Rightarrow$  can erect a sawtooth on  $[a_0, a]$ .

(= thm, since we can then inductively decrease  $h(L)$  by inserting sawtooths on negative edges).

PF-Lemma: Observe: given  $[a_{i-1}, a_i]$  negative edge and kn link  $l$ ,  
the 2 planes through  $l \& a_{i-1}$  /  $l \& a_i$  subdivide  $\mathbb{R}^3$  into 4 regions

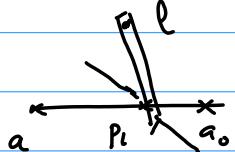


The edges  $[a_{i-1}, b_i]$  &  $[b_i, a_i]$  will be  $>0$  iff  
 $b_i$  is chosen in the region  $R(a_{i-1}, a_i)$

Case 1: if the projection of the edge  $[a_0, a]$  has no double pts (crossings).

→ choose  $b$  in the region  $R(a_0, a_1)$ , very far above (or below) the place of projection so plane  $(a_0, b, a)$  is almost  $\perp$  projection  
⇒ then the triangle  $\Delta(a_0, b, a)$  doesn't intersect  $L$   
(pass above or below any part of  $L$  that pass between  $[a_0, a]$  and  $L$ ).  
⇒ can end a sawtooth there (w/out subdividing  $[a_0, a]$ ).

Case 2: if  $\exists$  double pts:  $p_1, \dots, p_r \in [a_0, a]$  corresponding to crossings of the diagram

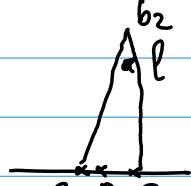


$P_1 :=$  plane through  $p_1$  and  $l$

Can draw a line  $\overset{\text{segment}}{\curvearrowright}$  in  $P_1$  joining  $p_2$  to  $l$  & avoiding  $L$  otherwise; Then make a very thin tooth which lies very close to this line.

(triangle  $\Delta(a_1, a_2, b_2)$  stays in nbd

of the chosen line segment & hence only intersects  $L$  along edge  $a_1 - a_2$ ).



proceed similarly near each  $p_i$ , & on remaining segments do as in Case 1.  $\square$

• Representing links by closed braids as a way to put more algebraic structure into things. As an illustration, can compute  $\pi_1$  of link complement easily for a closed braid:

Thm (Arkin 1925):  $\beta \in B_n$ ,  $\hat{\beta}$  = link obtained by closing  $\beta$   
 $\Rightarrow \pi_1(S^3 - \hat{\beta}) = \langle y_1, \dots, y_n \mid y_i = (y_i) \beta \quad \forall i = 1 \dots n \rangle$   
when  $(y_i) \beta = \text{action of } B_n \text{ on } F_n = \pi_1(R^2 - Q_n)$

Corollary: A group is a link gp iff it admits a presentation of the form  
 $\langle y_1, \dots, y_n \mid y_i = A_i(y_1 \dots y_n) y_{\tau(i)} A_i(y_1 \dots y_n)^{-1} \rangle$   
where  $\tau \in S_n$  permutation  
and  $A_i \in F_n$  s.t.  $\prod_{i=1}^n A_i y_{\tau(i)} A_i^{-1} = y_1 \dots y_n$  in  $F_n$ .

(using characterization of image of  $B_n \hookrightarrow \text{Aut}(F_n)$ )

Pf Thm: Let  $\varphi \in \text{Homeo}_+^+(\mathbb{R}^2, Q_n)$  agreeing  $\beta$  (supported in a disc  $D^2$ )  
& consider the 3mfld  $Y = (D^2 - Q_n) \times I/\sim$   $(z, 1) \sim (\varphi(z), 0)$   
- the mapping torus of  $\varphi|_{D^2 - Q_n}$

$Y$  is the complement of the link  $\hat{\beta}$  in the solid torus  $D^2 \times I/\sim = T$   
ie  $Y = T - \hat{\beta} \cong S^3 - (\hat{\beta} \cup \bar{l})$

unbr = closure of the axis



Clearly  $Y$  fibers over  $S^1$  with fiber  $D^2 - Q_n$

$$\Rightarrow \text{L.e.s. gives } 1 \rightarrow \pi_1(D^2 - Q_n) \xrightarrow{i_*} \pi_1(Y) \xrightarrow{p_*} \pi_1(S^1) \rightarrow 1$$

$$\qquad \qquad \qquad \pi_1(D^2 - Q_n)$$

Let  $y_i = i_{\pi_1(Y)}(x_i)$  ( $i=1 \dots n$ ), and  $t \in \pi_1(Y)$  loop gen, to  $\{z^3 \times I, z \in \partial D^2\}$ ,

( $t = \text{lift of the generator of } \pi_1(S^1)$ )

→ Then the exact sequence defines a presentation for

$\pi_1(Y) = \mathbb{Z} \times \pi_1(D^2 - Q_n)$  ( $\rightsquigarrow$  semidirect product of 2 free groups  $\Rightarrow$  just need to specify how to select t's to the beginning, by describing conjugation action of t on  $F_n$ )

generators =  $y_1, \dots, y_n, t$   
relations:  $ty_i t^{-1} = (y_i)\beta$

• Next,  $\pi_1(S^3 - \hat{\beta}) = \pi_1(Y)/\langle t \rangle$ , since  $t$  = meridian of  $\bar{l}$   
( $Y = (S^3 - \hat{\beta}) - \bar{l}\bar{l}$ ) on  $\partial(Y)$ .

(Van Kampen:  $(S^3 - \hat{\beta}) = Y \cup (S^3 - T)$  intersecting along torus  $\partial T$

so  $\pi_1(S^3 - \hat{\beta})$  generators  $y_1, \dots, y_n, t, l, L, M$  ( $M$  bounds disc  $\bar{l}, \bar{l}, L, \bar{L}$ )

rels:  $\begin{cases} ty_i t^{-1} = (y_i)\beta & (\text{from } Y) \\ L = l, M = 1 & (\text{from } \partial T \subset S^3 - T) \\ L = y_1 \dots y_n, M = t & (\text{from } \partial T \subset Y) \end{cases}$   $\Rightarrow \pi_1(S^3 - \hat{\beta}) = \pi_1(Y)/\langle t \rangle = \langle y_i \mid y_i = (y_i)\beta \rangle \checkmark$

gen of  $(S^3 - T)$  longitude & meridian on  $\partial T$  (axis) generate  $\pi_1(\partial T) = \mathbb{Z}^2$