

last time:

- λ/μ skew Young diagram, $|\lambda/\mu| = n$
 $\beta = (\beta_1, \dots, \beta_k)$ composition of n

Young's orthogonal form \Rightarrow

$$\chi_{\lambda/\mu}(\beta) = \sum_{T \in \text{SYT}(\lambda/\mu)} \prod_{i \in [n] \setminus \{\beta_1, \beta_1 + \beta_2, \dots, \beta_1 + \dots + \beta_k\}} \frac{1}{c_{i+1} - c_i}$$

the value of character of $V_{\lambda/\mu}$ on a permutation with cyclic type β

$T \in \text{SYT}(\lambda/\mu)$ content vector of T $c(T) = (c_1, \dots, c_n)$

- Murnaghan - Nakayama Rule:

$$\chi_{\lambda/\mu}(\beta) = \sum_{RT \text{ ribbon tableau of shape } \lambda/\mu \text{ and type } \beta} (-1)^{\text{ht}(RT)}$$

Moreover, we gave a multivariate generalization of the M.-N. rule

where we replaced c_i by $-x_i$, where x_i 's are variables:

$$\chi_{\lambda/\mu}^{(x_i)} = \sum_{RT \text{ ribbon tableau of shape } \lambda/\mu \text{ and type } \beta} \text{wt}(RT)$$

$$\text{wt}(RT) := \prod_{\text{ribbon in } RT} \text{wt}(\text{ribbon})$$

$$\text{wt} \left(\begin{array}{c} \text{skew Young diagram} \\ \text{with ribbons} \end{array} \right) := (-1)^{\#} \prod_{i \in \{c+1, \dots, c'-1\}} \frac{1}{x_i - x_{i+1}}$$

This M.-N. rule basically reduces to 3 lemmas about certain rational expressions in x_i 's.

For $i_1, \dots, i_n \in \mathbb{Z}$, let

$$\langle i_1, \dots, i_n \rangle := \frac{1}{x_{i_1} - x_{i_2}} \frac{1}{x_{i_2} - x_{i_3}} \dots \frac{1}{x_{i_{n-1}} - x_{i_n}}$$

Lemma 1. A, B ^{non-empty} disjoint sequences of int.

$$\sum_{C \in \text{Shuffle}(A, B)} \langle C \rangle = 0$$

all shuffles of $A \cup B$

Lemma 2. $A, B, \{1\}$ ^{non-empty} disjoint

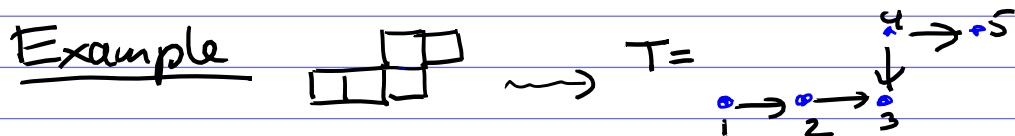
$$\sum_{C \in \text{Shuffle}(A, B)} \langle 1, C, 1 \rangle = 0$$

Lemma 3. (Tree relation)

Let T be any tree on vertices $\{1, \dots, n\}$ with directed edges.

$$\sum_{\substack{\sigma = \sigma_1 \dots \sigma_n \text{ permutation of } [n] \\ \sigma^{-1}(i) < \sigma^{-1}(j) \text{ for} \\ \text{any edge } i \rightarrow j \text{ in } T}} \langle \sigma_1 \dots \sigma_n \rangle = \prod_{\substack{i \rightarrow j \\ \text{edge in } T}} \langle ij \rangle$$

(For the M.N rule we need the "ribbon relation" which is a special case of Lemma 3 when T is a chain.)



$\sigma = (\sigma_1 \dots \sigma_n)$ correspond to

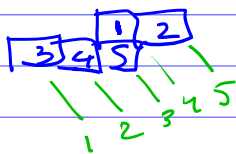
SYT's of ribbon shape

(i.e. linear extensions of the

poset given by directed tree)

where σ = content vector of tableau.

$$\langle 4, 5, 1, 2, 3 \rangle + \langle 4, 1, 5, 2, 3 \rangle + \dots$$



$$= \langle 12 \rangle \langle 23 \rangle \langle 43 \rangle \langle 45 \rangle = (-1) \langle 1, 2, 3, 4, 5 \rangle$$

\nearrow
wt(ribbon).

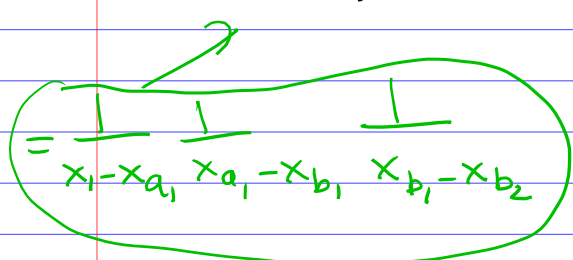
We'll prove these lemmas by induction. We need another lemma...

Lemma 4 $A, B, \{i\}$ disjoint

$$\sum_{C \in \text{Shuffle}(A, B)} \langle 1, C \rangle = \langle 1, A \rangle \langle 1, B \rangle$$

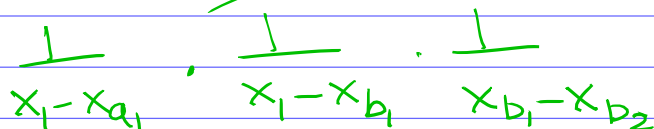
Example $A = (a_1)$ $B = (b_1, b_2)$

$$\langle 1, a_1, b_1, b_2 \rangle + \langle 1, b_1, a_1, b_2 \rangle$$



$$+ \langle 1, b_1, b_2, a_1 \rangle$$

$$= \langle 1, a_1 \rangle \langle 1, b_1, b_2 \rangle$$



Proof of Lemma 4, Induction on $|A| + |B|$.

Base: $A = B = \emptyset$

$$\text{LHS} = \langle 1 \rangle = 1, \quad \text{RHS} = \langle 1 \rangle \langle 1 \rangle = 1$$

(by convention, $\text{Shuffle}(\emptyset, \emptyset)$

consists of 1 element \emptyset ;

and $\prod_{\text{empty set}} = 1$)

Induction Step: Let

$$A = (a_1, A') \quad B = (b_1, B')$$

$$\sum_{C \in \text{Shuffle}(A, B)} \langle 1, C \rangle = \begin{array}{l} \text{any shuffle} \\ \text{of } A \cup B \\ \text{starts with} \\ a_1 \text{ or } b_1 \end{array}$$

$$= \sum_{C' \in \text{Shuffle}(A', B)} \langle 1, a_1 \rangle \langle a_1, C' \rangle +$$

$$+ \sum_{C'' \in \text{Shuffle}(A, B')} \langle 1, b_1 \rangle \langle b_1, C'' \rangle$$

ind. hypothesis

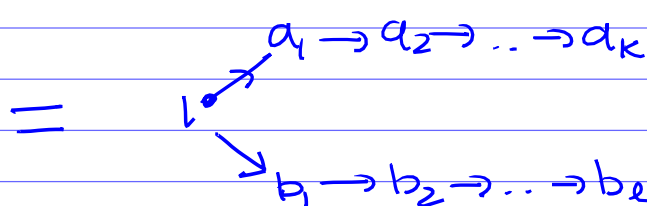
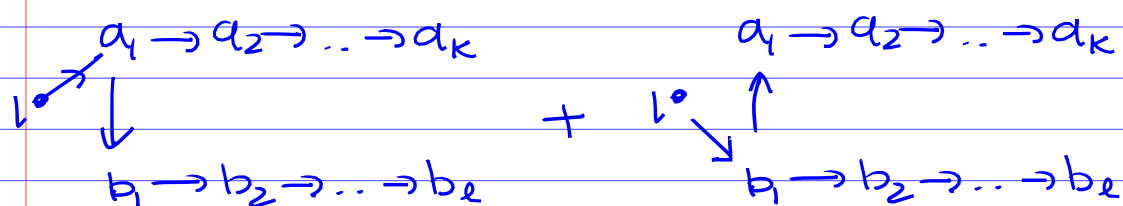
$$= \langle 1, a_1 \rangle \langle a_1, A' \rangle \langle a_1, B \rangle +$$

$$+ \langle 1, b_1 \rangle \langle b_1, A \rangle \langle b_1, B' \rangle$$

$$= (\langle 1, a_1, b_1 \rangle + \langle 1, b_1, a_1 \rangle) \langle A \rangle \langle B \rangle$$

$$= \langle 1, a_1 \rangle \langle 1, b_1 \rangle \langle A \rangle \langle B \rangle$$

$$= \langle 1, A \rangle \langle 1, B \rangle, \text{ as needed}$$



Here an edge $i \rightarrow j$ represents a term $\langle i, j \rangle := \frac{1}{x_i - x_j}$

□

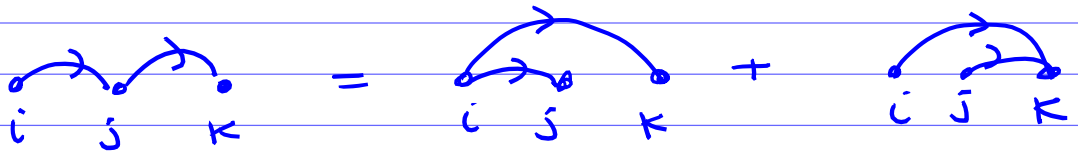
Key Relation :

follows from

$$x_i - x_k = (x_i - x_j) + (x_j - x_k).$$

$$\frac{1}{x_i - x_j} \cdot \frac{1}{x_j - x_k} = \frac{1}{x_i - x_k} \cdot \frac{1}{x_i - x_j} + \frac{1}{x_j - x_k} \cdot \frac{1}{x_i - x_k}$$

$$\langle ij \rangle \cdot \langle jk \rangle = \langle ik \rangle \langle ij \rangle + \langle jk \rangle \langle ik \rangle$$



Proof of Lemma 1

$$A = (a, A')$$

$$B = (b, B')$$

$$\sum_{C \in \text{shuffle}(A, B)} \langle C \rangle$$

$$= \sum_{C' \in \text{shuffle}(A', B)} \langle a, C' \rangle + \sum_{C'' \in \text{shuffle}(A, B')} \langle b, C'' \rangle$$

by lemma 4

$$= \langle a, A' \rangle \langle a, B \rangle + \langle b, A \rangle \langle b, B' \rangle$$

$$= (\langle a, b \rangle + \langle b, a \rangle) \langle A \rangle \cdot \langle B \rangle = 0$$

$$\begin{array}{c}
 a_1 \rightarrow a_2 \rightarrow \dots \rightarrow a_k \\
 \downarrow \\
 b_1 \rightarrow b_2 \rightarrow \dots \rightarrow b_k
 \end{array}
 +
 \begin{array}{c}
 a_1 \rightarrow a_2 \rightarrow \dots \rightarrow a_k \\
 \uparrow \\
 b_1 \rightarrow b_2 \rightarrow \dots \rightarrow b_k
 \end{array}
 = 0$$

□

Proof of Lemma 3 (tree relation)

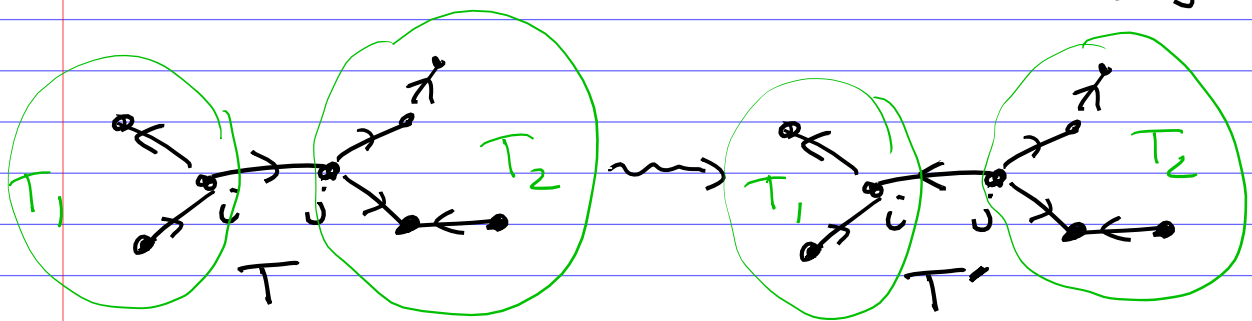
Induction on n (# vertices in a tree).

Need to show

$$\sum_{\sigma \in \text{lin. ext. of } T} \langle \sigma \rangle = \prod_{i \rightarrow j \text{ edges of } T} \langle ij \rangle$$

LHS RHS

Let T' is obtained from T by reversing direction of one edge $i \rightarrow j$



Clearly, $(\text{RHS for } T) = -(\text{RHS for } T')$

$$(\text{LHS for } T) + (\text{LHS for } T')$$

$$= \sum_{\sigma \in S_n} \langle \sigma \rangle$$

$\sigma^{-1}(a) < \sigma^{-1}(b)$ for

all edges $a \rightarrow b$ of

the forest $F := T - (\text{edge } i \rightarrow j)$

Since forest F has 2 connected components T_1 & T_2

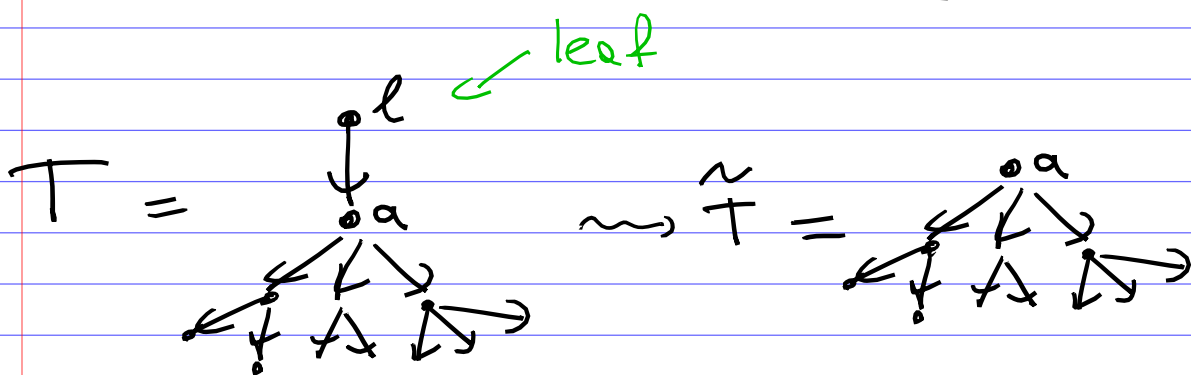
$$= \sum_{\substack{A \text{ lin. ext. of } T_1 \\ B \text{ lin. ext. of } T_2}} \sum_{C \in \text{Shuff}(A, B)} \langle C \rangle$$

$$= 0 \text{ (by lemma 1).}$$

So $(\text{LHS for } T) = -(\text{LHS for } T')$

So when we reverse directions of edges in T , both sides of the needed equality simultaneously change signs.

Let us pick a leaf l of T and direct all edges away from l



any lin. ext. of T starts as l, a, \dots

So LHS for $T = \langle l, a \rangle$. LHS for \tilde{T} ,

$$\tilde{T} = T \setminus \text{leaf } l$$

Clearly, RHS for $T = \langle l, a \rangle$ RHS for \tilde{T}

By induction, needed equality holds for \tilde{T} . Done. \square

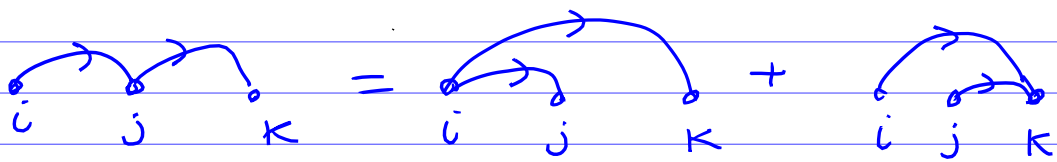
We proved all lemmas, except Lem. 2.

Exercise, Prove Lemma 2:

$$\sum_{C \in \text{Shuffle}(A, B)} \langle l, C, l \rangle = 0$$

for $|A|, |B| \geq 1$.

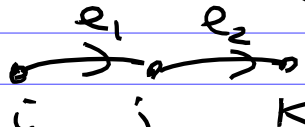
More on the "Key Relation":



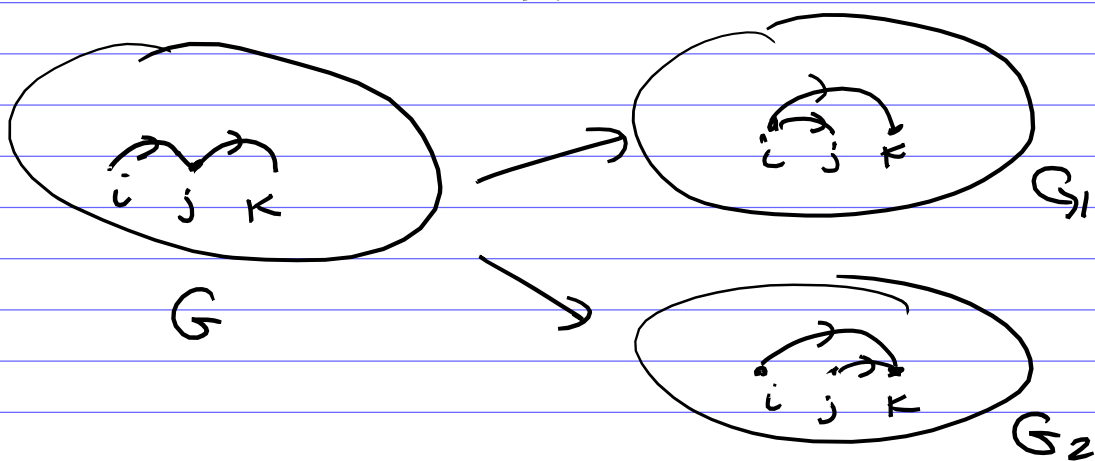
$$\langle ij \rangle \langle jk \rangle = \langle ik \rangle \langle ij \rangle + \langle jk \rangle \langle ik \rangle$$

Let's play the following game on graphs:

- Start with a graph G on vertices $1, \dots, n$ with edges directed as $i \rightarrow j$ for $i < j$

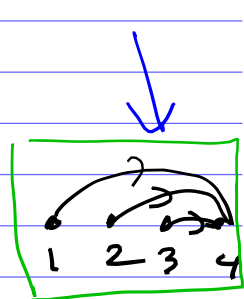
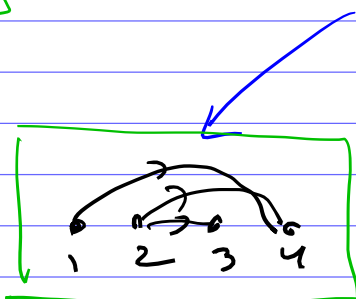
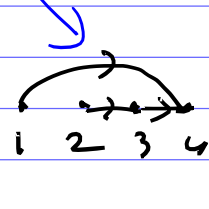
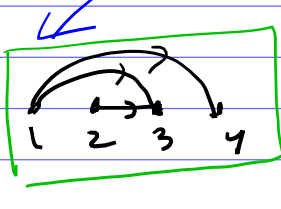
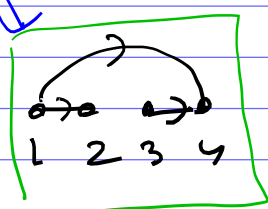
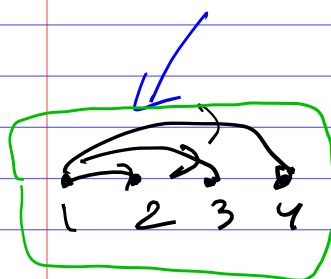
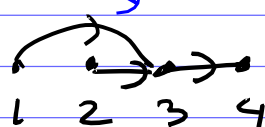
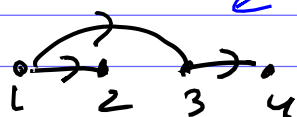
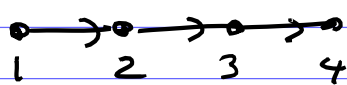
- If we can find 2 edges e_1, e_2 in G s.t. 

Then replace G by the sum of two graphs G_1 & G_2 (in the space of formal linear combinations of graphs) obtained from G by replacing the edges e_1 & e_2 with e_1 & $e_3 = i \rightarrow k$ or e_2 & e_3



- Then apply the same operations to graphs G_1 & G_2 and continue until cannot find a pair of edges $i \rightarrow j \rightarrow k$

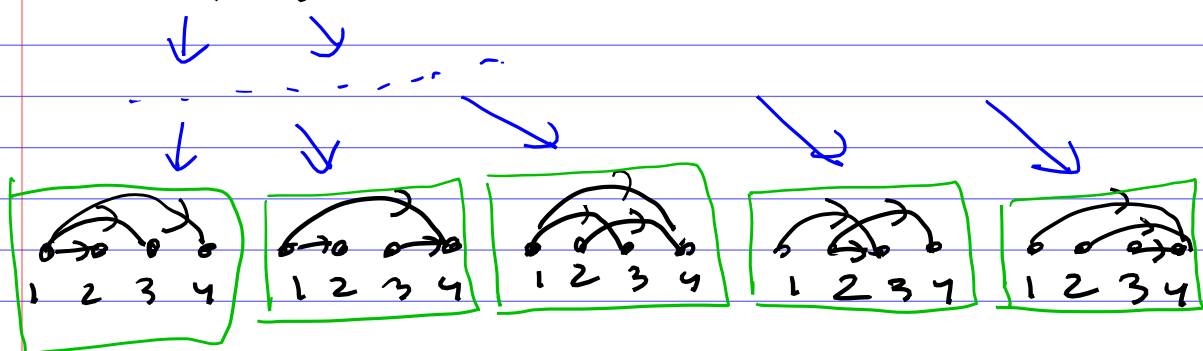
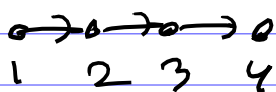
Example



Thus we get the identity:

$$\langle 1 \ 2 \ 3 \ 4 \rangle = \sum_{\substack{S \text{ endpoints} \\ T \text{ of this game}}} \langle T \rangle$$

There are many ways to play this game. Another way to play it produces:



But the number of end-points is the same...

Theorem For any initial graph G , any way to play the game produces the same number of end-points.

Example $G = \begin{array}{ccccccc} \bullet & \longrightarrow & \longrightarrow & \longrightarrow & \longrightarrow & & \\ & 1 & 2 & 3 & \dots & & n \end{array}$

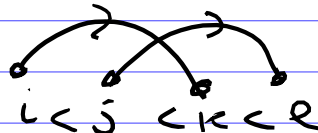
Theorem. (1) # end-points

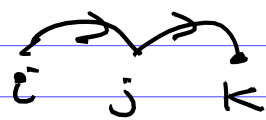
in this case is the

Catalan number $C_{n-1} = \frac{1}{n} \binom{2n-2}{n-1}$

(2) One way to play the game produces the sum over

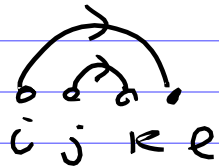
non-crossing alternating
trees on $[n]$.

• non-crossing: no 
 $i < j < k < l$

• alternating: no 
 $i \quad j \quad k$
(every vertex is either a source or a sink)

(3) Another way to play the game produces the sum over

non-nesting alternating trees on $[n]$

• non-nesting: no 
 $i \quad j \quad k \quad l$

A more non-trivial result...

Theorem For $G = K_n$,

the game produces

$C_{n-1} \cdot C_{n-2} \cdots C_1$ end-points

This theorem is equivalent to
Chen-Robbin-Yuen conjecture,
which was proved analytically
by Zeilberger in 1998.

[D. Zeilberger] Proof of a
conjecture of Chen-Robbins-Yuen.

Open problem (for 22 years)

Find a combinatorial proof
of this theorem.

This stuff is closely
related to Kostant's
partition function & flow
polytopes that we discussed
before...

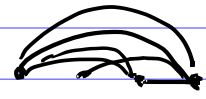
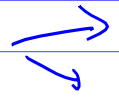
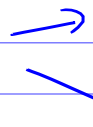
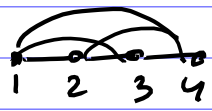
The Above Theorem \Leftrightarrow

Theorem $K(1, 2, \dots, n-1, -\binom{n}{2}) =$
 $= C_1 C_2 \cdots C_{n-1}.$

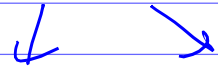
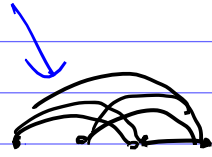
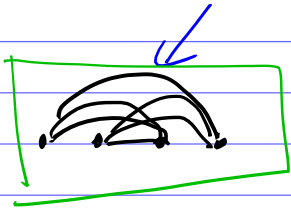
to value of Kostant's partition
function: $\stackrel{\text{def}}{=} \#$ number of ways
to express this vector as
a non-negative integer linear
combination of vectors
 $e_i - e_j, 1 \leq i < j \leq n.$

Example

$$G = K_4$$



one
end point →



This game will have $C_3 \cdot C_2 = 10$
end-points no matter how
you play it.

This stuff is also related to root polytopes

The root polytope

$$R_n := \text{conv}(0, e_i - e_j, 1 \leq i < j \leq n)$$

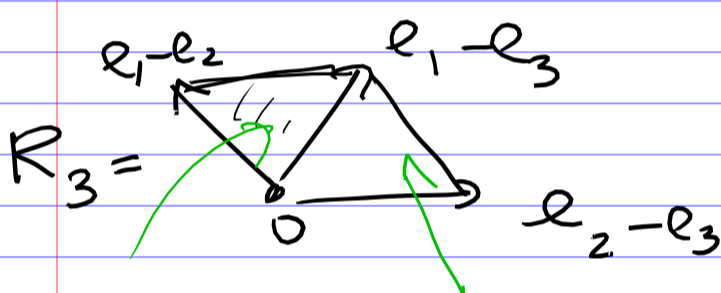
For a tree T on $[n]$

(with edges directed as $i \rightarrow j, i < j$)

$$R(T) := R_n \cap \langle \text{positive span of } e_i - e_j \text{ for edges } i \rightarrow j \text{ in } T \rangle$$

$$R_n = R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ 1 \ 2 \ 3 \ \dots \ n \end{array} \right)$$

Ex. $n=3$



$$R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ 1 \ 2 \ 3 \end{array} \right) \quad R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ 1 \ 2 \ 3 \end{array} \right)$$

$$\text{So } R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ 1 \ 2 \ 3 \end{array} \right) =$$

$$= R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ 1 \ 2 \ 3 \end{array} \right) \cup R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ 1 \ 2 \ 3 \end{array} \right)$$

a union of two polytopes with disjoint interior

More generally,

$$R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ i \ j \ k \end{array} \right) = R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ i \ j \ k \end{array} \right) \cup R \left(\begin{array}{c} \rightarrow \rightarrow \rightarrow \\ i \ j \ k \end{array} \right)$$

So playing the game on trees is equivalent to subdividing root polytopes into smaller pieces.

Theorem [GFP]

There are 2 triangulations of the root polytope R_n into unit simplices

$$R_n = \bigcup_{T \text{ non-crossing alternating tree on } [n]} R(T) \quad \text{the "non-crossing" triangulation}$$

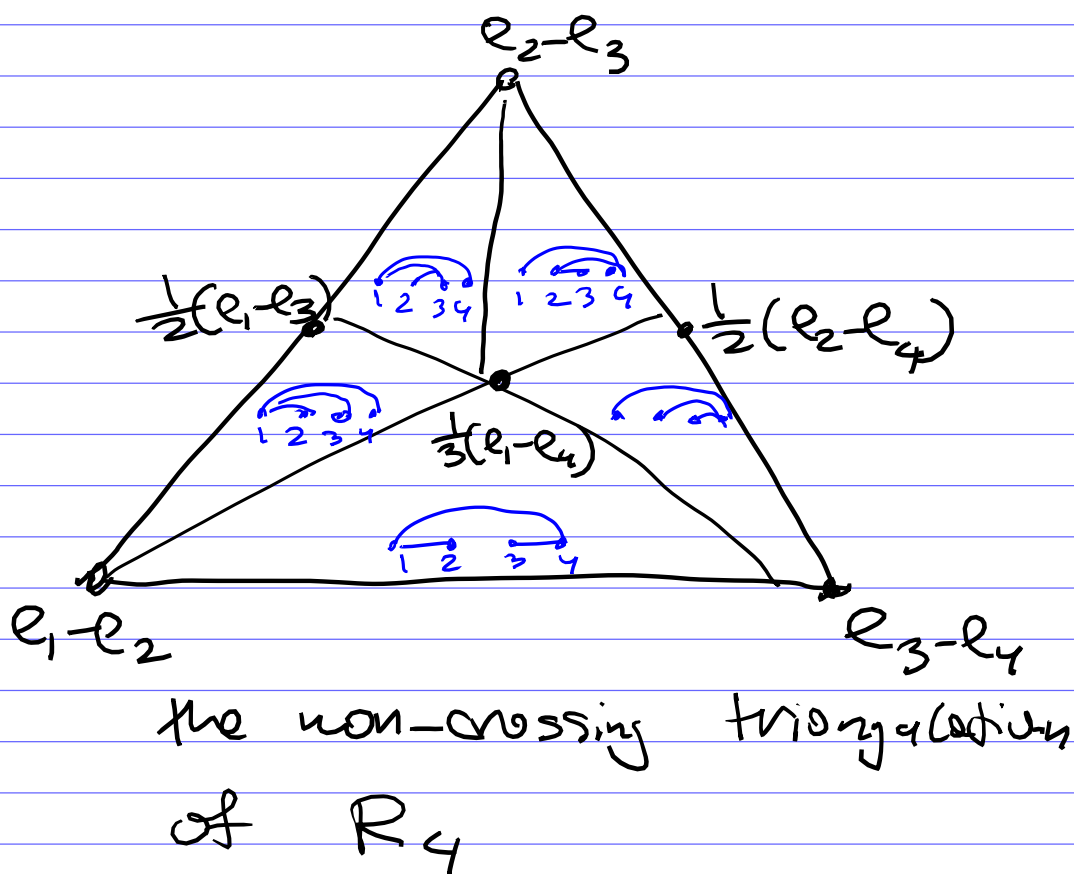
$$= \bigcup_{\tilde{T} \text{ non-nesting alternating tree on } [n]} R(\tilde{T}) \quad \text{the "non-nesting" triangulation}$$

In particular,

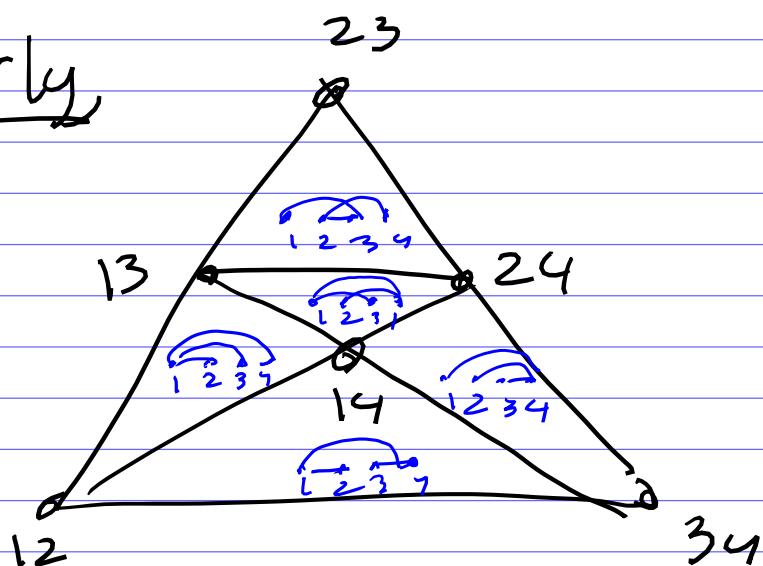
$$\text{Vol}(R_n) = \frac{1}{(n-1)!} C_{n-1}.$$

Example $n=4$,

R_4 is 3-dimensional, but we will represent it by the 2-dim section with the affine plane passing through the points $e_1 - e_2, e_2 - e_3, e_3 - e_4$



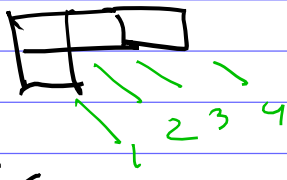
Similarly



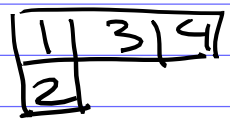
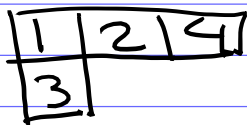
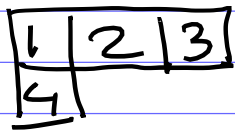
the non-nesting triangulation of R_4

Example Ribbon relation

for



3 SYTs :



$$\langle 2, 3, 4, 1 \rangle + \langle 2, 3, 1, 4 \rangle + \langle 2, 1, 3, 4 \rangle = - \langle 1, 2, 3, 4 \rangle$$

$(-1)^4$

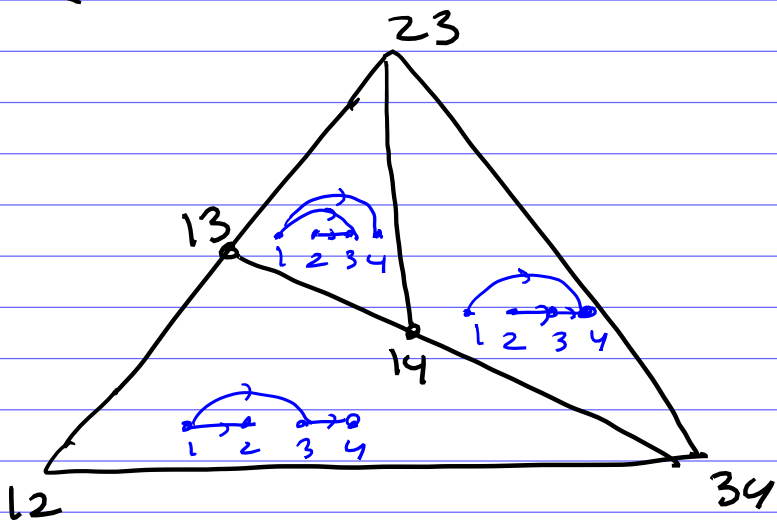
Graphically,



$$= \begin{array}{cccc} \rightarrow & \rightarrow & \rightarrow & \rightarrow \\ 1 & 2 & 3 & 4 \end{array}$$

Geometrically, $R_4 = R(\begin{array}{cccc} \rightarrow & \rightarrow & \rightarrow & \rightarrow \\ 1 & 2 & 3 & 4 \end{array})$

$$= R(\begin{array}{c} \curvearrowright \\ 1 \ 2 \ 3 \ 4 \end{array}) \cup R(\begin{array}{c} \curvearrowright \\ 1 \ 2 \ 3 \ 4 \end{array}) \cup R(\begin{array}{c} \curvearrowright \\ 1 \ 2 \ 3 \ 4 \end{array})$$



Orlik-Terao Algebra

(for the braid arrangement
of type A_{n-1})

$OT_n :=$ the algebra of
rational functions on \mathbb{C}^n
generated by $\langle ij \rangle = \frac{1}{x_i - x_j}$
for $1 \leq i \neq j \leq n$.

Equiv. $OT_n =$ the commutative
algebra over \mathbb{C}
with generators $\langle ij \rangle$,
 $i, j \in [n]$, $i \neq j$, and relations:

- $\langle ij \rangle = -\langle ji \rangle$
- $\langle ij \rangle \langle jk \rangle = \langle ik \rangle \langle ij \rangle + \langle jk \rangle \langle ik \rangle$

Basically, all relations that
we discussed today are
relation in the Orlik-Terao
algebra OT_n .

Clearly, the symmetric group S_n acts on \mathcal{OT}_n by permutations of the coordinates x_i .

$$w \in S_n : \langle ij \rangle \mapsto \langle w(i) w(j) \rangle$$

Some interesting S_n -invariant subspaces of \mathcal{OT}_n

• Tree space

T_n = the space spanned by $\langle T \rangle := \prod_{(ij) \text{ edge of } T} \langle ij \rangle$
 T is a tree on $[n]$

• Forest space

F_n = the space spanned by $\langle F \rangle := \prod_{(ij) \text{ edge of } F} \langle ij \rangle$
 F is a forest on $[n]$

Clearly the forest space is graded by # edges in forests.

$$F_n = F_n^0 \oplus F_n^1 \oplus \dots \oplus F_n^{n-k}$$

F_n^k the space spanned by forests on $[n]$ with k edges.

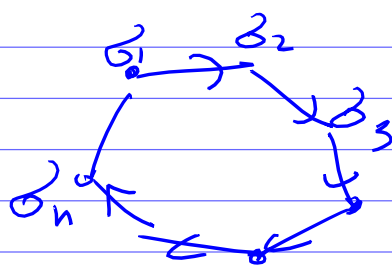
$$\text{Clearly } T_n = F_n^{n-1}.$$

• Perke-Taylor space

PT_n = the subspace of \mathcal{OT}_n spanned by Perke-Taylor factors

$$PT(\sigma) = \langle \sigma_1 \sigma_2 \dots \sigma_n \sigma_1 \rangle$$

$\sigma = \sigma_1 \dots \sigma_n$ is a permutation of $1, 2, \dots, n$.



Each of these spaces T_n , F_n , F_n^k , PT_n is not just a space but a representation of S_n .

Problem Investigate these representations. Find their dimensions, characters, decomposition into irreducibles, etc.

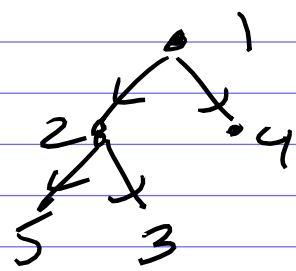
Theorem, • $\dim T_n = (n-1)!$

• $\dim F_n = n!$

• $F_n \cong$ the regular representation of S_n

• 2 linear bases of T_n (& F_n):

- all increasing trees (or forests),



oriented away from a root at 1, the

labels increase as we go away from the root.

- all chains

$\langle 1, i_1, i_2, \dots, i_{n-1} \rangle$

i_1, \dots, i_{n-1} a permutation of $2, 3, \dots, n$.

Example: F_3

$$F_3^0 \cong V \quad \boxed{\square} \\ \text{dim} = 1$$

identity rep
of S_3

$$F_3^1 \cong V \oplus V \quad \boxed{\square} \oplus \boxed{\square} \quad \text{dim} = 3$$

$$F_3^2 \cong T_2 \cong V \oplus V \quad \boxed{\square} \oplus \boxed{\square} \quad \text{dim} = 2$$

Theorem The Hilbert
polynomial of F_n is

$$\sum_{k=0}^{n-1} \text{dim } F_n^k t^k =$$

$$= (1+t)(1+2t) \dots (1+(n-1)t).$$