

18.217 LECTURE 6

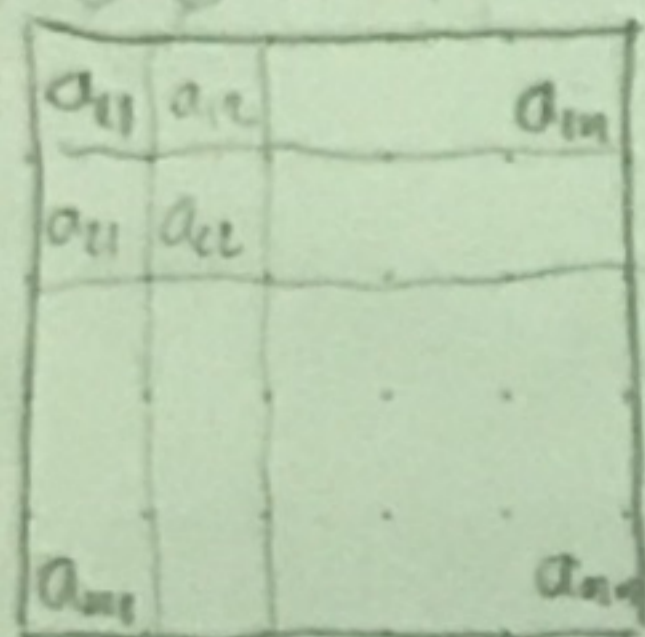
Last time:

$$RSK: \left\{ \begin{matrix} \boxed{a_{ij} \geq 0} \\ n \times n \end{matrix} \right\} \rightarrow \{(P, Q)\}$$

\updownarrow Self. Tsetlin patterns

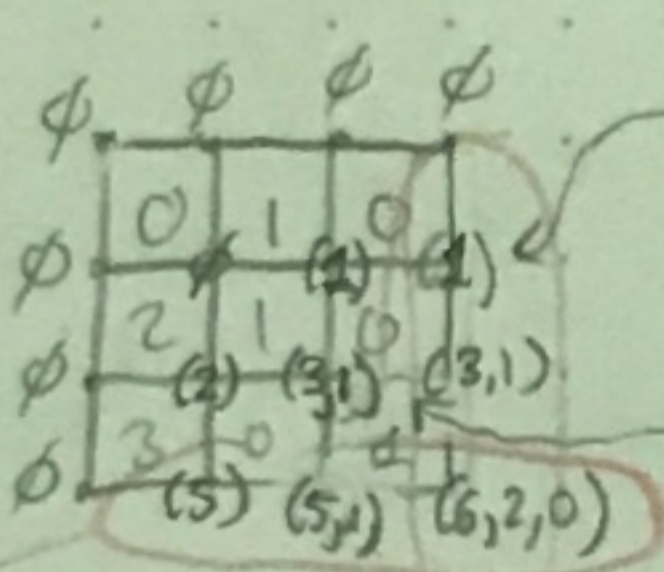
gen. growth diagrams

$$\left\{ \begin{matrix} \boxed{RPP'_s} \\ n \times n \end{matrix} \right\}$$



$$\begin{matrix} \lambda & & \nu \\ & a & \\ \mu & & \end{matrix} \quad \mathcal{R}(\lambda, \mu, \nu, a) \text{ (toggle rule)}$$

Ex. $A = \begin{bmatrix} 0 & 1 & 0 \\ 2 & 1 & 0 \\ 3 & 0 & 1 \end{bmatrix}$

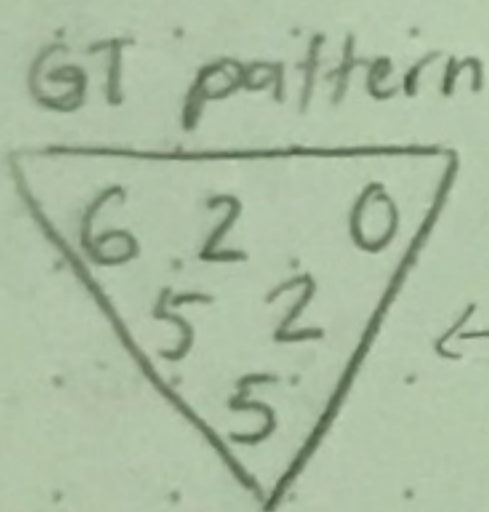


? Toggle $(0, 0, \dots)$, $(0, 0, \dots)$ and $(1, 0, \dots)$
 $\leftrightarrow (1, 0, \dots)$

? Toggle $(0, 0, \dots)$, $(1, 0, \dots)$, $(2, 0, \dots)$

$$\begin{matrix} 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{matrix}$$

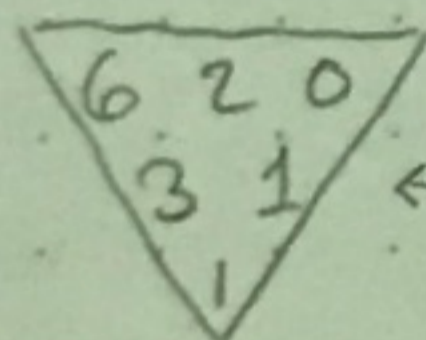
add 2 to largest part



$P =$

1	1	1	1	1	3
2	2				

Column GT pattern:



$Q =$

1	2	2	3	3	3
2	3				

$A \rightarrow (P, Q)$ of shape λ .

Thm Greene's Thm. Lets you find λ non-recursively

$A = (a_{ij}) \rightarrow$ multiset of pairs (j) repeated a_{ij} times.

"generalized perm" $(\begin{matrix} i_1 \leq i_2 \leq \dots \leq i_N \\ j_1 \leq j_2 \leq \dots \leq j_N \end{matrix})$

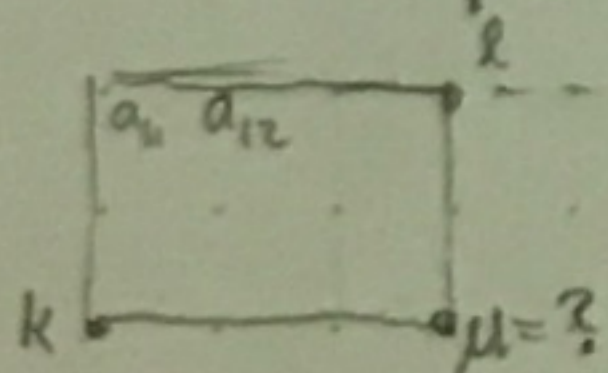
Then $\lambda_1 =$ size of maximal ^{weakly} increasing subseq. in j_1, \dots, j_n

$\lambda_1 + \lambda_2 =$ max cardinality subset of $\{1, \dots, N\}$ that can be covered by 2 weakly increasing subseq.

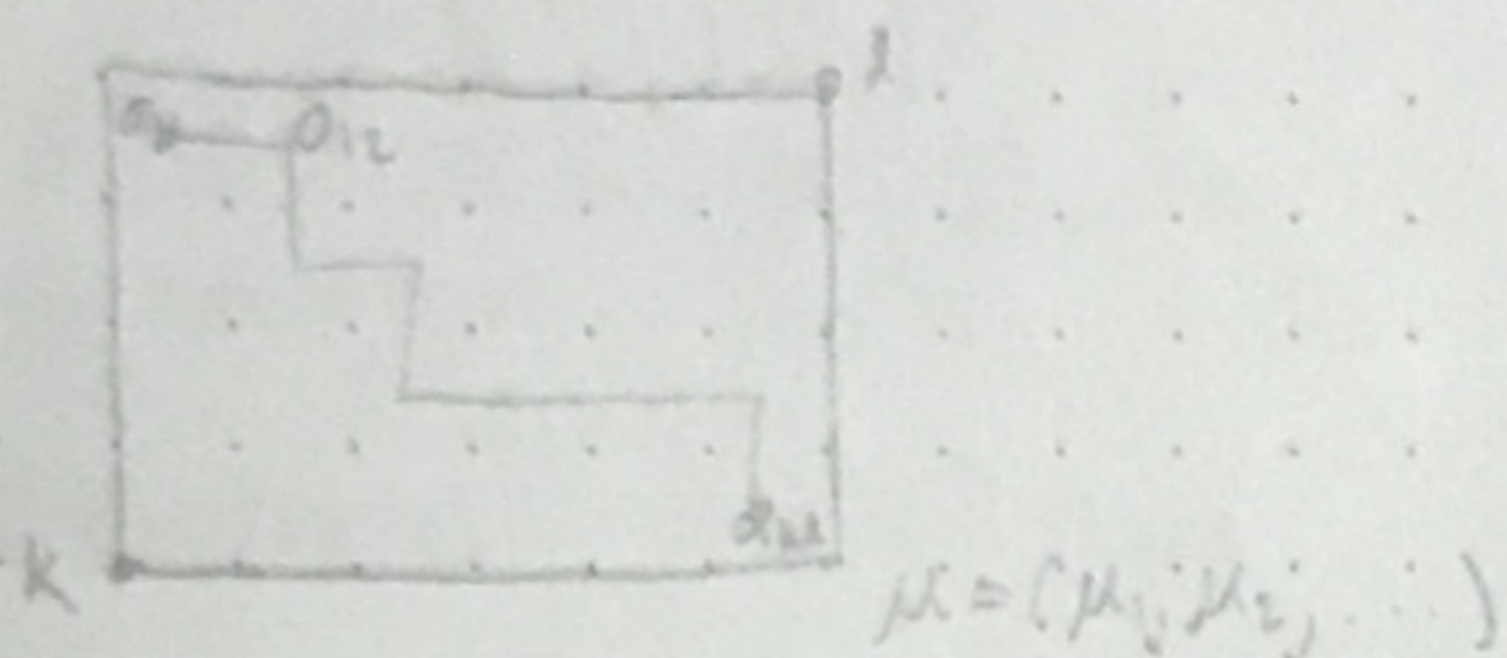
(Note: Finding max. seq. then finding longest incr. subseq. from what is left may not work)

$\lambda_1 + \dots + \lambda_k = \dots$ that can be covered with k weakly increasing subseq.

But how can we rephrase to see λ in terms of A (instead of permutation) non-recursively.



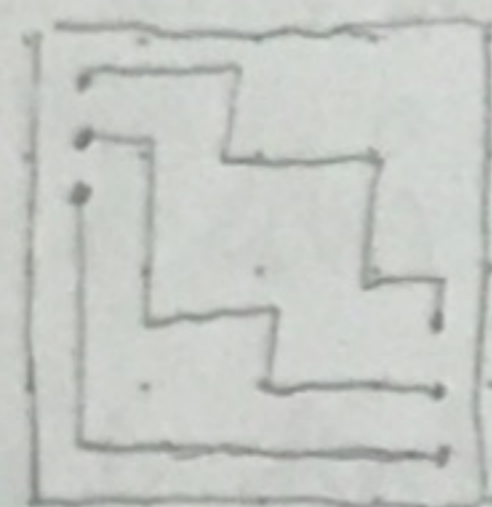
$$A^{k,l} = k \times l \text{ submatrix of } A$$



$\mu_\lambda = \max P$ lattice path from $(0,1)$ to (k,l)
 where value given by
 $\sum_{(i,j) \in P} a_{ij}$

Thm: $\mu = (\mu_1, \dots, \mu_r)$ $r = \min(k, l)$

where $\mu_1 + \mu_2 + \dots + \mu_s = \max S \subseteq [k] \times [l]$ covered by s non-crossing lattice paths connecting pts $(1,1), (2,1), \dots, (s,1)$ with $(k, l-s+1), \dots, (k, l)$



where max means largest value of $\sum_{(i,j) \in S} a_{ij}$

Exercise: Prove this.

Rational Algebras
 operations:

- " \cdot " times
- " $/$ " divide
- " $+$ " plus

Tropical semi-alg
 operations:

- " $+$ " plus
 - " $-$ " minus
 - "max" maximum
- tropicalize
 de tropicalize

$\min(a, b) := -\max(a, b)$

The above thm could be detropicalized to something about minors of determinants.
 but more on that some other time maybe

Schur Polynomials
 $S_\lambda(x_1, \dots, x_n)$

Schur functions
 $S_\lambda(x_1, x_2, \dots)$

← infinitely many terms

$\mathbb{Z}[x_1, \dots, x_n]^{S_n}$ = symmetric polys in n vars

def: 4 different formulas of S_λ :

1. Combinatorial definition

$S_\lambda(x_1, \dots, x_n) = \sum_{\substack{T \text{ a SSYT} \\ \text{w/ entries in } n}} X^{\text{weight}(T)}$

2. Classical definition (Weyl character formula)

$S_\lambda = \det(\dots) / \det(\dots)$

3. S_λ = certain Schubert polynomial

4. Coming later

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Schur Polynomials / functions

$$S_\lambda(x_1, \dots, x_n) \in \Delta_n = \mathbb{C}[x_1, \dots, x_n]^{S_n}$$

$$S_\lambda(x_1, x_2, \dots) \in \Delta := \lim_{n \rightarrow \infty} \Delta_n = \left\{ \text{set of formal power series in } x_1, x_2, \dots \text{ of bounded degree, invariant under perms of } x_1, x_2, \dots \right\}$$

Combinatorial def: $S_\lambda(x_1, \dots, x_n) = S_\lambda^{\text{comb}}(x_1, \dots, x_n) = \sum_{T \text{ SSYT of shape } \lambda, \text{ filled w/ entries in } \{1, \dots, n\}} X^{\text{weight}(T)}$

where $X^{\text{weight}(T)} = \prod_{i=1}^n x_i^{(\# \text{ of } i\text{'s in } T)}$

Ex. $S_0(x_1, \dots, x_n) = x_1 + \dots + x_n$

(RSK) \Rightarrow Cauchy formula

Cauchy formula: $\sum_{\lambda \text{ perm w/ at most } \min(m,n) \text{ parts}} S_\lambda(x_1, \dots, x_m) S_\lambda(y_1, \dots, y_n) = \prod_{i=1}^m \prod_{j=1}^n \frac{1}{1 - x_i y_j}$

LHS = $\sum_{\text{pairs } P, Q \text{ of SSYT}} X^{\text{weight}(P)} Y^{\text{weight}(Q)}$

RHS = $\prod_{i,j} \sum_{a_{ij} \geq 0} (x_i y_j)^{a_{ij}} = \sum_{\substack{m \times n \text{ matrix} \\ A = (a_{ij})}} X^{\text{columns of } A} Y^{\text{rows of } A}$

And then these are equal by RSK

Classical def of S_λ :

$$\alpha = (\alpha_1, \dots, \alpha_n) \in \mathbb{Z}_{\geq 0}^n, \quad a_\alpha := \det \begin{pmatrix} x_1^{\alpha_1} & \dots & x_n^{\alpha_1} \\ \vdots & & \vdots \\ x_1^{\alpha_n} & \dots & x_n^{\alpha_n} \end{pmatrix}$$

WLOG assume that $\alpha = (\alpha_1 > \alpha_2 > \dots > \alpha_n)$

called δ in Lie Theory $\rightarrow \delta = (n-1, n-2, \dots, 1, 0)$

$$a = \lambda + \delta, \quad \lambda = (\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n)$$

$$a_\delta = \prod_{1 \leq i < j \leq n} (x_i - x_j) \quad (\text{Vandermonde})$$

Def: $S_\lambda^{\text{class}}(x_1, \dots, x_n) = \frac{a_{\lambda+\delta}}{a_\delta}$ (a symmetric poly since numerator & denominator both antisymmetric)

Ex. $n=2, \lambda = (1, 0) = \square, \lambda + \delta = (2, 0)$

$$S_\lambda^{\text{class}}(x_1, x_2) = \frac{\begin{vmatrix} x_1^2 & x_2^2 \\ 1 & 1 \end{vmatrix}}{\begin{vmatrix} x_1 & x_2 \\ 1 & 1 \end{vmatrix}} = \frac{x_1^2 - x_2^2}{x_1 - x_2} = x_1 + x_2$$

Thm: $S_\lambda^{\text{comb}} = S_\lambda^{\text{class}}$

Symmetric group S_n

generators: s_1, s_2, \dots, s_n

$s_i = (i, i+1)$ - adjacent transpositions

Relations: 1.) $s_i^2 = \text{id}$

2.) $s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1}$

3.) $s_i s_j = s_j s_i$ when $|i-j| \geq 2$

Can write permutations $w = s_{i_1} s_{i_2} \dots s_{i_\ell}$

def: This is a reduced decomposition of w if ℓ is as small as possible
 $\ell = \ell(w)$ is the length of w

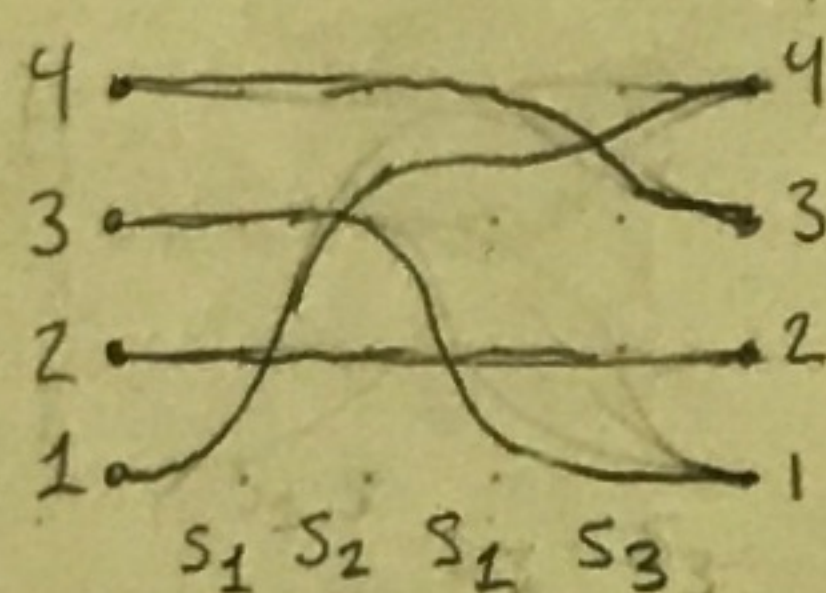
Lemma 1: $\ell(w) = \#$ of inversions in $w := \# \{ (i, j) \mid 1 \leq i < j \leq n, w(i) > w(j) \}$

Lemma 2: Any 2 reduced decompositions of w can be obtained from each other by a sequence of moves (2) and (3).

Note: This is true in the more general setting of Coxeter groups as well

Wiring Diagrams

Ex. $w = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 4 & 2 & 1 & 3 \end{pmatrix}$



$w = s_3 s_1 s_2 s_1$

decomposition is reduced iff its wiring diagram has no double crossings

Can see lemma 1 & 2 from wiring diagram

Def: Divided difference operators are operators

$\partial_1, \partial_2, \dots, \partial_{n-1}$ acting on $\mathbb{C}[x_1, \dots, x_n]$ by

$$\partial_i: f(x_1, \dots, x_n) \mapsto \frac{1}{x_i - x_{i+1}} (1 - s_i)(f) = \frac{f(x_1, \dots, x_i, x_{i+1}, \dots, x_n) - f(x_1, \dots, x_{i+1}, x_i, \dots, x_n)}{x_i - x_{i+1}}$$

Lemma: Operators satisfy relations

(1) $\partial_i^2 = 0$

(2) $\partial_i \partial_{i+1} \partial_i = \partial_{i+1} \partial_i \partial_{i+1}$

(3) $\partial_i \partial_j = \partial_j \partial_i$ when $|i-j| \geq 2$

Notation: For $w = s_{i_1} \dots s_{i_\ell}$ reduced decomp, $\partial_w := \partial_{i_1} \dots \partial_{i_\ell}$ (depends only on w)

$w_0 := \begin{pmatrix} 1 & 2 & \dots & n \\ n & n-1 & \dots & 1 \end{pmatrix}$, the longest perm in S_n

Schur as Schubert Polynomials

Def: $S_\lambda^{\text{schub}} := \partial_{w_0} (x^{\lambda+\delta})$ $\lambda = (\lambda_1, \dots, \lambda_n)$, $\delta = (n-1, n-2, \dots, 0)$

Thrm: $S_\lambda^{\text{schub}}(x_1, \dots, x_n) = S_\lambda^{\text{class}}(x_1, \dots, x_n)$

$$\text{Thrm: } \partial_{w_0} = \frac{1}{\prod_{1 \leq i < j \leq n} (x_i - x_j)} \left(\sum_{w \in S_n} (-1)^{l(w)} w \right)$$

$$\text{Ex. } n=3 \\ \partial_1 \partial_2 \partial_1 = \frac{1}{x_1 - x_2} (1 - s_1) \frac{1}{x_2 - x_3} (1 - s_3) \frac{1}{x_1 - x_2} (1 - s_1)$$

$$= \frac{1}{(x_1 - x_2)(x_2 - x_3)(x_1 - x_2)} (1 - s_1 - s_2 + s_1 s_2 + s_2 s_1 - s_1 s_2 s_1)$$

$$\partial_{w_0} (x^{\lambda + \delta}) = \frac{1}{\text{Vand}} \cdot \underbrace{\sum_{w \in S_n} (-1)^{l(w)} w}_{a_{\lambda + \delta}} (x^{\lambda + \delta})$$