

RSK (cont'd)Insertion Algorithm

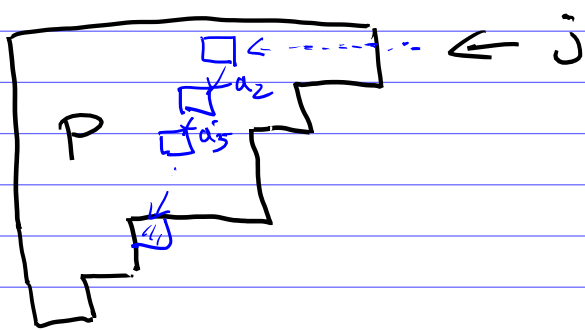
P - a semi-standard Young tableau

j - positive integer

$P \leftarrow j$: the SSYT w/ 1 move box, obtained as follows:

Algorithm

- (1) $r := 1$ (row index), $a_r := j$
- (2) IF $a_r \geq$ largest entry of row r ,
then add 1 box in the end
of row r filled with a_r ,
and STOP.
- (3) OTHERWISE, find
the smallest entry a_{r+1} of
row r s.t. $a_{r+1} > a_r$.
Replace a_{r+1} with a_r
- (4) Set $r := r + 1$
- (5) GO TO (2)



Ex.

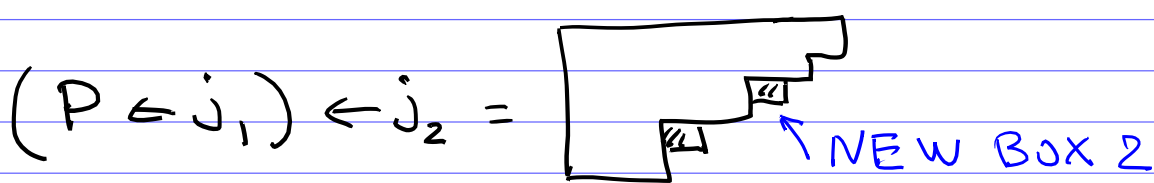
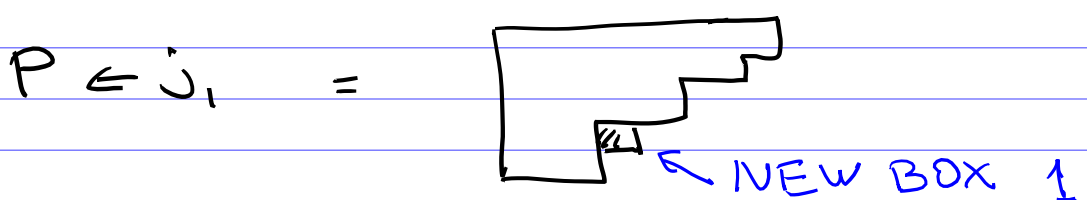
1	1	1	2	3	3	5
2	2	2	3	4	5	
3	4	4	6	7	7	
5	5					
6	7					

$j=2$

$P \leftarrow j$:

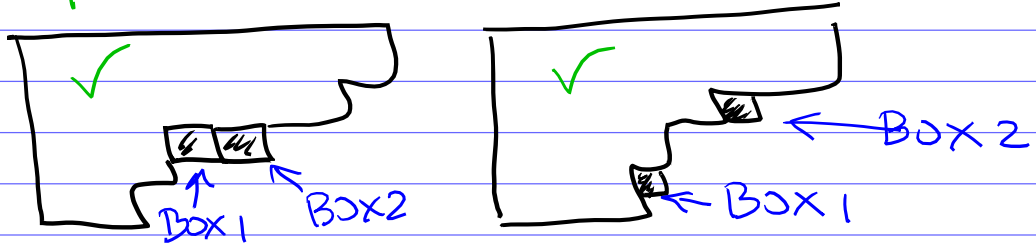
1	1	1	2	2	3	5
2	2	2	3	3	5	
3	4	4	4	7	7	
5	5	6				
6	7					

Lemma. If $j_1 \leq j_2$

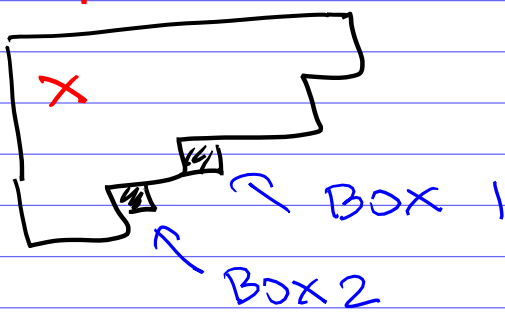


Then NEW BOX 2 is located immediately to the right, or above of NEW BOX 1

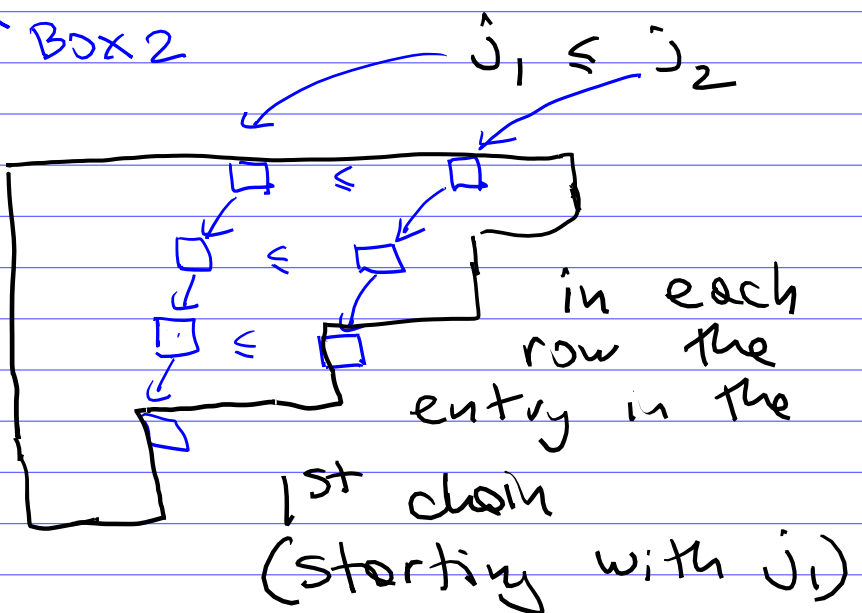
possible:



impossible:



Proof



is strictly to the left and less than or equal to the entry in the 2nd chain (starting with j_2).

\Rightarrow The second chain ends before or in the same row as the first chain. \square

RSK : $A = (a_{ij}) \rightsquigarrow (P, Q)$

Convert A into biword

pair of SSYTs
of same
shape λ

$\begin{pmatrix} i_1 \\ j_1 \end{pmatrix} \begin{pmatrix} i_2 \\ j_2 \end{pmatrix} \dots \begin{pmatrix} i_n \\ j_n \end{pmatrix}$ with

a_{ij} pairs $\begin{pmatrix} i \\ j \end{pmatrix}$ ordered

lexicographically: $\begin{pmatrix} i \\ j \end{pmatrix} < \begin{pmatrix} i' \\ j' \end{pmatrix}$ if

$i < i'$ or
 $(i = i' \ \& \ j < j')$

insertion tableau:

$P = ((\emptyset \leftarrow j_1) \leftarrow j_2) \leftarrow \dots \leftarrow j_n$

recording tableau:

Q : if at k -th insertion
we added a box to P ,
then we add a box to
 Q located at the same
position, filled with i_k .

Ex $A = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$

biword $w = \begin{pmatrix} 1 \\ 2 \end{pmatrix} \begin{pmatrix} 2 \\ 1 \end{pmatrix} \begin{pmatrix} 2 \\ 1 \end{pmatrix} \begin{pmatrix} 2 \\ 4 \end{pmatrix} \begin{pmatrix} 3 \\ 3 \end{pmatrix}$

P: $\emptyset \xrightarrow{2} \boxed{2} \xrightarrow{1} \begin{array}{|c|} \hline 1 \\ \hline 2 \\ \hline \end{array} \xrightarrow{1} \begin{array}{|c|c|} \hline 1 & 1 \\ \hline 2 & \\ \hline \end{array} \xrightarrow{4} \begin{array}{|c|c|c|} \hline 1 & 1 & 4 \\ \hline 2 & & \\ \hline \end{array} \xrightarrow{3} \begin{array}{|c|c|c|} \hline 1 & 1 & 3 \\ \hline 2 & 4 & \\ \hline \end{array}$

Q: $\emptyset \xrightarrow{1} \boxed{1} \xrightarrow{2} \begin{array}{|c|} \hline 1 \\ \hline 2 \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 2 & \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|c|} \hline 1 & 2 & 2 \\ \hline 2 & & \\ \hline \end{array} \xrightarrow{2} \begin{array}{|c|c|c|} \hline 1 & 2 & 2 \\ \hline 2 & 3 & \\ \hline \end{array}$

$\begin{pmatrix} 0 & 1 & 0 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \xrightarrow{\text{RSK}} \left(\begin{array}{|c|c|c|} \hline 1 & 1 & 3 \\ \hline 2 & 4 & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 1 & 2 & 2 \\ \hline 2 & 3 & \\ \hline \end{array} \right)$

A

P

Q

weights: $(2, 1, 1, 1)$ $(1, 3, 1)$

column sums
of A

row sums
of A

Clearly, $A \rightsquigarrow$ pair of SSYT's
of the same shape

In order to show that
this is a bijection, we
need to construct inverse
procedure.

Inverse Algorithm:

$$(P, Q) \mapsto \text{biword} \rightsquigarrow A$$

(1) Find the horizontal strip in Q filled with largest entries of Q

(2) "Uninsert" entries of P tableau located in the same positions as in the strip of Q (found in (1)) in the order right to left

this follow from Lemma

(3) Repeat steps (1) & (2) until the tableaux are empty.

Ex.

$$P = \begin{array}{|c|c|c|} \hline 1 & 1 & 3 \\ \hline 2 & 4 & \\ \hline \end{array} \quad Q = \begin{array}{|c|c|c|} \hline 1 & 2 & 2 \\ \hline 2 & 3 & \\ \hline \end{array}$$

the largest entry in previous row which is < 4

$$\left(\begin{array}{|c|c|c|} \hline 1 & 1 & 4 \\ \hline 2 & & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 1 & 2 & 2 \\ \hline 2 & & \\ \hline \end{array} \right)$$

need to uninsert entries in these positions from right to left

$$\left(\begin{array}{|c|} \hline 2 \\ \hline \end{array}, \begin{array}{|c|} \hline 1 \\ \hline \end{array} \right)$$

bi word

$$\begin{pmatrix} 1 \\ 2 \end{pmatrix} \begin{pmatrix} 2 \\ 1 \end{pmatrix} \begin{pmatrix} 2 \\ 1 \end{pmatrix} \begin{pmatrix} 2 \\ 4 \end{pmatrix} \begin{pmatrix} 3 \\ 3 \end{pmatrix}$$

matrix

$$A = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

We proved

Theorem RSK is a bijection
between matrices A &
pairs (P, Q) of SSYT's
of some shape λ s.t.
column/row sums of A
are weights of P & Q .

This proved Cauchy
identity

$$\prod_{i,j} \frac{1}{1-x_i y_j} = \sum_{\lambda} s_{\lambda}(x) s_{\lambda}(y)$$

How about dual Cauchy?

$$\prod_{i,j} (1+x_i y_j) = \sum_{\lambda} s_{\lambda}(x) s_{\lambda'}(y)$$

There is a similar
construction, called the
dual RSK that gives
a bijection:

A $\xleftrightarrow{\text{dual RSK}}$ (P, Q)
matrix filled
with 0's & 1's pairs of SSYT's
of conjugate
shapes
shape $(P) = \lambda$
shape $(Q) = \lambda'$

column/row sums of A
= weights of P & Q .

RSK has many other important properties.

Theorem If $A \xrightarrow{\text{RSK}} (P, Q)$

then $A^T \xrightarrow{\text{RSK}} (Q, P)$.

Can be proved. But not obvious from the classical construction of RSK.

We'll give another construction where this symmetry is manifest.

Theorem $A \xrightarrow{\text{RSK}} (P, Q)$ of shape λ .

Then $\lambda_1 =$ the length of a "longest weakly increasing" subsequence of w

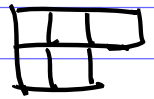
length of 1st column of λ

$\lambda'_1 =$ the length of a "longest strictly decreasing" subseq. of w

Ex. $A = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \xrightarrow{\text{RSK}} (P, Q)$

$w = (1 \ 2 \ 2 \ 2 \ 3)$
 $(2 \ \boxed{1} \ \boxed{1} \ 4 \ \boxed{3})$

of shape



$\lambda_1 = 3$

$\lambda'_1 = 2$

Note:

Subsequence does not need to be consecutive

a longest weakly incr. subseq.

Problem. What is the length of a longest increasing subsequence in a permutation (distribution of the lengths, etc)

This was one of the reasons why Robinson-Schensted was initially introduced.

Let's talk about permutations, (Schensted's case)

We have $S_n \ni w \xleftrightarrow{\text{Schubert}} (P, Q)$

pair of SYTs
of same
shape λ
with $|\lambda| = n$

$w = \begin{pmatrix} 1 & 2 & \dots & n \\ w_1 & w_2 & \dots & w_n \end{pmatrix}$ permutation

λ_1 : longest incr. subseq

$$w_{i_1} < w_{i_2} < \dots < w_{i_{\lambda_1}}$$

$$i_1 < \dots < i_{\lambda_1}$$

λ'_1 : longest decreasing
subseq.

$$w_{j_1} > w_{j_2} > \dots > w_{j_{\lambda'_1}}$$

$$j_1 < j_2 < \dots < j_{\lambda'_1}$$

Symmetry If $w \leftrightarrow (P, Q)$

then $w^{-1} \leftrightarrow (Q, P)$

Let $f_\lambda := \#\{\text{SYT's of shape } \lambda\}$

Corollary

$$(1) \sum_{\lambda \vdash n} f_\lambda^2 = n!$$

later we'll give explicit formula for f_λ called the hook lengths formula

$$(2) \sum_{\lambda \vdash n} f_\lambda = \#\{\text{involutions in } S_n\} \\ \equiv \#\{w \in S_n \mid w^{-1} = w\}$$

Exercise
(will be in P set 1)

$$\equiv \sum_{k=0, \dots, \lfloor \frac{n}{2} \rfloor} \binom{n}{2k} (2k-1)!!$$

where $(2k-1)!! = 1 \cdot 3 \cdot 5 \cdots (2k-1)$.

Proof of (2)

$$w \mapsto (P, Q), \quad w^{-1} \mapsto (Q, P)$$

$P=Q \iff w = w^{-1}$, i.e. w is an involution in S_n

□

Applications to Ramsey Theory

Corollary (Erdős-Szekeres Theorem)

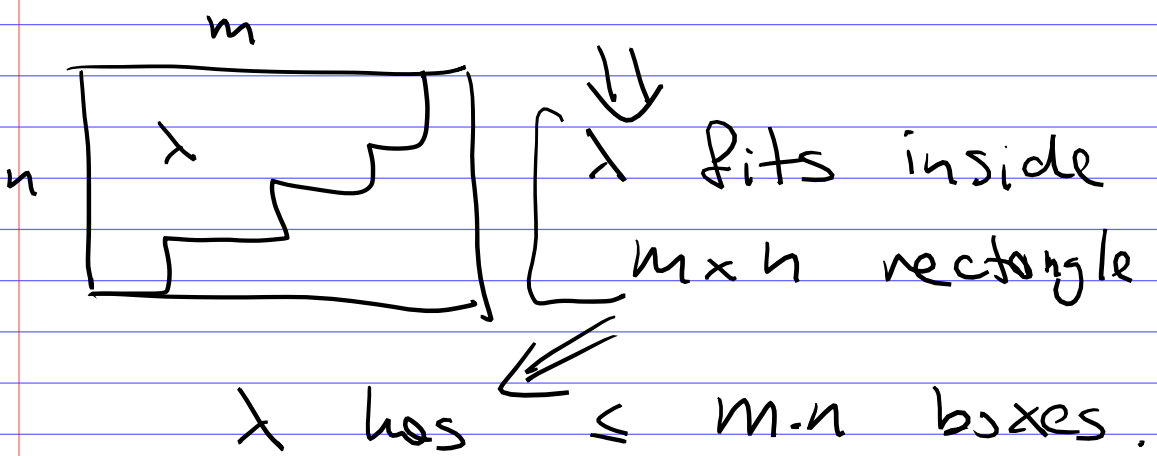
Fix $m, n \geq 1$, $N = m \cdot n + 1$.

Any permutation w_1, w_2, \dots, w_N either has an increasing subsequence of length $m+1$ or a decreasing subseq. of length $n+1$.

Proof. Suppose not:
longest incr. subseq. has length $\leq m$
& longest decr. subseq. $-1- \leq n$

$w \xrightarrow{\text{Schensted}} (P, Q)$ of shape λ

$$\lambda_1 \leq m \quad \& \quad \lambda'_1 \leq n$$



But we assumed that

$$N = |\lambda| = m \cdot n + 1. \quad \square$$

There is another nice argument for Erdős-Szekeres Thm based on pigeonhole principle.

$\lambda_1 \leftrightarrow$ increasing subseq.

How about $\lambda_2, \lambda_3, \dots$?

Greene's Theorem

$w \xrightarrow{\text{RSK}} (P, Q)$ of shape λ
 \uparrow (a permutation or a biword)

Then $\forall k$

$$\lambda_1 + \dots + \lambda_k = \max \# \left(\begin{array}{l} \text{union of} \\ k \text{ weakly} \\ \text{increasing} \\ \text{subseq. of } w \end{array} \right)$$

$$= \max \# (I_1 \cup \dots \cup I_k)$$

each I_1, \dots, I_k is increasing

$$\lambda'_1 + \dots + \lambda'_k = \max \# \left(\begin{array}{l} \text{union of} \\ k \text{ strictly} \\ \text{decr. subseq.} \end{array} \right)$$

$$= \max \# (D_1 \cup \dots \cup D_k)$$

each D_1, \dots, D_k is decreasing

Example. $w = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 3 & 5 & 1 & 8 & 7 & 2 & 6 & 4 \end{pmatrix}$

shape of P & Q ?

$$w = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 3 & 5 & 1 & 8 & 7 & 2 & 6 & 4 \end{pmatrix}$$

$$\lambda_1 = 3$$

$$\lambda_1 + \lambda_2 = 6$$

$$\lambda_1 + \lambda_2 + \lambda_3 = 7$$

$$\lambda_1 + \lambda_2 + \lambda_3 + \lambda_4 = 8$$

$$\lambda = \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & & \\ \hline \end{array}$$

$$w = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 3 & 5 & 1 & 8 & 7 & 2 & 6 & 4 \end{pmatrix}$$

$$\lambda'_1 = 4$$

$$\lambda'_1 + \lambda'_2 = 6$$

$$\lambda'_1 + \lambda'_2 + \lambda'_3 = 8$$

$$\lambda' = \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & \square & \square & \\ \hline \square & \square & & \\ \hline \end{array}$$

Warning: $\lambda_1 + \lambda_2 = \max \# (I_1 \cup I_2)$

I_1, I_2 are incr. subseq.

but sometimes it is impossible

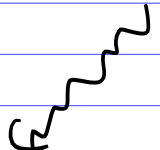
to find such I_1 & I_2 s.t.

I_1 is a max. incr. subseq.

How to prove these
properties of RSK
(Symmetry, incr. & decr.
subsequences, ...)?

There is another construction
of RSK where these
properties become clear.

We want to break
Schensted insertion steps
into smaller, "more
elementary" steps.

this leads to Fomin's
growth diagrams (for SYTs
&
toggle operations (for SSYTs
Kirillov - Berenstein, & ...)

Let's start with
permutations & SYT's.

$$S_n \xleftrightarrow{\text{Schensted}} \{ (P, Q) \}$$

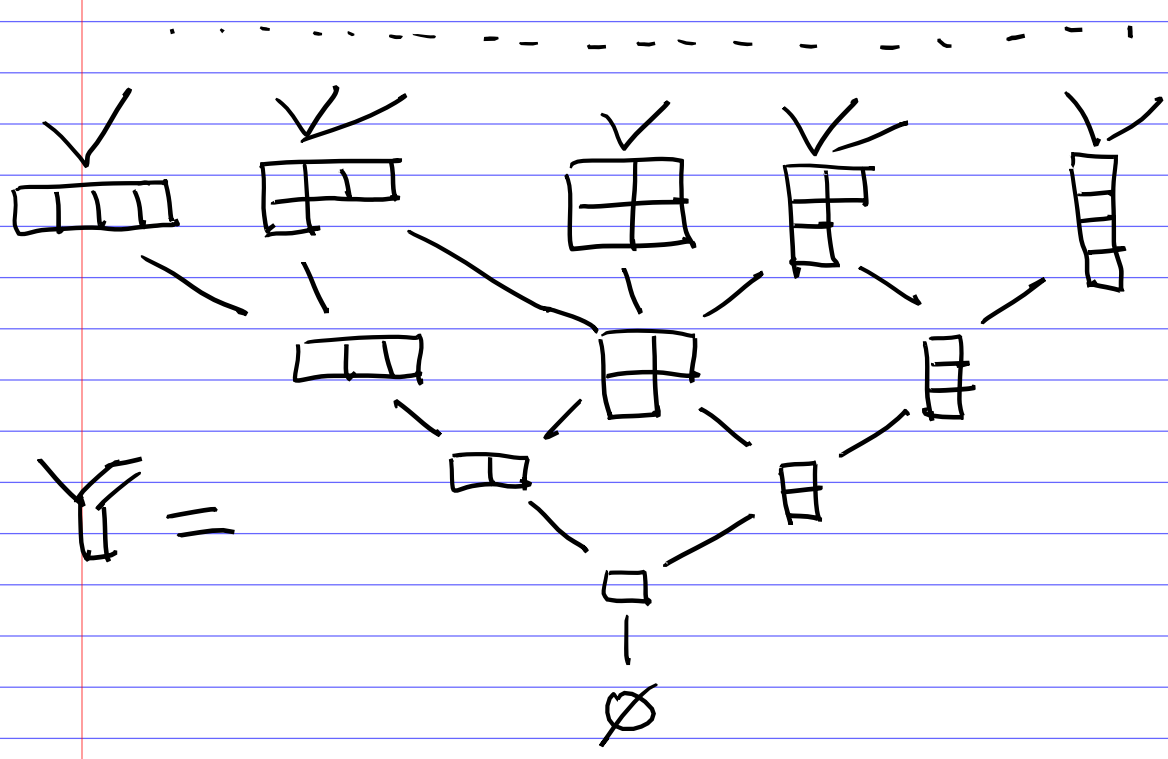
pair of SYT's
of shape $\lambda \vdash n$.

Young's lattice \mathcal{Y}

All Young diagrams ordered
by inclusion $\lambda \leq \mu$ if
 $\lambda \subseteq \mu$.

Covering relations in \mathcal{Y}

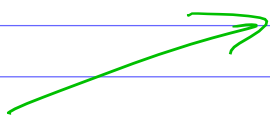
$$\lambda < \mu \text{ if } \lambda \subseteq \mu \text{ \& } |\mu/\lambda| = 1$$



An SYT corresponds to a saturated increasing chain is \mathbb{Y} from \emptyset to λ

Ex. $\emptyset < \square < \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array} < \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array} < \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \end{array} < \begin{array}{|c|c|c|c|c|} \hline \square & \square & \square & \square & \square \\ \hline \end{array}$

1	3
2	4
5	

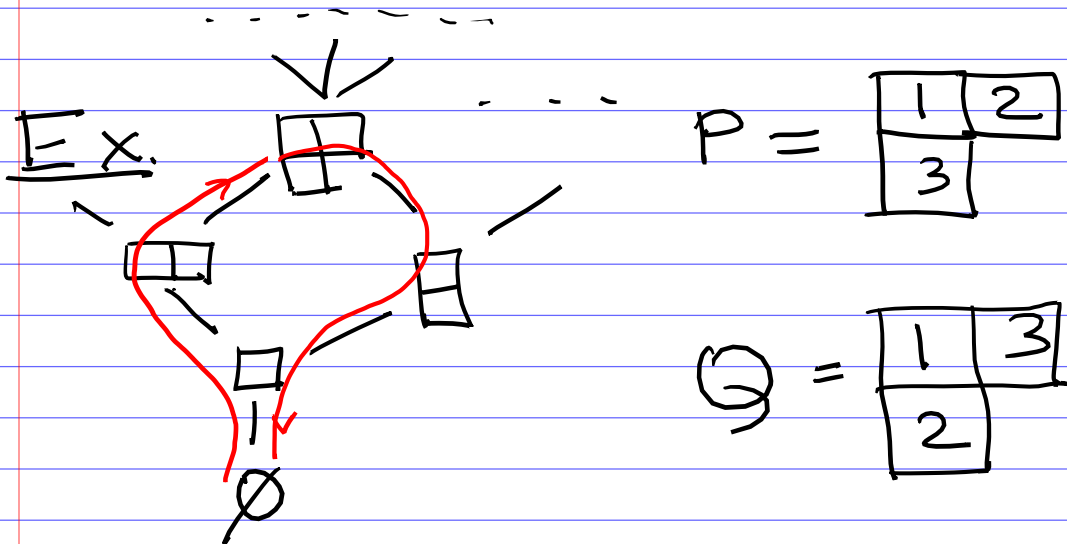


A pair (P, Q) of SYT's of the same shape λ

corresponds to a path in \mathbb{Y}

$\emptyset < \square < \dots < \lambda > \dots > \emptyset$

from \emptyset to \emptyset with
 n "UP" steps followed by
 n "DOWN" steps



How about arbitrary paths
in \mathbb{Y} from \emptyset to \emptyset ?

Up & Down operators

acting on $\mathbb{Z}[\mathbb{Y}]$
(the space of formal linear
combinations of Young diagrams)

$\mathbb{Z}[\mathbb{Y}]$ has linear basis
given by λ 's

Some elts of $\mathbb{Z}[\mathbb{Y}]$:

$$\square, \quad 2\square + 25 \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array} - \begin{array}{|c|} \hline \square \\ \hline \end{array}$$

$$3 \cdot \emptyset + 100 \cdot \square, \text{ etc.}$$

$U, D : \mathbb{Z}[\mathbb{Y}] \rightarrow \mathbb{Z}[\mathbb{Y}]$
two linear operators s.t.

$$U : \lambda \mapsto \sum_{\mu : \mu \triangleright \lambda} \mu$$

(μ obtained from
 λ by adding a
single box)

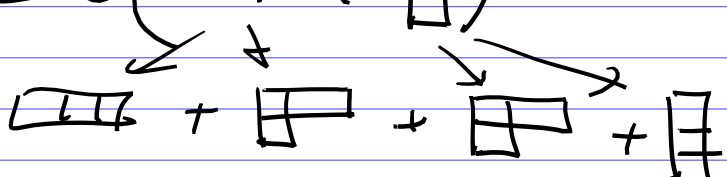
$$D : \lambda \mapsto \sum_{\mu : \mu \triangleleft \lambda} \mu$$

μ is obtained from
 λ by removing a
single box.

$$\underline{\text{Ex.}} \quad U: \square \mapsto \square\square + \begin{array}{|c|} \hline \square \\ \hline \end{array}$$

$$D: \square \mapsto \emptyset$$

$$U^2(\emptyset) = U(\square) = \square\square + \begin{array}{|c|} \hline \square \\ \hline \end{array}$$

$$U^3(\emptyset) = U(\square\square + \begin{array}{|c|} \hline \square \\ \hline \end{array}) =$$


$$= \square\square + 2 \begin{array}{|c|} \hline \square \\ \hline \end{array} + \begin{array}{|c|} \hline \square \\ \hline \end{array}$$

etc

$$U^n(\emptyset) = \sum_{\lambda \vdash n} f_\lambda \lambda$$

$$D^3 U^3(\emptyset) = D^3(\square\square + 2 \begin{array}{|c|} \hline \square \\ \hline \end{array} + \begin{array}{|c|} \hline \square \\ \hline \end{array})$$

$$= \emptyset + 2(2 \cdot \emptyset) + \emptyset$$

$$= 6 \emptyset.$$

How about

$$D^n U^n(\emptyset) = ?$$

Theorem

$$D^n U^n(\emptyset) = n! \emptyset.$$

This is equivalent to

$$\sum_{\lambda \vdash n} f_{\lambda}^2 = n!$$



There are $n!$ of paths in \mathbb{Y} from \emptyset to \emptyset with n "Up" steps followed by n "Down" steps.

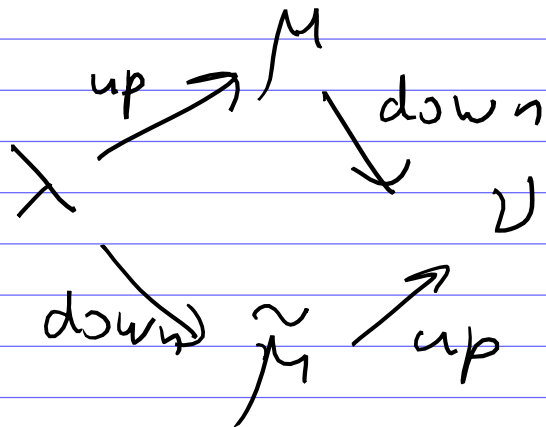
We already know this from RSK.

But is there a simpler proof?

Lemma $DU - UD = I$

or $[D, U] = I$

Proof



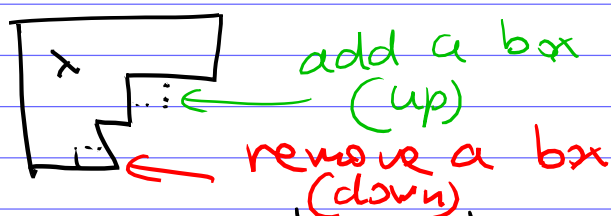
The coeff. of ν in

$$(DU - UD)(\lambda) =$$

$$\# \text{ such } \mu\text{'s} - \# \text{ such } \tilde{\mu}\text{'s}$$

2 cases:

(I) $\lambda \neq \nu$



we want to add a box to λ & then remove a different box.

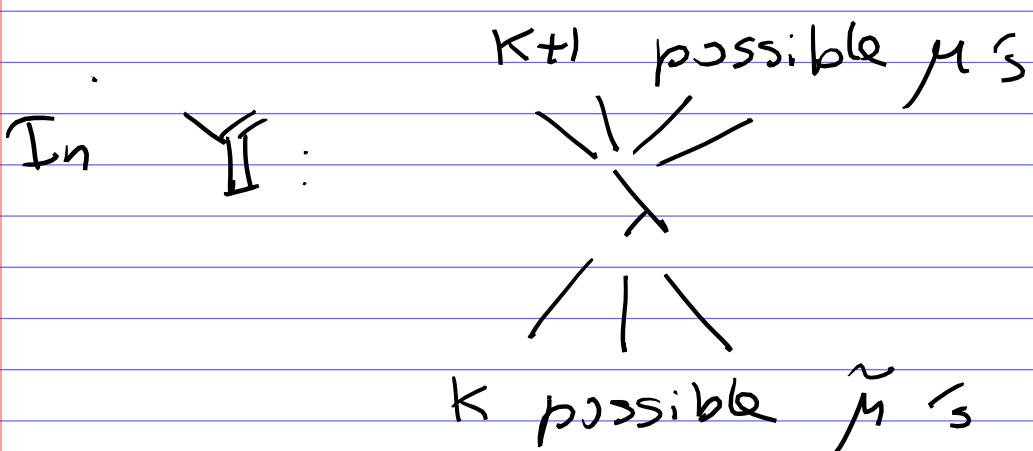
We can do these operations in a different order (first remove & then add)

$$\text{So } [\nu] (DU - UD)(\lambda) = 0 \quad \text{if } \nu \neq \lambda.$$

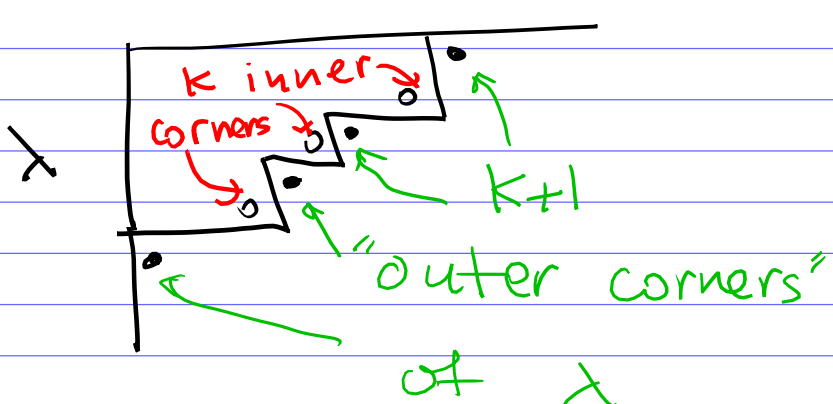
the coeff. of ν

Case II $\lambda = \nu$

$$\# \left\{ \begin{array}{l} \text{Add a box to } \lambda \text{ \& then} \\ \text{remove the same box?} \end{array} \right\} - \# \left\{ \begin{array}{l} \text{Remove a box from } \lambda \\ \text{\& then add the same box?} \end{array} \right\}$$



Why?



$$\text{So } [\lambda] (DU - UD)(\lambda) = (k+1) - k = 1$$

We obtain $(DU - UD)(\lambda) = \lambda.$

Claim the relations

$$(1) [D, U] = I$$

$$(2) D \emptyset = 0$$

Imply that $D^n U^n (\emptyset) = n! \emptyset$

Ex $DDUU(\emptyset) =$

$$= (DU \overset{\curvearrowright}{DU} + \overset{\curvearrowright}{DU}) (\emptyset)$$

$$= \left(\underset{=0}{DU} \underset{=0}{UD} + \overset{\curvearrowright}{DU} + \underset{=0}{U} \overset{\curvearrowright}{D} + I \right) \emptyset$$

$$= \left(\underset{=0}{UD} + I + I \right) (\emptyset) =$$

$$= 2I \quad \left(\begin{array}{l} \text{each } D \text{ is} \\ \text{"moving" to the right} \\ \text{until it "annihilates"} \\ \text{with some } U \end{array} \right)$$

There are $n!$ ways to match all n D 's with n U 's into "annihilating pairs".

Another argument

U, D satisfy the same relations as the operators

$$X : f(x) \mapsto x \cdot f(x)$$

$$\frac{d}{dx} : f(x) \mapsto f'(x)$$

acting on polynomial ring $\mathbb{Z}[x]$

$$\text{and } \frac{d}{dx}(1) = 0.$$

So the calculation

$$\text{of } D^n U^n(\varnothing) = ? \quad \varnothing$$

is equivalent to

$$\left(\frac{d}{dx}\right)^n X^n(1) = \left(\frac{d}{dx}\right)^n (x^n)$$

$$= n! \cdot 1 \Rightarrow ? = n!$$

□

This is why Stanley gave the following definition

Def. A poset P is called a differential poset if

(1) P has a unique minimal element $\hat{0}$.

(2) The "up" & "down" operators acting on $\mathbb{Z}[P]$ satisfy

$$[D, U] = I$$

Theorem. For any differential poset

$$D^n U^n(\hat{0}) = n! \hat{0}$$