# G-Parking Functions and Monomial Ideals 

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## Definitions

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Theorem (Cayley)
The complete graph $K_{n+1}$ has $(n+1)^{n-1}$ spanning trees.

## G-Parking Functions

- $G$ is a digraph on vertices $\{0,1, \ldots, n\}$
- For a nonempty subset $I \subseteq\{1, \ldots, n\}$, and vertex $i \in I$, let $d_{l}(i)$ denote the number of edges from $i$ to vertices outside $I$



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- A $G$-parking function is an $n$-tuple $\left(b_{1}, \ldots, b_{n}\right)$ such that for any nonempty subset $I \subseteq\{1,2, \ldots, n\}$, there exists $i \in I$ such that $b_{i}<d_{l}(i)$


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- Example: $(0,1,1)$ is a $G$-parking function, where $G$ is the graph above


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- Classical parking functions are the special case $G=K_{n+1}$
- Chebikin and Pylyavskyy constructed an explicit bijection


## Chebikin-Pylyavskyy Bijection

- For every subtree $T$ of $G$ rooted at 0 , assign an order $\pi(T)$ to $T$ 's vertices. Let $i<_{\pi(T)} j$ denote $i$ being smaller than $j$ in this order


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- An choice of orders $\Pi(G)$ is a proper set of tree orders if for each subtree $T$ rooted at 0 :
- if an edge $(i, j) \in T$, then $i>_{\pi(T)} j$
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- Example: Breadth-first search order


$$
0<_{\pi(T)} 1<_{\pi(T)} 3<_{\pi(T)} 2
$$

## Chebikin-Pylyavskyy Bijection

- Fix a proper set of tree orders $\Pi(G)$
- For each spanning tree $T$, let $e(T, i)$ be the edge out of $i$ in $T$
- Given a subtree $T$ and order $\pi(T)$, for each vertex $i$, order the edges from $i$ to $T$ such that $\left(i, j_{1}\right)<_{\pi(T)}\left(i, j_{2}\right)$ if $j_{1}<_{\pi(T)} j_{2}$


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Theorem (Chebikin, Pylyavskyy)
Map each spanning tree $T$ to $\left(b_{1}, \ldots, b_{n}\right)$, where $b_{i}$ is the number of edges $e$ from $i$ such that $e<_{\pi(T)} e(T, i)$. This mapping is a bijection between $G$ 's spanning trees rooted at 0 and G-parking functions.

## Chebikin-Pylyavskyy Bijection - An Example



- G-parking functions: $(0,0,0),(1,0,0),(0,1,0),(0,0,1),(1,1,0)$, (0,1,1)


## Chebikin-Pylyavskyy Bijection - An Example



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- Spanning trees:



## Monomial Ideals

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- For each nonempty subset $I \subseteq\{1, \ldots, n\}$, define

$$
m_{l}=\prod_{i \in I} x_{i}^{d_{l}(i)}
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and let the ideal $\mathcal{I}_{G}=\left\langle m_{l}\right\rangle$ as $/$ ranges over all nonempty subsets of $\{1, \ldots, n\}$

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- $\left(b_{1}, \ldots, b_{n}\right)$ is a $G$-parking function if and only if $x_{1}^{b_{1}} \cdots x_{n}^{b_{n}}$ does not vanish in $\mathbb{K}\left[x_{1}, \ldots, x_{n}\right] / \mathcal{I}_{G}$


## Monomial Ideals - An Example



- $\mathcal{I}_{G}=\left\langle x_{1}^{2}, x_{2}^{2}, x_{3}^{2}, x_{1}^{2} x_{2}, x_{1} x_{3}, x_{2} x_{3}^{2}, x_{1} x_{2}^{0} x_{3}\right\rangle$


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- Non-vanishing monomials: $1, x_{1}, x_{2}, x_{3}, x_{1} x_{2}, x_{2} x_{3}$
- G-parking functions: $(0,0,0),(1,0,0),(0,1,0),(0,0,1),(1,1,0)$, $(0,1,1)$


## Almost-G-Parking Functions

- For each nonempty subset $I=\left\{i_{1}<\cdots<i_{k}\right\} \subseteq\{1, \ldots, n\}$, define

$$
\hat{m}_{I}=x_{i_{1}} \prod_{i \in I} x_{i}^{d_{l}(i)}
$$

and let the ideal $\hat{\mathcal{I}}_{G}=\left\langle\hat{m}_{I}\right\rangle$ as $/$ ranges over all nonempty subsets of $\{1, \ldots, n\}$

- $\left(b_{1}, \ldots, b_{n}\right)$ is an almost-G-parking function if $x_{1}^{b_{1}} \cdots x_{n}^{b_{n}}$ does not vanish in $\mathbb{K}\left[x_{1}, \ldots, x_{n}\right] / \hat{\mathcal{I}}_{G}$


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- $\left(b_{1}, \ldots, b_{n}\right)$ is an almost-G-parking function if $x_{1}^{b_{1}} \cdots x_{n}^{b_{n}}$ does not vanish in $\mathbb{K}\left[x_{1}, \ldots, x_{n}\right] / \hat{\mathcal{I}}_{G}$
- Theorem (Postnikov, Shapiro, Shapiro) When $G=K_{n+1}$, the number of almost-G-parking functions equals the number of (undirected) spanning forests of $G$.


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- We explicity construct a bijection between almost-G-parking functions and spanning forests of $G$ whose connected components are rooted at their smallest vertices


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- For every subtree $T$ of $G$ rooted at its numerically smallest vertex, assign an order $\hat{\pi}(T)$. Let $i<_{\hat{\pi}(T)} j$ denote $i$ being smaller than $j$ in this order
- A choice of orders $\hat{\Pi}(G)$ is a super-proper set of tree orders if for each subtree $T$ rooted at its numerically smallest vertex:
- if an edge $(i, j) \in T$, then $i>_{\hat{\pi}(T)} j$
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- Example: Breadth-first search order


## Almost-G-Parking Functions and Spanning Forests

- Given a super-proper set of tree orders $\hat{\Pi}(G)$, for every spanning forest $F$ of $G$ whose connected components are rooted at their numerically smallest vertices, assign an order $\pi(F)$ such that $i<_{\pi(F)} j$ if:
- the root of $i$ 's connected component is smaller than the root of $j$ 's connected component, or
- $i$ and $j$ are in the same connected component $T$, and $i<_{\hat{\pi}(T)} j$


## Almost-G-Parking Functions and Spanning Forests

- Fix a super-proper set of tree orders $\hat{\Pi}(G)$
- Let $e(F, i)$ be the edge out of $i$ in the spanning forest $F$, if it exists
- Given a subforest $F$ and order $\pi(F)$, for each vertex $i$, order the edges from $i$ to $F$ such that $\left(i, j_{1}\right)<_{\pi(F)}\left(i, j_{2}\right)$ if $j_{1}<_{\pi(F)} j_{2}$


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- Map each spanning forest $F$ whose connected components are rooted at their numerically smallest vertices to ( $b_{1}, \ldots, b_{n}$ ), where $b_{i}$ is:
- the number of edges from $i$ to vertices smaller than $i$ in $\pi(F)$, if $i$ is the root of its connected component
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## Theorem

This mapping is a bijection between almost-G-parking functions and $G$ 's spanning forests whose connected components are rooted at their numerically smallest vertices.

## An Example



This corresponds to the almost- $G$-parking function $(0,2,1)$

## Modified Monomial Ideals

For each nonempty $I \subseteq\{1, \ldots, n\}$, choose any $k_{I} \in I$ and let

$$
\hat{m}_{l}^{\prime}=x_{k_{l}} \prod_{i \in I} x_{i}^{d_{l}(i)}
$$

Let $\hat{\mathcal{I}}_{G}^{\prime}=\left\langle\hat{m}_{I}^{\prime}\right\rangle$ as $I$ ranges over all nonempty subsets of $\{1, \ldots, n\}$, and let $\hat{\mathcal{A}}_{G}^{\prime}=\mathbb{K}\left[x_{1}, \ldots, x_{n}\right] / \hat{\mathcal{I}}_{G}^{\prime}$
Theorem
If $G=K_{n+1}$, then $\operatorname{dim} \hat{\mathcal{A}}_{G}^{\prime}$ is independent of the choices of $k_{l}$.

## Future Directions

## Conjecture

$\operatorname{dim} \hat{\mathcal{A}}_{G}^{\prime}$ is independent of the choices of $k_{I}$ for all choices of $k_{I}$ that preserve the monotonicity of the ideal $\hat{\mathcal{I}}_{G}^{\prime}$ (i.e. if $I \subset J$, then for any $\left.i \in I, \operatorname{deg}_{x_{i}} \hat{m}_{l}^{\prime} \geq \operatorname{deg}_{x_{i}} \hat{m}_{J}^{\prime}\right)$.
It would also be interesting to find a combinatorial interpretation of ideals in which the $m_{l}$ are modified by multiplication by more than one variable

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