Minimal Percolating Sets with Time-Dependent Bootstrap Percolation

Yuyuan Luo Mentor: Dr. Laura Schaposnik

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Introduction

Idea

Bootstrap percolation: A deterministic process on a network where nodes become infected once a certain number r of their neighbors are infected.

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 - Orientational ordering processes of magnetic alloys
 - Failure of units in structured collection of computer memory
 - Modeling infectious diseases

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- Applications:
 - Fluid flow in porous areas
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 - Failure of units in structured collection of computer memory
 - Modeling infectious diseases
- Advantages when used to model infectious diseases
 - Individual behavior
 - Spatial aspects
 - Mixing of individuals



t = 0
r = 2



t = 1r = 2



t = 2r = 2



Definition

Suppose we have a graph G(V, E), some positive integer r, and a set $A_0 \subset V$ of initially infected nodes. Then in the r-bootstrap percolation process, the set of infected vertices at time t + 1 is given by

$$A_{t+1} = A_t \cup \{v \in V(G) : |N(v) \cap A_t| \ge r\}$$

where N(v) is the set of adjacent vertices to v.

Notation

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A percolating set A is **minimal** if for any $v \in A$,

$$P(A \setminus \{v\}) \neq V.$$

Example: Minimal Percolating Set



Figure: In this tree, having nodes 2, 5, 6, 7, 8 initially infected is sufficient to ensure that the whole tree is infected at some finite time for r = 2. Further, since each of them has only 1 neighbor, removing any will cause the remaining set to not percolate.

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Motivation

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Instead of r, use F(t).

Definition

Consider a graph G(V, E), a function F(t), and a set $A_0 \subset V$ of initially infected nodes. Then in the F(t)-bootstrap percolation process the set of infected vertices at time t + 1 is

$$A_{t+1} = A_t \cup \{v \in V(G) : |N(v) \cap A_t| \ge \lceil F(t) \rceil\}.$$

Problem Statement

Obtain an algorithm that finds a smallest minimal percolating set for any given tree T(V, E) and percolation function F(t).

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Suppose we have a set of initially infected nodes A_0 . A node v is **isolated** if it is impossible to infect it by time $t_f(v)$ by adding any one other node to A_0 that is not itself.

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Let v be a node that is not isolated. We denote by $t_p(v)$ the largest time t at which F(t) is not larger than the number of infected neighbors of v at t.

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- Step 1. Initialize tree: Let A₀ = Ø. For v ∈ V set t_f(v) arbitrarily large, and set it to true for needing to be considered.
- Step 2. Percolate using current A_0 . Stop the algorithm if $P(A_0) = V$: the resulting A_0 is a smallest minimal percolating set.
- Step 3. Infection. For an unconsidered v furthest away from the root (isolated first).
 - if v is isolated, add v to A_0 .
 - otherwise, set t_f of the parent of v to $t_p(v) 1$ if it is smaller than the current t_f of the parent.

• Step 4. Go to step 2.



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F(t) = t

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Size of Minimal Percolating Sets on Perfect Trees with Height 4



F(t)=4 F(t) = 4t^2

Figure: The size of smallest minimal percolating sets on perfect trees with height 4, with a constant and a non-constant percolation function 10 / 12

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Minimal Percolating Sets

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- Devise a similar algorithm for largest minimal percolating set.
- Study the size of largest and smallest minimal percolating sets in specific trees and lattices, with specific types of percolation functions.

- Dr. Laura Schaposnik
- Dr. Tanya Khovanova
- My parents
- MIT PRIMES program