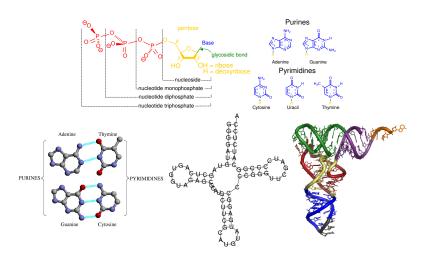
RNA Structure and RNA Structure Prediction





Definitions

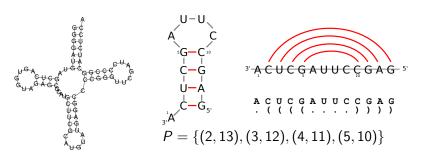
Definition (RNA Structure)

Let $S \in \{A, C, G, U\}^*$ be an RNA sequence of length n = |S|. An RNA structure of S is a set of base pairs

$$P \subseteq \{(i,j) \mid 1 \le i < j \le n, S_i \text{ and } S_j \text{ complementary}\}$$

such that the degree of P is at most one, i.e.

for all
$$(i,j), (i',j') \in P : (i = i' \Leftrightarrow j = j')$$
 and $i \neq j'$.



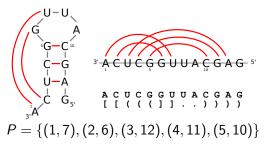
Definitions II

Definition (Crossing)

Two base pairs (i, j) and (i', j') are crossing iff

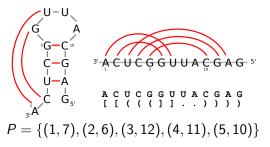
$$i < i' < j < j'$$
 or $i' < i < j' < j$.

An RNA structure P (of an arbitary RNA sequence S) is crossing iff P contains (at least) two crossing base pairs. Otherwise, P is called non-crossing or nested.



Remarks

- Synonyms: $(i,j) \in P$ is a "base pair", "bond", "arc"
- Usually, assume minimal allowed size of base pair (aka loop length) m. Then: additional constraint i - i > m in def of RNA structure.
- Crossing base pairs form "pseudoknots" crossing structures contain pseudoknots. The terms pseudoknot-free and non-crossing are synonymous for RNA structures.
- As defined "RNA structure" describes the secondary structure of an RNA. We will look at tertiary structure only later.



Prediction of RNA (Secondary) Structure

Definition (Problem of RNA non-crossing Secondary Structure Prediction by Base Pair Maximization)

IN: RNA sequence S

OUT: a non-crossing RNA structure P of S that maximizes |P| (i.e. the number of base pairs in P).

Remarks:

- By dropping the non-crossing condition, we can define the general base pair maximization problem. The general problem can be solved by maximum matching.
- Maximizing base pairs for non-crossing structures will help to understand the more realistic case of minimizing energy. For ernergy minimization, predicting general structures is NP-hard.
- RNA structure prediction is often (less precisely) called RNA folding.

Nussinov Algorithm — Matrix definition

Let S be and RNA sequence of length n.

The Nussinov Algorithm solves the problem of RNA non-crossing secondary structure prediction by base pair maximization with input S.

Definition (Nussinov Matrix)

The *Nussinov matrix* $N = (N_{ij}) \underset{i-1 \le j \le n}{1 \le i \le n}$ of S is defined by

$$N_{ij} := \max\{|P| \mid P \text{ is non-crossing RNA } ij\text{-substructure of } S\}$$

where we use:

Definition (RNA Substructure)

An RNA structure P of S is called *ij-substructure of* S iff $P \subseteq \{i, \ldots, j\}^2$.



Nussinov Algorithm — Recursive computation of $N_{i,j}$

Init: (for $1 \le i \le n$)

$$N_{ii} = 0$$
 and $N_{ii-1} = 0$

Recursion: (for $1 \le i < j \le n$)

$$N_{ij} = \max egin{cases} N_{ij-1} \ \max_{\substack{i \leq k < j \ S_k, S_j complementary}} N_{ik-1} + N_{k+1j-1} + 1 \end{cases}$$

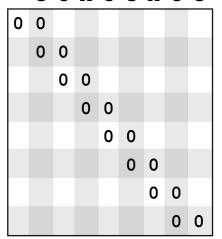
Remarks:

- case 2 of recursion covers base pair (i,j) for k=i; then: N_{ik-1} (initialized with 0!) is max. number of base pairs in empty sequence.
- solution is in $N_{1,n}$
- Recursion furnishs a DP-Algorithm for computing the Nussinov matrix (including $N_{1,n}$) in $O(n^3)$ time and $O(n^2)$ space.
- How to guarantee minimal loop length?
- What happens without restriction non-crossing?
 - Are there other decompositions?



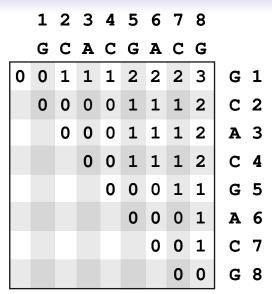
Nussinov Algorithm — Example

1 2 3 4 5 6 7 8 G C A C G A C G





Nussinov Algorithm — Example





Nussinov Algorithm — Traceback

Determine one non-crossing RNA structure P with maximal |P|.

pre: Nussinov matrix N of S:

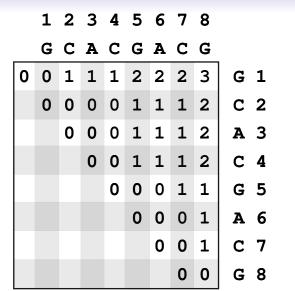


Idea:

- start with entry at upper right corner N_{1n}
- determine recursion case (and the entries in N) that yield maximum for this entry
- trace back the entries where we recursed to

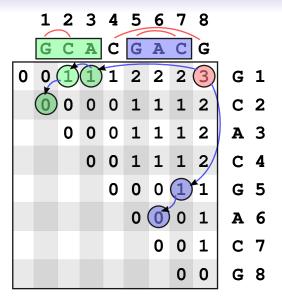


Nussinov Algorithm — Traceback Example





Nussinov Algorithm — Traceback Example





Nussinov Algorithm — Traceback Pseudo-Code

```
CALL: traceback(1, n)
Procedure traceback(i, j)
  if i < i then
     return
  else if N_{ii} = N_{ii-1} then
    traceback(i, i-1):
     return
  else
    for all k : i \le k < j, S_k and S_i complementary do
       if N_{ii} = N_{i,k-1} + N_{k+1,i-1} + 1 then
          print (k,j);
         traceback(i, k-1); traceback(k+1, j-1);
          return
       end if
     end for
  end if
```

Remarks

- Complexity of trace-back $O(n^2)$ time
- How to get all optimal non-crossing structures?
- How to trace-back non-recursively?
- How to output / represent structures?
 - Dot-bracket
 - 2D-layout
 - Tree-like

Limitations of the Nussinov Algorithm

- Base pair maximization does not yield biologically relevant structures:
 - no stacking of base pairs considered
 - loop sizes not distinguished
 - no special scoring of multi-loops
- only one structure predicted
 - base pair maximization can not differnciate structures sufficiently well: possibly many optima
 - no sub-optimal solutions
- crossing structures cannot be predicted

However:

- shows pattern of RNA structure prediction by DP (simple+instructive)
- energy minimization (Zuker) will have similar algorithmic structure
- "only one solution"-problem can be overcome (suboptimal: Wuchty
- prediction of (restricted) crossing structure can be seen as extension



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