# Efficient Calculation of Determinants of Symbolic Matrices with Many Variables 

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## Iliit itill

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## Vectors and Volumes



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## Determinants

$2 \times 2$ matrices:

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\operatorname{det}\left(\begin{array}{ll}
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$n \times n$ matrices:

$$
\operatorname{det}(A)=\sum_{\sigma \in S_{n}}\left[\operatorname{sgn}(\sigma) \prod_{i=1}^{n} A_{i, \sigma(i)}\right]
$$

## Motivation

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We are interested in matrices with polynomial entries.

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Minor expansion requires $\sum_{i=2}^{n} i\binom{n}{i} \in \Theta\left(2^{n} n\right)$ polynomial multiplications.

## Gaussian Elimination

$$
\operatorname{det}\left(\begin{array}{cccc}
a & b & c & d \\
0 & f & g & h \\
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e-\frac{e}{a} a & f-\frac{e}{a} b & g-\frac{e}{a} c & \ldots \\
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\begin{aligned}
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\end{array}\right) \\
A^{(1)}=A, \\
A_{i, j}^{(k+1)}=A_{i, j}^{(k)}-\frac{A_{i, k}^{(k)}}{A_{k, k}^{(k)}} A_{k, j}^{(k)}
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\end{gathered} A_{i, j}^{(k+1)}=\frac{A_{i, j}^{(k)} A_{k, k}^{(k)}-A_{i, k}^{(k)} A_{k, j}^{(k)}}{A_{k-1, k-1}^{(k-1)}} .
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\end{array}\right.
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Fraction-free Gaussian elimination requires $\sum_{i=1}^{n} \Theta\left(i^{2}\right) \in \Theta\left(n^{3}\right)$ polynomial
multiplications and divisions.

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cost ratio $=\frac{\operatorname{cost} \text { of ME }}{\operatorname{cost} \text { of FFGE }}$


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- Absolute value of determinant is invariant under row swaps.


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- Product of one more than number of terms for each entry of a row, $\prod_{i=1}^{n}\left(\operatorname{nterms}\left(r_{i}\right)+1\right)$.


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- Crossover points between algorithms.
- Machine learning.


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