# Gleason's theorem and unentangled orthonormal bases 

Nolan R. Wallach
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## Gleason's theorem

- Let $\mathcal{H}$ be a separable Hilbert space with unit sphere $S$. Then

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f: S \rightarrow \mathbb{R}_{\geq 0}
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is called a frame function if for every orthonormal basis of $\mathcal{H},\left\{e_{i}\right\}$,

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- If $\operatorname{dim} \mathcal{H} \geq 3$ then every frame function is of this form.
- Gleason's proof uses a little representation theory, a reduction to 3 real dimensions and geography of the 2 sphere.


## A little quantum mechanics

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- To a pure state, $v$, we form the linear map $v \otimes v^{\dagger}$. A mixed state is a limit of convex combinations of pure states. Thus mixed state is a positive semidefinite trace class operator of trace 1 .


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- To a pure state, $v$, we form the linear map $v \otimes v^{\dagger}$. A mixed state is a limit of convex combinations of pure states. Thus mixed state is a positive semidefinite trace class operator of trace 1.
- So Gleason's theorem gives an operational interperatation of mixed states and has been used argue against hidden variables in quantum mechanics.


## Two dimensions

- We assume that $\operatorname{dim} \mathcal{H}=2$. We note that if $v \in \mathcal{H}$ is a unit vector then there is a unique, up to phase, unit vector $\widehat{v}$ orthogonal to it. This yields a map $\mathbb{P}^{1} \rightarrow \mathbb{P}^{1},[v] \longmapsto[\widehat{v}]$.


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- Fix such an $X$ and $g: X \rightarrow[0,1]$. Then if we define $f(v)=g([v])$ and $f(\widehat{v})=1-g([v])$ then we have defined a frame function and this type of function is the most general one.


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- If $m>1, d_{i}>1$ then a randomly chosen state will be entangled. Since the dimension of the set of states is $d_{1} \cdots d_{n}-1$ and the dimension of the set of product states is $d_{1}+\ldots+d_{m}-n+1$. Thus if $m>1$ and all $d_{i}>1$ almost all states are entangled.


## Unentangled Gleason theorem

- If $\left\{\phi_{i j}\right\}_{0 \leq j<\operatorname{dim} \mathcal{H}_{i}}$ is an orthonormal basis of $\mathcal{H}_{i}$ then the orthonormal basis $\left\{\phi_{1 i_{1}} \otimes \cdots \otimes \phi_{n i_{n}}\right\}$ is called a product basis. If $\left\{\varphi_{k}\right\}$ is an orthonormal basis and each basis element is a product vector then we will call the basis an unentangled basis.


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- Suppose that $\mathcal{H}=\mathcal{H}_{1} \otimes \mathcal{H}_{2} \otimes \cdots \otimes \mathcal{H}_{n}$ and we consider $\Sigma=\left\{\phi_{1} \otimes \cdots \otimes \phi_{n} \mid\left\|\phi_{i}\right\|=1, \phi_{i} \in \mathcal{H}_{i}\right\}$ and consider the functions $f: \Sigma \rightarrow \mathbb{R}_{\geq 0}$ such that for each unentangled basis $\left\{\varphi_{k}\right\}$ we have $\sum f\left(\varphi_{k}\right)=1$. That is, unentangled frame functions.


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- If $\operatorname{dim} \mathcal{H}_{i} \geq 3$ for all $i=1, \ldots, n$ then there exists $T$ a mixed state in $\mathcal{H}$ such that $f(\varphi)=\langle T \varphi \mid \varphi\rangle$ for all product states $\varphi$.
- Suppose $\mathcal{H}=\mathbb{C}^{2} \otimes V$ a tensor product Hilbert space. We assert that in this context there is an unentangled frame function that is not given by a mixed state.
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- We first describe all unentangled orthonormal bases. Let $z_{i}=a_{i} \otimes b_{i}$ be such an orthonormal basis. If $i \neq j$ and if $\left\langle a_{i} \mid a_{j}\right\rangle \neq 0$ then $\left\langle b_{i} \mid b_{j}\right\rangle=0$.
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- There is an orthogonal decomposition $V=\oplus V_{j}$ and for each $j$ two orthonormal bases $v_{j i}$ and $w_{j i}$ of $V_{j}$ and $c_{j}$ a unit vector in $\mathbb{C}^{2}$ such that the the basis is

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\cup_{j}\left(\left\{c_{j} \otimes v_{j i}\right\}\right) \cup\left(\left\{\widehat{c}_{j} \otimes w_{j i}\right\}\right)
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Using this result we easily see

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- If $f$ is a frame function for $\mathbb{C}^{2}$ and if $g$ is one for $V$ then $f \otimes g$ is an unentangled one for $\mathbb{C}^{2} \otimes V$.


## Qubits

- We now concentrate on the case when all of the spaces $\mathcal{H}_{i}$ are equal to $\mathbb{C}^{2}$. The case that is most important in quantum information theory. We will denote the Hilbert space $Q B_{n}=\otimes^{n} \mathbb{C}^{2}$. From here on I am discussing joint work with Jiri Lebl and Asif Shakeel.


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- A UOB for $Q B_{n}$ is $2^{n}$ vectors $z_{j}=a_{1}^{j} \otimes a_{2}^{j} \otimes \cdots \otimes a_{n}^{j}, j=1, \ldots, 2^{n}$ with $a_{i}^{j}$ a unit vector in $\mathbb{C}^{2}$ and for each $j \neq k$ there exists $i$ such that up to phase $a_{i}^{k}=\widehat{a_{i}^{j}}$. Thus if we think of a UOB projectively $\left[z_{j}\right]=\left[a_{1}^{j}\right] \otimes\left[a_{2}^{j}\right] \otimes \cdots \otimes\left[a_{n}^{j}\right]$ (i.e. using the Segre imbedding of $\left(\mathbb{P}^{1}\right)^{n}$ into $\mathbb{P}^{2^{n}-1}$ ). So the notion of UOB is purely combinatorial.


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- Let $S_{i}$ be the set $\left\{\left[a_{i}^{j}\right] \mid j=1, \ldots, 2^{n}\right\}$. We have seen that the involution $S_{i} \rightarrow S_{i}$ given by $u \longmapsto \hat{u}$ is fixed point free. We can therefore choose $U_{i} \subset S_{i}$ a fundamental domain. We write $U_{i}=\left\{u_{i 1}, \ldots, u_{i k_{i}}\right\}$. We also note that the results above imply that the number of $j$ such that $\left[a_{i}^{j}\right]=u$ is the same as the number of $j$ such that $\left[a_{i}^{j}\right]=\hat{u}$.
- For each $j$ we put together an $n$-bit string $b(j)$ with $i$-th bit 0 if $\left[a_{i}^{j}\right]=u \in U_{i}$ or 1 if it is $\hat{u}$ with $u \in U_{i}$. Set $v_{j}=e_{b_{1}} \otimes e_{b_{1}} \otimes \cdots \otimes e_{b_{n}}$ with $e_{0}, e_{1}$ the standard orthonormal basis of $\mathbb{C}^{2}$ and $b=b(j)$.
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- The vectors are orthonormal. Thus there is a permutation, $\sigma$, of $1, \ldots, 2^{n}$ so that $b(\sigma j)$ is the base two expansion of $j-1$ padded by 0 's on the left. We will say that a UOB is in normal order if $\sigma$ is the identity. Assume that the UOB is in normal order.
- Set $k=\sum k_{i}$. We give a coloring of the hypercube graph with $k$ colors corresponding to the UOB. Recall that the $n$-th hypercube graph has vertices $0, \ldots, 2^{n-1}$ and edges $[i, j]$ with $i \neq j$ having all binary digits the same except for one $I=I(i, j)=I(j, i)$. We set $m=k_{1}+\ldots+k_{l-1}+p$ if

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\left[a_{l(i, j)}^{i}\right]=u,\left[a_{l(i, j)}^{j}\right]=\hat{u}
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- Set $\alpha_{i}^{j}=m$ if the edge eminating from $j$ in the direction $i$ has color $m$.

$$
a \otimes b, a \otimes \hat{b}, \hat{a} \otimes c, \hat{a} \otimes \hat{c},\left[\begin{array}{llll}
1 & 1 & 1 & 1 \\
2 & 2 & 3 & 3
\end{array}\right]
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- $\left[\begin{array}{llll}1 & 1 & 1 & 1 \\ 2 & 2 & 3 & 3\end{array}\right] \Leftrightarrow\left(u_{1}, u_{2}, u_{3}\right) \longmapsto$
$\left\{u_{1} \otimes u_{2}, u_{1} \otimes \widehat{u_{2}}, \widehat{u_{1}} \otimes u_{3}, \widehat{u_{1}} \otimes \widehat{u_{3}}\right\}$ a parametrized family of UOB parametrized by $\left(\mathbb{P}^{1}\right)^{3}$.
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- If we consider the locus of this parameterized subset of the and the one that comes by rotating the square by $\frac{\pi}{2}$ the union of the sets is all UOB in 2 qubits. The intersection of these sets is the locus parameterization that comes from using the red as above but replacing (say) the green by blue.
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- To get a better idea of the pattern consider $Q_{3}$. In this case the maximal number of colors that will come from a UOB is 7 and up to similar rotations to the case above the only coloring with 7 colors coming from a UOB is
cyan, red, green,orange,yellow, brown,blue $\left[\begin{array}{cccccccc}1 & 1 & 5 & 5 & 1 & 1 & 5 & 5 \\ 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 \\ 3 & 3 & 4 & 4 & 7 & 7 & 6 & 6\end{array}\right]$



- This example indicates a surprising (to us complication). Notice that it has 6 colors. We also note (not easy) that adding a color (in other words taking some colors that appear multiple times and changing a proper subset of the edges of each of these colors into the new color) is not possible.
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- We can go from a coloring of $Q_{n}$ with $k$ colors to a map, $\Phi$, of $\left(\mathbb{P}^{1}\right)^{k}$ to $\left(\otimes^{n} \mathbb{C}^{2}\right)^{2^{n}}$ as follows: Let $j$ be a vertex of $Q_{n}$ then for every $1 \leq i \leq n$ there is exactly one vertex that differs in exactly that digit from $j, j_{i}$. If the edge $\left[j, j_{i}\right]$ is colored with the $l_{i}$-th color. Then we take the $j$-th component of $\Phi\left(u_{1}, \ldots, u_{k}\right)$ to be

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v_{1} \otimes v_{2} \otimes \cdots \otimes v_{n}, v_{i}=\left\{\begin{array}{l}
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- We will say that the coloring is useful if it yields a family of UOB. We will say that a useful family is maximal if adding a color makes it useless.


## Theorems

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- The useful colorings with $2^{n}-1$ colors are constructed recursively as follows: There is one direction, say $i$, with exactly one color on every edge in that direction coordinates which we call vertical. We then have two $n-1$ cubes $Q^{0}$ ( $i$-th coordinate 0 ) and $Q^{1}$ each of these has $2^{n-1}-1$ with a total of $2^{n}-2$ colors (since adding 1 to the total number of colors of the two subcubes is $2^{n}-1$. So we see that each has one direction in each subcube with all colors the same and continue.


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- Note that the directions in the subcube can be different. If we take them to be the same then we have the examples above for $n=2,3,4$.
- If $n=3$ then up to permutation of order and permutation of vectors the the two types we gave are all of the maximal useful colorings.


