Analysis of the One Line Factoring Algorithm on Large Semiprimes

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May 18th, 2019 MIT PRIMES Conference What is a factoring algorithm?

Find a divisor of N.

Divide by every prime in $[1...\sqrt{N}]$

Example: N = 119

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119/7 = 17 is an integer!

Simple Factoring Algorithm: Fermat

Factoring ${\cal N}$

- Let $a := \lceil \sqrt{N} \rceil$
- Let $b := a^2 N$
- Repeat until b is a square: Increase a by 1 (a := a + 1) $b := a^2 - N$
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b := 11² - 119 = 2

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- $b := 11^2 119 = 2$
- b (2) is not a square:
 a := a + 1 = 12
 b := 12² 119 = 25

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Works because of square difference $x^2 - y^2 = (x + y)(x - y)$

One Line Factoring Algorithm?

Slower than the leading algorithms

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Much less space required

One line of PARI/GP...

 $OLF(x) =; i = 1; while (i < x, if (issquare (ceil (sqrt(i * x))^2%x), return (gcd(x, floor (ceil (sqrt(i * x)) - sqrt((ceil (i * x))^2)%)))); i++)$

The One Line Factoring Algorithm

Repeat for k = 1 to k = N:

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$$k = 1$$
 to $k = N$:

• Let
$$m := \left\lceil \sqrt{N \cdot k} \right\rceil^2 \% N$$

• If *m* is a square:
Factor is
$$\operatorname{GCD}(N, \left\lceil \sqrt{N \cdot k} \right\rceil - \sqrt{m})$$

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Example: N = 119

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- When k = 1, m = 2
- When k = 2, m = 18

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- When k = 2, m = 18
- When k = 3, m = 4

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- Example: N = 119
- When k = 1, m = 2
- When k = 2, m = 18
- When k = 3, m = 4
- Factor: GCD(119, $\left\lceil \sqrt{119 \cdot 3} \right\rceil - \sqrt{4}$)

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- Let $m := \left[\sqrt{N \cdot k}\right]^2 \% N$
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- Example: N = 119
- When k = 1, m = 2
- When k = 2, m = 18
- When k = 3, m = 4
- Factor: $\begin{array}{c} \operatorname{GCD}(119, \left\lceil \sqrt{119 \cdot 3} \right\rceil - \sqrt{4}) \\ \operatorname{GCD}(119, 17) = 17 \end{array}$

Factor numbers N = pq

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Applications in cryptography (like RSA)



Factoring pq:



Factoring pq:

 X-coordinate is prime p,
 Y-coordinate is prime q,
 p, q are first

1600 primes



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• Green: Smaller prime returned



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- If *p*,*q* are close: Smaller prime returned



Factoring pq:

• X-coordinate is prime p, Y-coordinate is prime q,

p, q are first 1600 primes

- Green: Smaller prime returned
- If *p*,*q* are close: Smaller prime returned
- Probability of green is $\sim 50\%$



Performance of OLF on semiprimes



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Performance of OLF on semiprimes

Number of iterations to factor *pq*:

• X-coordinate is prime p, Y-coordinate is prime q,

p, q are first 1600 primes



Performance of OLF on semiprimes

Number of iterations to factor *pq*:

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• Points colored from black to white; Whiter means more iterations required



The algorithm required trying **every number** from k = 1 to (at most) k = N

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Can we skip some k?

- The algorithm required trying **every number** from k = 1 to (at most) k = N
- Can we skip some k?
- What if we just use squarefree k?



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OLF on Large Semiprimes

• Green: Squarefree approach faster (fewer iterations)



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- Distinct regions where this is more efficient



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- Distinct regions where this is more efficient
- Better on roughly $\sim 35.5\%$ of semiprimes



How many iterations to factor a general integer?



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- k^{th} bar: Amount of integers that requires kiterations to factor
- Decreases rapidly



How many iterations to factor a general integer?



- Decreases rapidly
- Therefore, skipping k will not **always** help.

How many iterations to factor semiprimes?



• However, the picture is different if only factoring semiprimes

How many iterations to factor semiprimes?



- However, the picture is different if only factoring semiprimes
- Many k not used.

How many iterations to factor semiprimes?



• What causes the strange bands?

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- When can we skip k in the general algorithm (not just semiprimes)?

- What causes the strange bands?
- Can we precisely definine when the lower prime is returned?
- Prove the semiprime iterations conjecture.
- When can we skip k in the general algorithm (not just semiprimes)?
- Anything else to make it faster!

• Mentor Yichi Zhang

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- My Family