Due: 4/18/2018

Description

These problems are related to the material covered in Lectures 14-17. Your solutions are to be written up in latex and submitted as a pdf-file with a filename of the form SurnamePset4.pdf via e-mail to drew@math.mit.edu by noon on the date due. Collaboration is permitted/encouraged, but you must identify your collaborators, and any references consulted. If there are none, write "Sources consulted: none" at the top of your problem set. The first person to spot each non-trivial typo/error in any of the problem sets or lecture notes will receive 1-5 points of extra credit.¹

Instructions: Pick any combination of problems to solve that sum to 100 and write up your answers in latex (be sure to include Problem 4, which is a survey).

Problem 1. Modular forms on $SL_2(\mathbb{Z})$ (49 points)

Throughout this problem $\Gamma := \mathrm{SL}_2(\mathbb{Z})$ denotes the full modular group. Let Γ_{∞} be the stabilizer of ∞ , and let

$$\mathcal{F} := \{ z \in \mathbf{H} : |z| \ge 1 \text{ and } |\operatorname{Re}(z)| \le 1/2 \}$$

denote the standard fundamental domain for Γ . For even k > 2 we define the *Eisenstein series* $G_k(z)$ as the Poincaré series (see (2.6.2) in [1]):

$$G_k(z) := F_k(z; \phi_0, \chi_0, \Gamma_\infty, \Gamma) := \sum_{\gamma \in \Gamma_\infty \setminus \Gamma} \overline{\chi_0(\gamma)}(\phi_0|_k \gamma)(z),$$

where $\phi_0 = 1$ is the identity function and χ_0 is the trivial character. We also define the normalized Eisenstein series

$$E_k(z) := -\frac{B_k}{2k}G_k(x),$$

where $B_k \in \mathbb{Q}$ is the kth Bernoulli number (note: this is **not** the E_k defined in [1, §4.1]).

(a) Let $k \geq 4$ be even. Prove that $E_k(z)$ has the following q-expansion at ∞ :

$$E_k(z) = \frac{-B_k}{2k} + q + \sum_{n=2}^{\infty} \sigma_{k-1}(n)q^n,$$

where $q := e^{2\pi i z}$ and $\sigma_e(n) := \sum_{d|n} d^e$.

(b) Let $\rho = e^{\pi i/3}$. Prove that for any $k \in \mathbb{Z}$ and nonzero $f \in M_k(\Gamma)$ we have

$$\operatorname{ord}_{\infty}(f) + \frac{1}{2}\operatorname{ord}_{i}(f) + \frac{1}{3}\operatorname{ord}_{\rho}(f) + \sum_{\tau \in \mathcal{F} = \{i, \rho\}} \operatorname{ord}_{\tau}(f) = \frac{k}{12}.$$

Conclude that dim $M_k = 0$ unless $k \ge 4$ is even.

¹This is a long and completely new problem set; there are surely many points to be gained.

- (c) Let $\Delta := (G_4^3 G_6^2)/1728$. Prove that for even $k \geq 4$ the map $f \mapsto \Delta f$ defines an isomorphism $M_{k-12} \xrightarrow{\sim} S_k$. Conclude that $\dim M_k = 1$ and $\dim S_k = 0$ for k = 4, 6, 8, 10, 14.
- (d) Prove that if $4a + 6b = k \equiv 0 \mod 12$ then $3|a \text{ and } G_4^a G_6^b / G_6^{k/6} = (G_4^3 / G_6^2)^{a/3}$.
- (e) Prove that for all even integers $k \geq 2$ we have

$$\dim M_k(\Gamma) = \begin{cases} 0 & \text{if } k \text{ is odd or negative,} \\ \lfloor k/12 \rfloor & \text{if } k \equiv 2 \bmod 12 \text{ is positive and even.} \\ \lfloor k/12 \rfloor + 1 & \text{if } k \not\equiv 2 \bmod 12 \text{ is positive and even.} \end{cases}$$

(this is Corollary 4.1.4 in [1], you are being asked to give an alternative proof).

- (f) Prove that $\{G_4^a G_6^b : a, b \in \mathbb{Z}_{\geq 0} \text{ with } 4a + 6b = k\}$ is a basis for $M_k(\Gamma)$ for all $k \in \mathbb{Z}$. (this is Theorem 4.1.8 in [1], you are being asked to give an alternative proof).
- (g) Prove that the space $S_k(\Gamma)$ has a unique basis $\{f_1, \ldots, f_d\}$ such that $a_i(f_j) = \delta_{ij}$ for $1 \leq i, j \leq d$, where f_j has q-expansion $f_j = \sum_{i \geq 0} a_i(f_j)q^i$ at ∞ , and show that $f_j \in \mathbb{Z}[[q]]$ (hint: use the E_k and Δ). This basis for $S_k(\Gamma)$ is called the *Miller basis*.
- (h) Let $\{f_1, f_2\}$ be the Miller basis for $S_{38}(\Gamma)$. Compute the integers $a_3(f_i)$ for $1 \le i \le 2$ and express each f_i in terms of the basis for $M_{38}(\Gamma)$ given by part (f).
- (i) Let $\mathbb{T} := \mathbb{T}_{\mathbb{Z}}(\Gamma_0(1), \Delta_0(1))$ be the Hecke algebra for $\Gamma = \Gamma_0(1)$ over \mathbb{Z} , with $\Delta_0(1)$ the set of integer matrices in $\mathrm{GL}_2^+(\mathbb{Q})$ (as in (4.5.1) of [1]). Let $T(n) \in \mathbb{T}$ be the Hecke operator $T(n) := \sum_{\det \alpha = n} \Gamma \alpha \Gamma$, where the sum is over double cosets in $\Delta_0(1)$. Show that for every even integer $k \geq 4$ and integer $n \geq 1$ we have $T(n)(E_k) := E_k | T(n) = \sigma_{k-1}(n) E_k$, and compute the matrix of T(2) on the Miller basis for $S_{38}(\Gamma)$ (you can use sage to check your answer).

Problem 2. Computing Bernoulli numbers (49 points)

Recall from Problem 3 of Problem Set 8 that the Bernoulli numbers B_n arise as the coefficients of the exponential generating function

$$\frac{x}{e^x - 1} = \sum_{n > 0} B_n \frac{x^n}{n!}$$

and satisfy the recursive formula defined by $B_0 = 1$ and

$$B_n = -\frac{1}{n+1} \sum_{k=0}^{n-1} \binom{n+1}{k} B_k.$$

For a nonzero rational number r = a/b with $a, b \in \mathbb{Z}$ coprime we define its height $h(x) := \max(\log |a|, \log |b|)$.

(a) Let $n \geq 2$ be an even integer. Prove that $v_p(B_n) \geq -1$ for all primes p and that $v_p(B_n) = -1$ if and only if p-1 divides n. In other words, when written in lowest terms the denominator of the rational number B_n is $\prod_{(p-1)|n} p$.

(b) Let $n \ge 2$ be an even integer. Using the identity $\zeta(n) = (-1)^{n/2-1} \frac{(2\pi)^n B_n}{2 \cdot n!}$ and the bounds

$$\sqrt{2\pi}n^{n+\frac{1}{2}}e^{-n}e^{\frac{1}{12n+1}} < n! < \sqrt{2\pi}n^{n+\frac{1}{2}}e^{-n}e^{\frac{1}{12n}}$$

determine an effective upper bound for $h(B_n)$ as a function of n.

- (c) Let M(b) denote the time to multiply two b-bit integers (the Schönhage-Strassen algorithm yields $M(b) = O(b \log b \log \log b)$). Estimate the complexity of computing B_n using the recursive formula above, assuming that it takes $M(b) \log b$ time to reduce a ratio of b-bit integers to a rational number in lowest terms (this is achieved by the fast Euclidean algorithm). Your answer should be an asymptotic upper bound (in big O-notation) on the number of bit operations required, as a function of n.
- (d) Describe an algorithm to compute B_n using Newton iteration to compute 1/f modulo x^{n+1} , where $f(x) = (e^x 1)/x = \sum_{n \geq 0} \frac{x^n}{(n+1)!}$, and estimate its complexity. Note that you can use Kronecker substitution to multiply two polynomials in $\mathbb{Z}[x]$ of degree at most d with coefficients of height b in time $O(\mathsf{M}(bd\log d))$ by replacing x with a suitable power of 2.
- (e) Let $n \geq 2$ be an even integer, let $d = \prod_{(p-1)|n} p$ be the denominator of B_n and let A be an approximation to $2 \cdot n!/(2\pi)^n$ to b-bits of precision, such that $2^b > e^{h(B_n)}$. Let $P := \lceil (Ad)^{1/(n-1)} \rceil$ and let $r := \prod_{p \leq P} (1-p^{-n})^{-1}$. Show that if $a = (-1)^{n/2+1} \lceil drA \rceil$ then $B_n = a/d$.
- (f) Estimate the complexity of computing B_n using the algorithm suggested by (e).
- (g) Illustrate the 3 algorithms to compute B_n considered in this problem for n=8.

Problem 3. Hecke operators on $\Gamma_1(N)$ (49 points)

Fix $N \in \mathbb{Z}_{>0}$. In lecture, we considered the Hecke algebra $\mathbb{T}(\Gamma_0(N), \Delta_0(N))$, where

$$\Gamma_0(N) := \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \operatorname{SL}_2(\mathbb{Z}) : c \equiv 0 \bmod N, a \perp N \right\},$$

$$\Delta_0(N) := \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in M_2^+(\mathbb{Z}) : c \equiv 0 \bmod N, a \perp N \right\},$$

see [1, 4.5]. We now want to consider the Hecke algebra $\mathbb{T}(\Gamma_1(N), \Delta_1(N))$, where

$$\Gamma_1(N) := \{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \operatorname{SL}_2(\mathbb{Z}) : c \equiv 0 \bmod N, a \equiv d \equiv 1 \bmod N \},$$

$$\Delta_1(N) := \{ \alpha = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in M_2^+(\mathbb{Z}) : c \equiv 0 \bmod N, a \equiv 1 \bmod N \}.$$

We also define the diamond operator $\langle d \rangle$ for each $d \in (\mathbb{Z}/N\mathbb{Z})^{\times}$, which acts on modular forms of weight k via $f \mapsto f|_k \sigma_d$, where $\sigma_d \in \mathrm{SL}_2(\mathbb{Z})$ satisfies

$$\sigma_d \equiv \begin{pmatrix} 1/d & 0 \\ 0 & d \end{pmatrix} \bmod N.$$

Throughout this problem χ denotes a Dirichlet character of modulus N, and for any integer m we write $m|N^{\infty}$ to indicate that every prime divisor of m is also a prime divisor of N.

(a) Show that

$$M_k(\Gamma_0(N), \chi) = \{ f \in M_k(\Gamma_1(N)) : \langle d \rangle f = \chi(d) f \text{ for all } d \in (\mathbb{Z}/N\mathbb{Z})^{\times} \},$$

so $\langle d \rangle$ acts on $M_k(\Gamma_0(N), \chi)$ via $f \mapsto \chi(d)f$. Conclude that $\langle d \rangle$ acts on $M_k(\Gamma_1(N))$ and this action does not depend on the choice of σ_d .

We now define Hecke operators $T(n), T(a,d) \in \mathbb{T}(\Gamma_1(N), \Delta_1(N))$, with $a|d \perp N$, via

$$T(n) := \sum_{\det \alpha = n} \Gamma_1(N) \alpha \Gamma_1(N), \qquad T(a,d) := \Gamma_1(N) \sigma_a \begin{pmatrix} a & 0 \\ 0 & d \end{pmatrix} \Gamma_1(N),$$

where the sum for T(n) is over double cosets of $\Gamma_1(N)$ in $\Delta_1(N)$.

- (b) Show that $\mathbb{T}(\Gamma_1(N), \Delta_1(N))$ is commutative and for any $\alpha \in \Delta_1(N)$ we can choose $\alpha_1, \ldots, \alpha_n \in \Delta_1(N)$ so that $\Gamma_1(N)\alpha\Gamma_1(N) = \sqcup_i\Gamma_1(N)\alpha_i = \sqcup_i\alpha_i\Gamma_1(N)$.
- (c) Prove that every Hecke operator $\Gamma_1(N)\alpha\Gamma_1(N)$ with $\alpha\in\Delta_1(N)$ can be uniquely expressed in the form

$$T(m)T(a,d) = T(a,d)T(m),$$

with $m|N^{\infty}$ and $a|d \perp N$.

- (d) Show that $\mathbb{T}(\Gamma_1(N), \Delta_1(N)) = \langle T(p), T(q,q) : p, q \text{ prime}, q \perp N \rangle$ and that we also have $\mathbb{T}_{\mathbb{Q}}(\Gamma_1(N), \Delta_1(N)) = \langle T(n) \rangle$.
- (e) For each $n \in \mathbb{Z}_{>0}$ let $\Delta^n := \{ \alpha \in \Delta_1(N) : \det \alpha = n \}$ so that $\Delta_1(N) = \sqcup_{n \geq 0} \Delta^n$. For each positive integer $a \perp N$ fix a choice of $\sigma_a \in \mathrm{SL}_2(\mathbb{Z})$. Show that for each $n \in \mathbb{Z}_{>0}$ we have

$$\Delta^{n} = \bigsqcup_{\substack{ad=n\\a \perp N}} \bigsqcup_{0 \le b < d} \Gamma_{1}(N) \sigma_{a} \begin{pmatrix} a & b\\ 0 & d \end{pmatrix}.$$

Conclude that for each $f \in M_k(\Gamma_0(N), \chi)$ we have

$$f|_k T(n) = n^{k-1} \sum_{\substack{ad=n\\a \mid N}} \sum_{0 \le b < d} \chi(a) d^{-k} f\left(\frac{az+b}{d}\right),$$

and

$$f|_k T(d,d) = d^{k-2} \chi(d) f.$$

(f) Show that when we restrict the action of $\mathbb{T}(\Gamma_1(N), \Delta_1(N))$ to $M_k(\Gamma_0(N), \chi)$ we have

$$T(m)T(n) = \sum_{d \mid \gcd(m,n)} d^{k-1}\chi(d)T(mn/d^2),$$

and derive the identity

$$\sum_{n\geq 1} T(n)n^{-s} = \prod_{p} \left(1 - T(p)p^{-s} + \chi(p)p^{k-1-2s}\right)^{-1}$$

where both of the equalities above are identities of operators on $M_k(\Gamma_0(N), \chi)$.

(g) Prove the following identity of operators on $M_k(\Gamma_1(N))$:

$$p^{k-1}\langle p\rangle = T(p)^2 - T(p^2),$$

valid for all primes p.

Problem 4. Survey (2 points)

Complete the following survey by rating each problem you attempted on a scale of 1 to 10 according to how interesting you found it (1 = "mind-numbing," 10 = "mind-blowing"), and how difficult you found it (1 = "trivial," 10 = "brutal"). Also estimate the amount of time you spent on each problem to the nearest half hour.

	Interest	Difficulty	Time Spent
Problem 1			
Problem 2			
Problem 3			

Please rate each of the following lectures that you attended, according to the quality of the material (1="useless", 10="fascinating"), the quality of the presentation (1="epic fail", 10="perfection"), the pace (1="way too slow", 10="way too fast", 5="just right") and the novelty of the material to you (1="old hat", 10="all new").

Date	Lecture Topic	Material	Presentation	Pace	Novelty
-4/2	Hecke algebras				
-4/4	Modular groups				
4/9	Hecke algebras for modular groups				
4/11	Eigenforms and L-functions				

Please feel free to record any additional comments you have on the problem sets and the lectures, in particular, ways in which they might be improved.

References

[1] T. Miyake, *Modular forms*, Springer, 2006.