

LECTURE 23 Wed 10/30

Arnold-Orlik-Solomon Alg. A_n

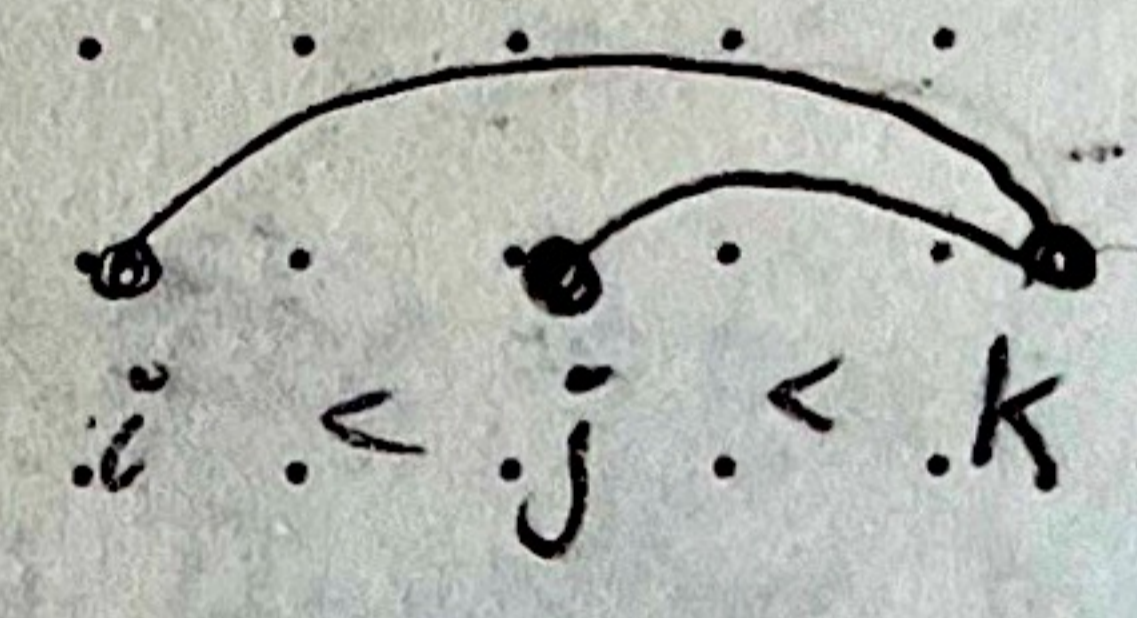
Anti-commutative generators $e_{ij} = -e_{ji} \quad i \neq j \in [n]$

relations: $e_{ij}e_{ik} - e_{ij}e_{jk} + e_{ik}e_{jk} = 0 \quad \forall i < j < k$

NBC basis: $e_F := \prod_{(i,j) \text{ edge of } F} e_{ij}$ (w.r.t. lex order) ← assume terms are in lex order.
 \forall increasing forests $F \subset K_n$.

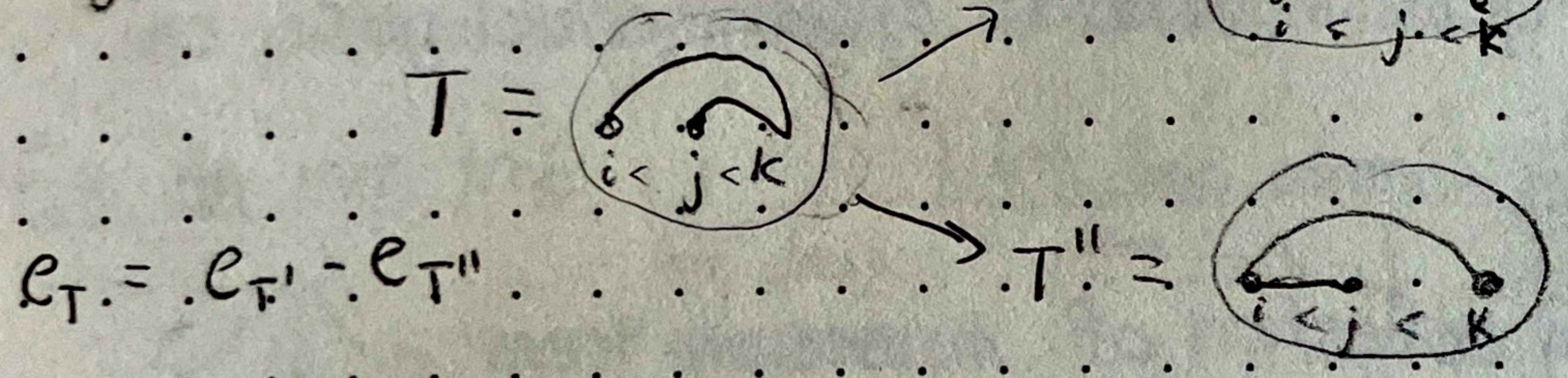
How to write any monomial in e_{ij} 's in the NBS-basis?

Lemma: F is increasing if it avoids the pattern.



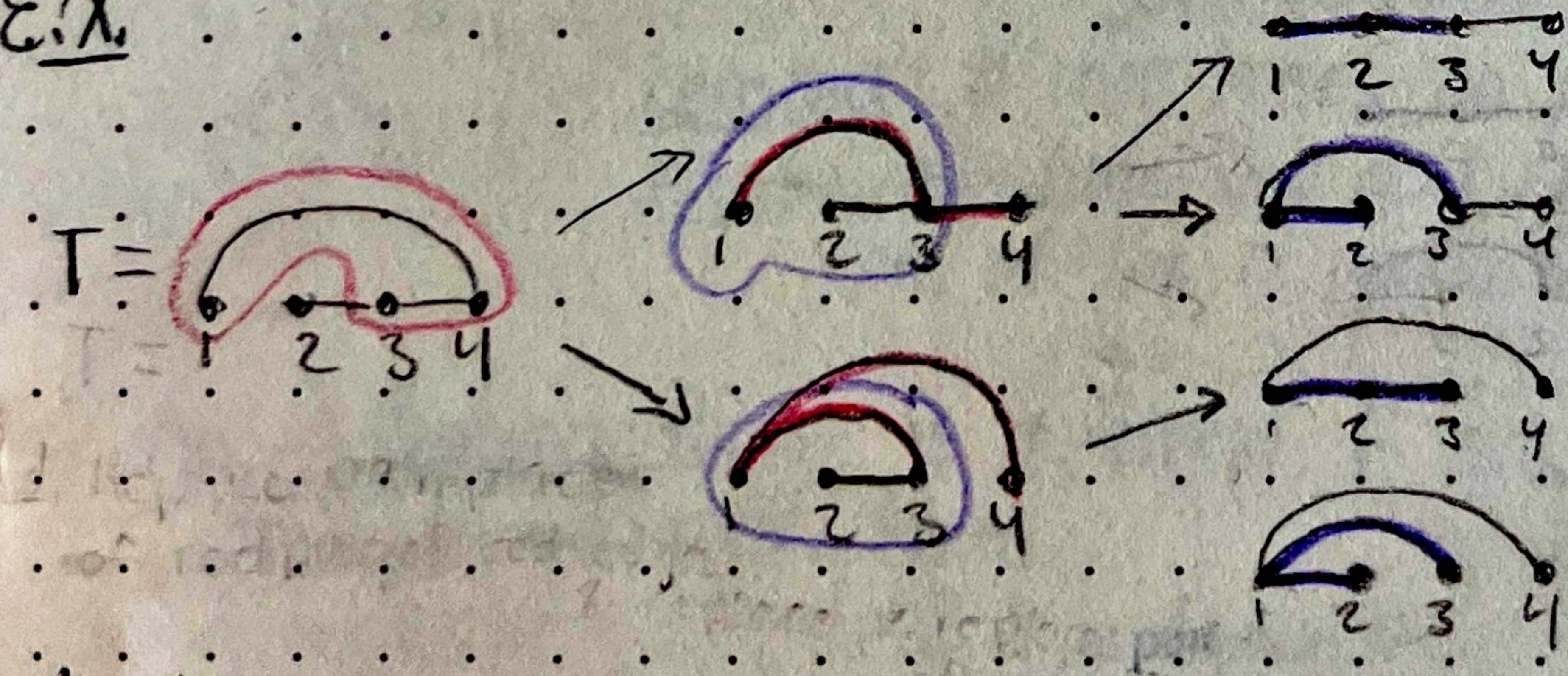
To write elts. in terms of basis, play

A game on trees:



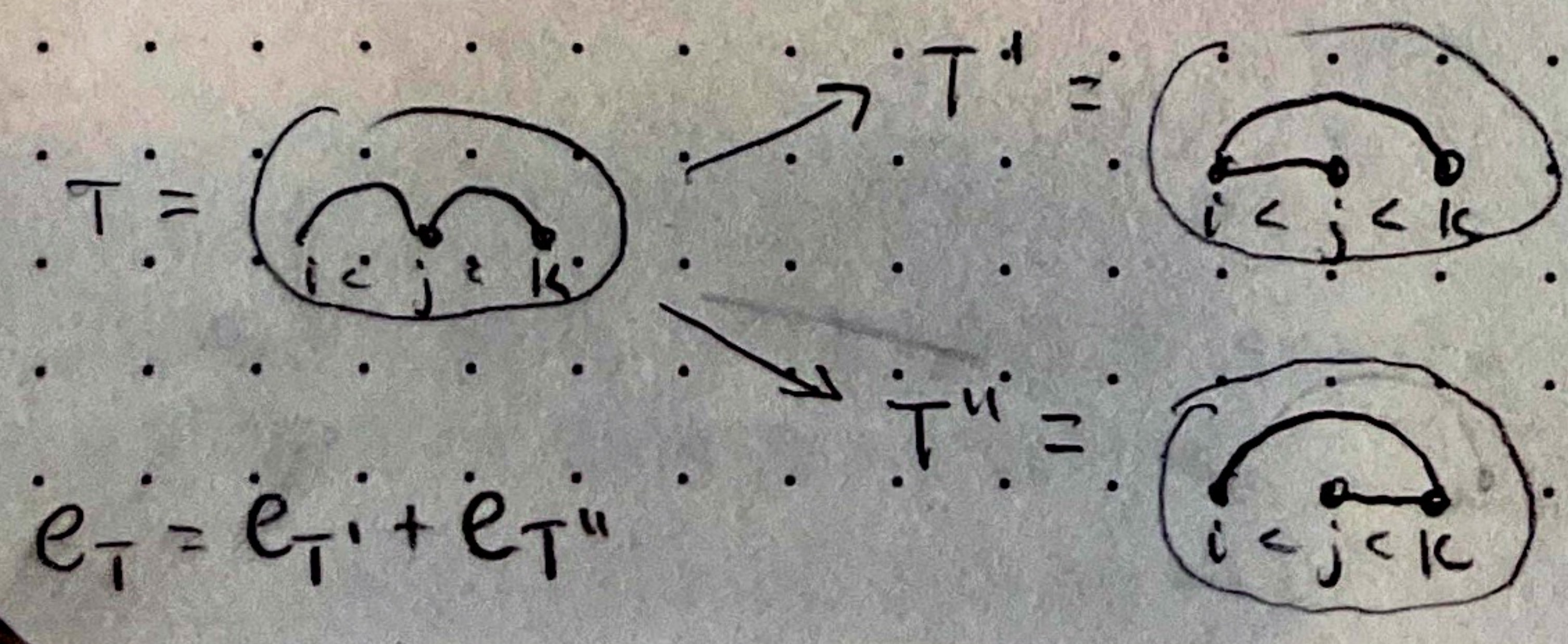
Game corresponds to replacing $e_{ik}e_{jk}$ with $e_{ije_{jk}} - e_{ij}e_{ik}$.

Ex.



Another game on trees:

Now replace $e_{ij}e_{jk} = e_{ij}e_{ik} + e_{ik}e_{jk} \quad i < j < k$

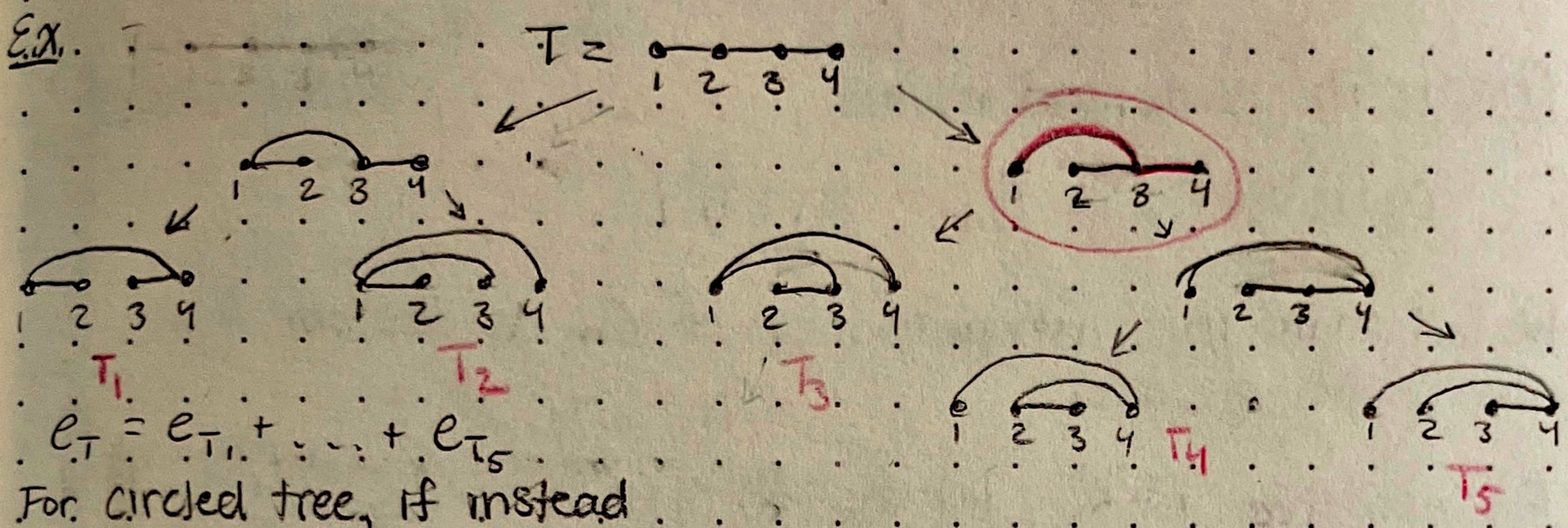


A certain set of alternating trees (forbidding pattern).

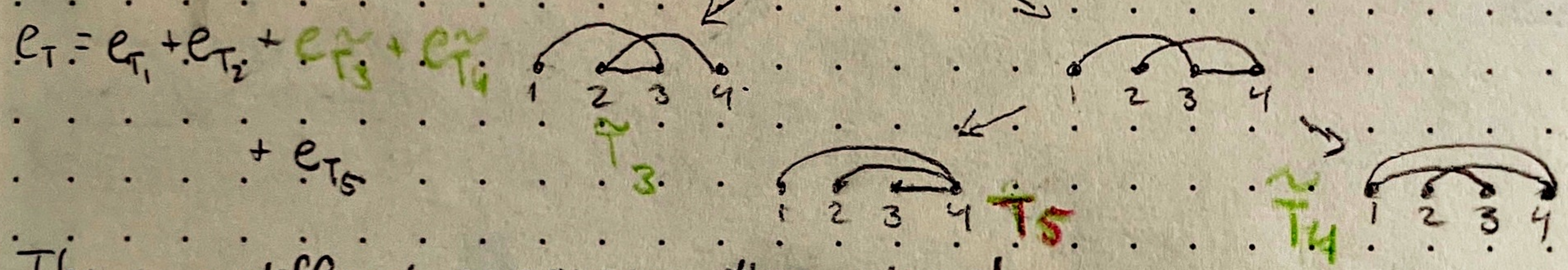
This game will end after finitely many steps, \Rightarrow

Prop: The set $\{e_T \mid T \text{ is an alternating forest}\}$ linearly spans A_n .

Q: Is this also a basis?



For circled tree, if instead we played game on edges



This is a different expansion than above!

\Rightarrow Alternating forests span, but are not a basis

Prop:

$$e_{\text{path}} = \sum_{T \text{ NCA} \subset K_n} e_T = \sum_{T \text{ NNA} \subset K_n} e_T$$

\uparrow non-crossing alternating trees \uparrow non-nesting alternating trees

terms = $\frac{1}{n+1} \binom{2n}{n} = \text{Catalan \# } C_n$

Can we also see geometrically why these #'s are the same?

Root polytopes

$\vec{e}_1, \dots, \vec{e}_n$ std coord vectors in \mathbb{R}^n

$$R_n = \text{conv}(\vec{0}, \vec{e}_i - \vec{e}_j \mid 1 \leq i < j \leq n)$$

Ex. $n=3$ $\vec{e}_1 - \vec{e}_2$ $\vec{e}_1 - \vec{e}_3$ (These are positive roots of type A rt. system)



Let $\text{Vol}(R_n) =$ the volume of its projection to \mathbb{R}^{n-1} under the map $(x_1, \dots, x_n) \mapsto (x_1, \dots, x_{n-1})$

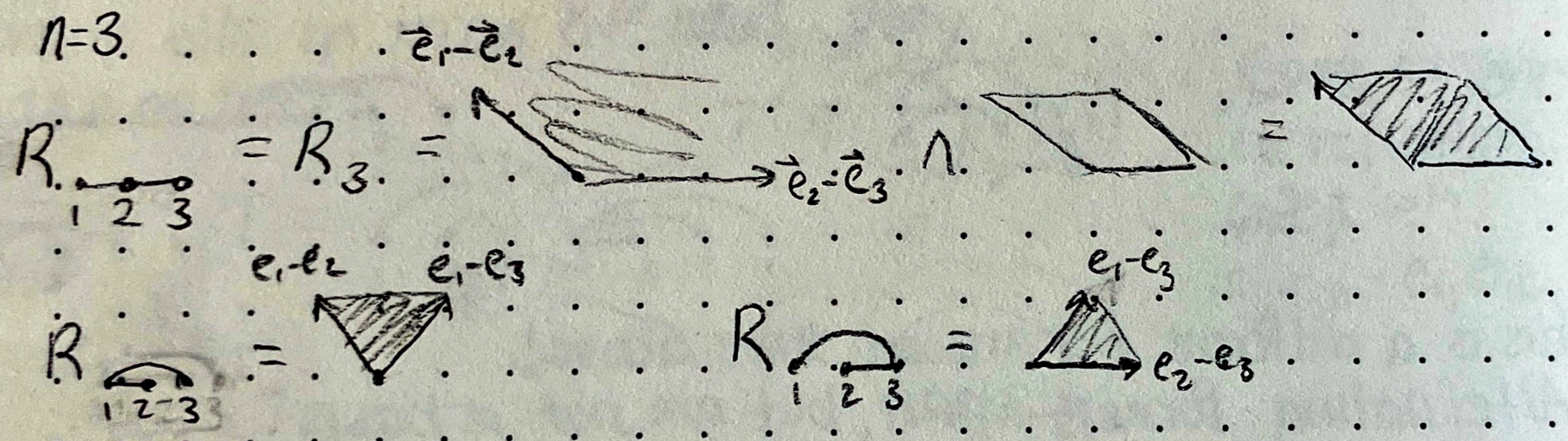
Thrm: The normalized volume

$$(n-1)! \text{Vol}(R_n) = C_n = \frac{1}{n+1} \binom{2n}{n}$$

We will relate this interpretation of C_n to the one from our game on trees

For any tree $T \subset K_n$, let $R_T = R_n \cap$ (the cone generated by vectors $\vec{e}_i - \vec{e}_j$ for edge (i,j) of T)

Ex. $n=3$



Lemma: If $T \begin{matrix} \nearrow T' \\ \searrow T'' \end{matrix}$ (in the second game) then

$$R_T = R_{T'} \cup R_{T''} \text{ and } R_{T'} \cap R_{T''} = \text{the common facet of } R_{T'} \text{ \& } R_{T''}$$

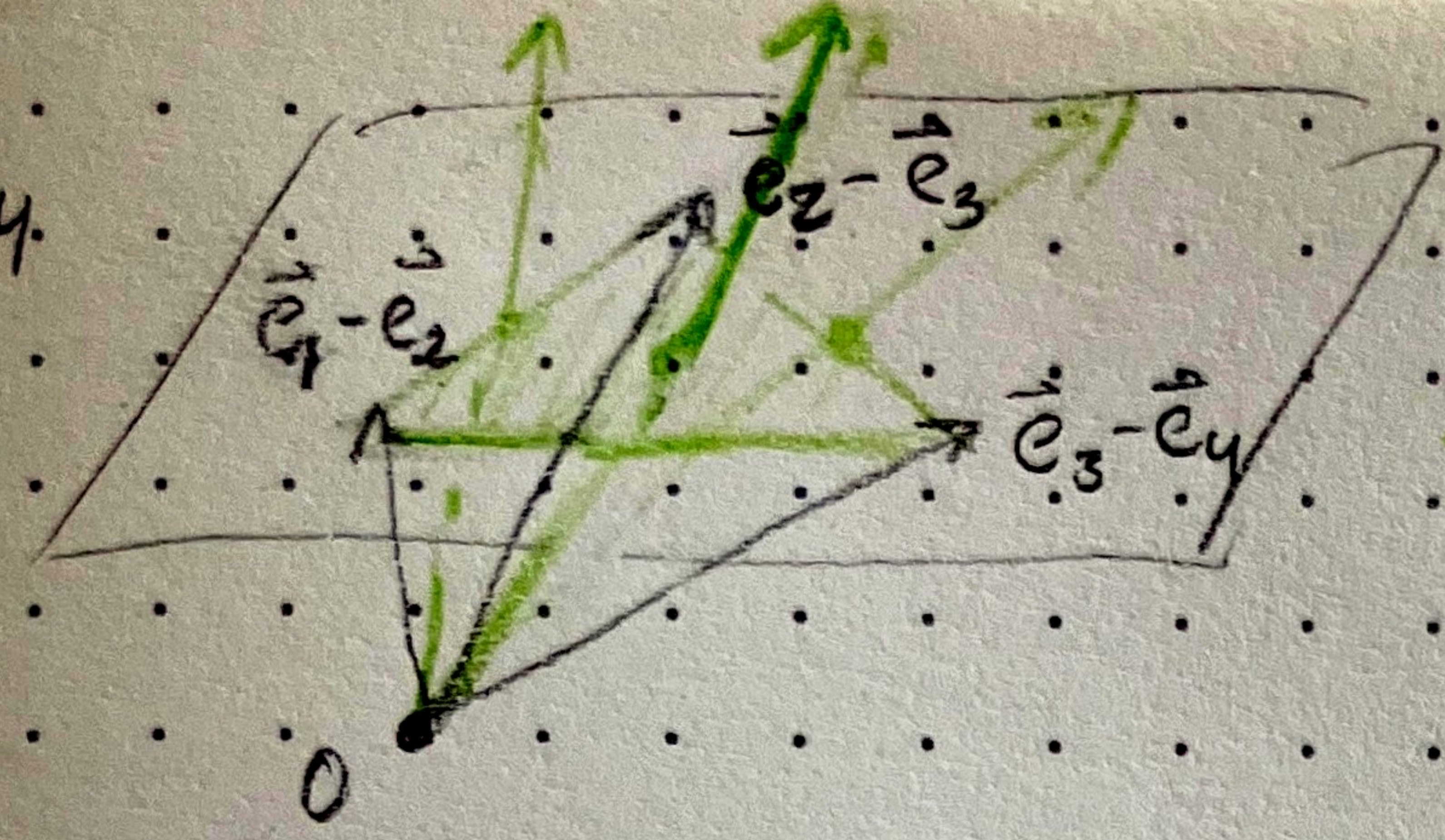
$$\text{Vol}(R_T) = \text{Vol}(R_{T'}) + \text{Vol}(R_{T''})$$

Lemma: If T is an alternating tree, then R_T is unit simplex (i.e. simplex of volume 1)

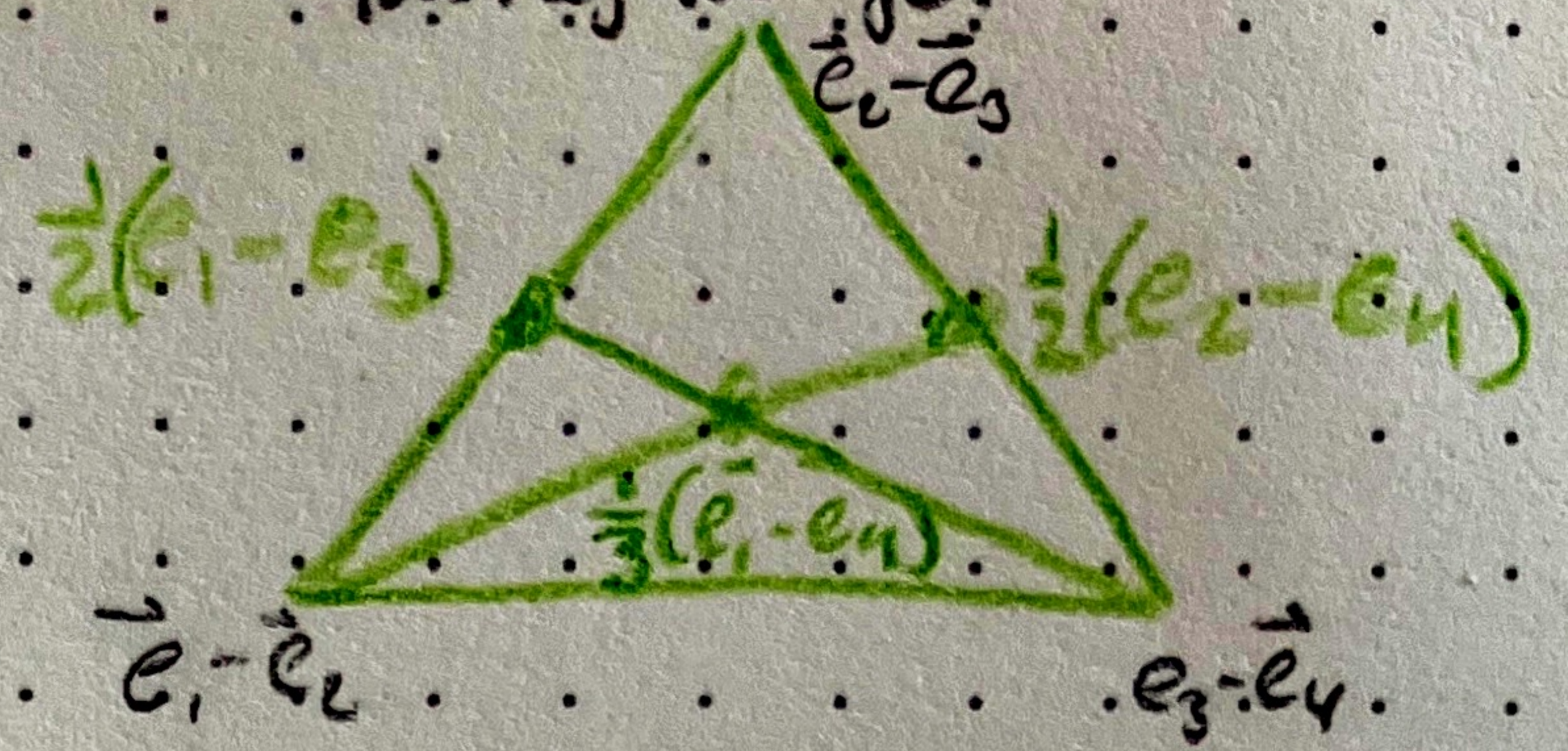
Cor: # endpoints in the game = $(n-1)! \text{Vol}(R_T) = C_n$

Ex R4

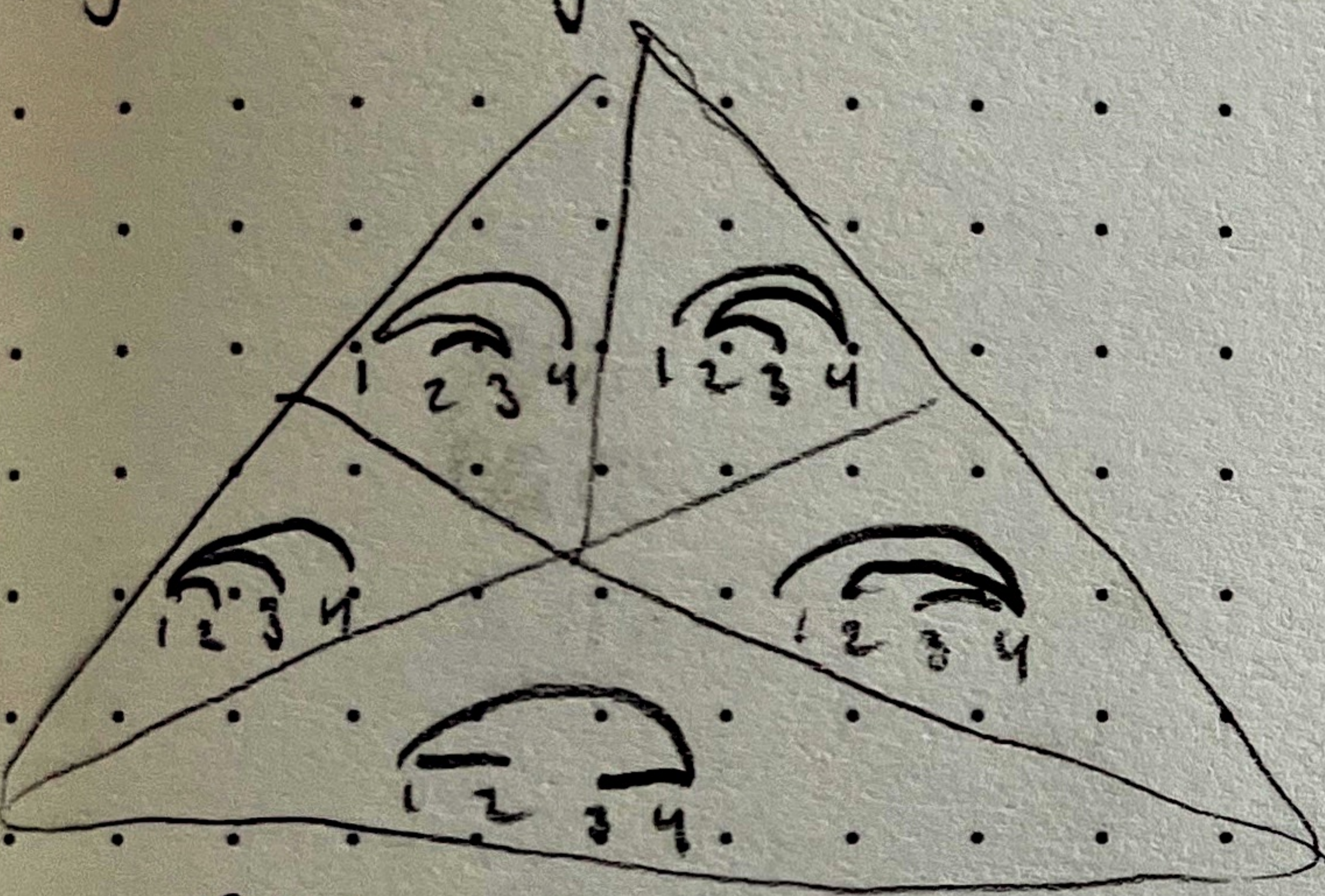
\mathbb{R}^3



Just looking at green triangle on plane, we get

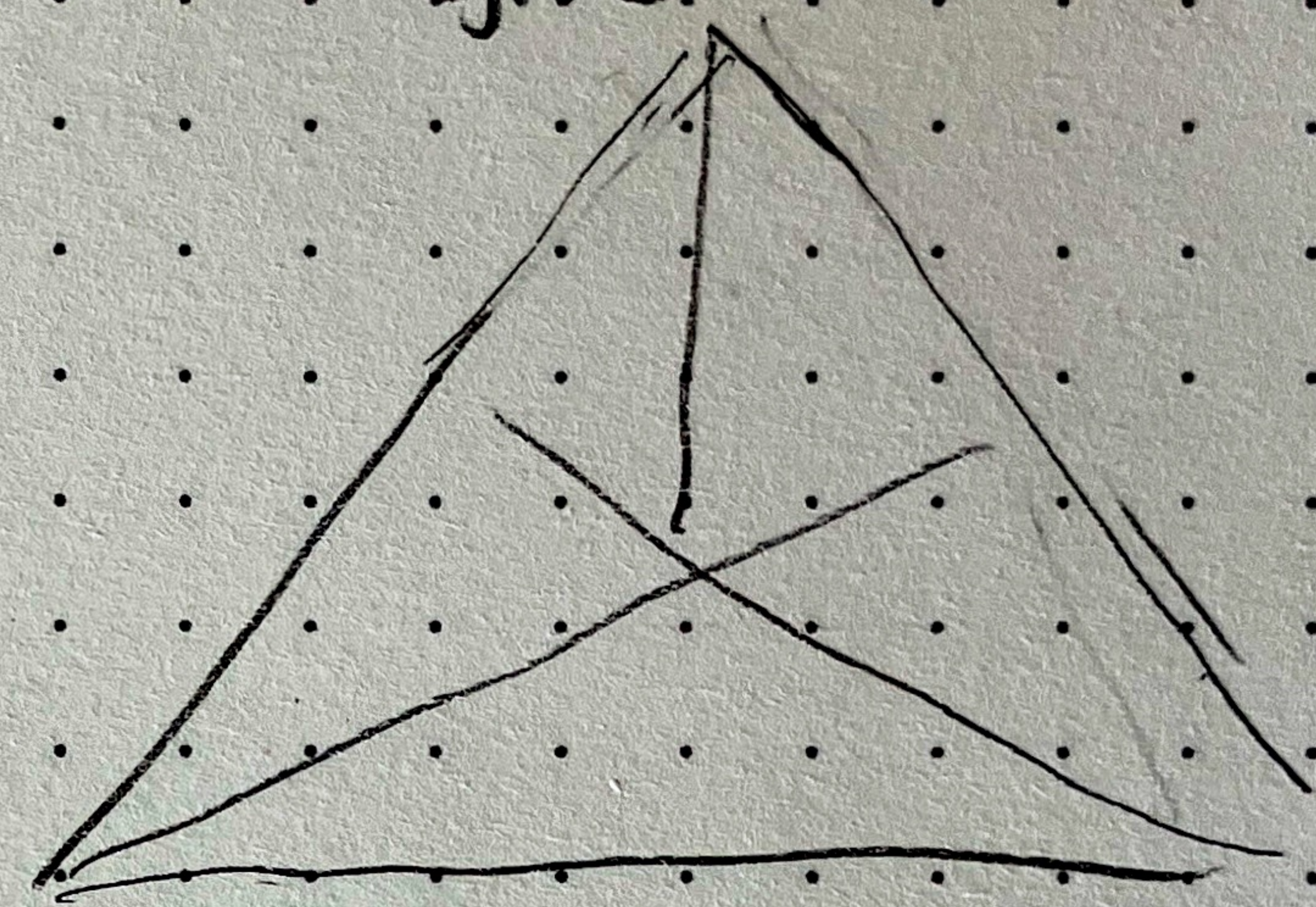


One way to play the game gives triangulation



non-crossing

Another way to play gives



non-nesting

somewhat if you look closely enough at it, it's related to the associahedron.

There might be something on this in the next pset.