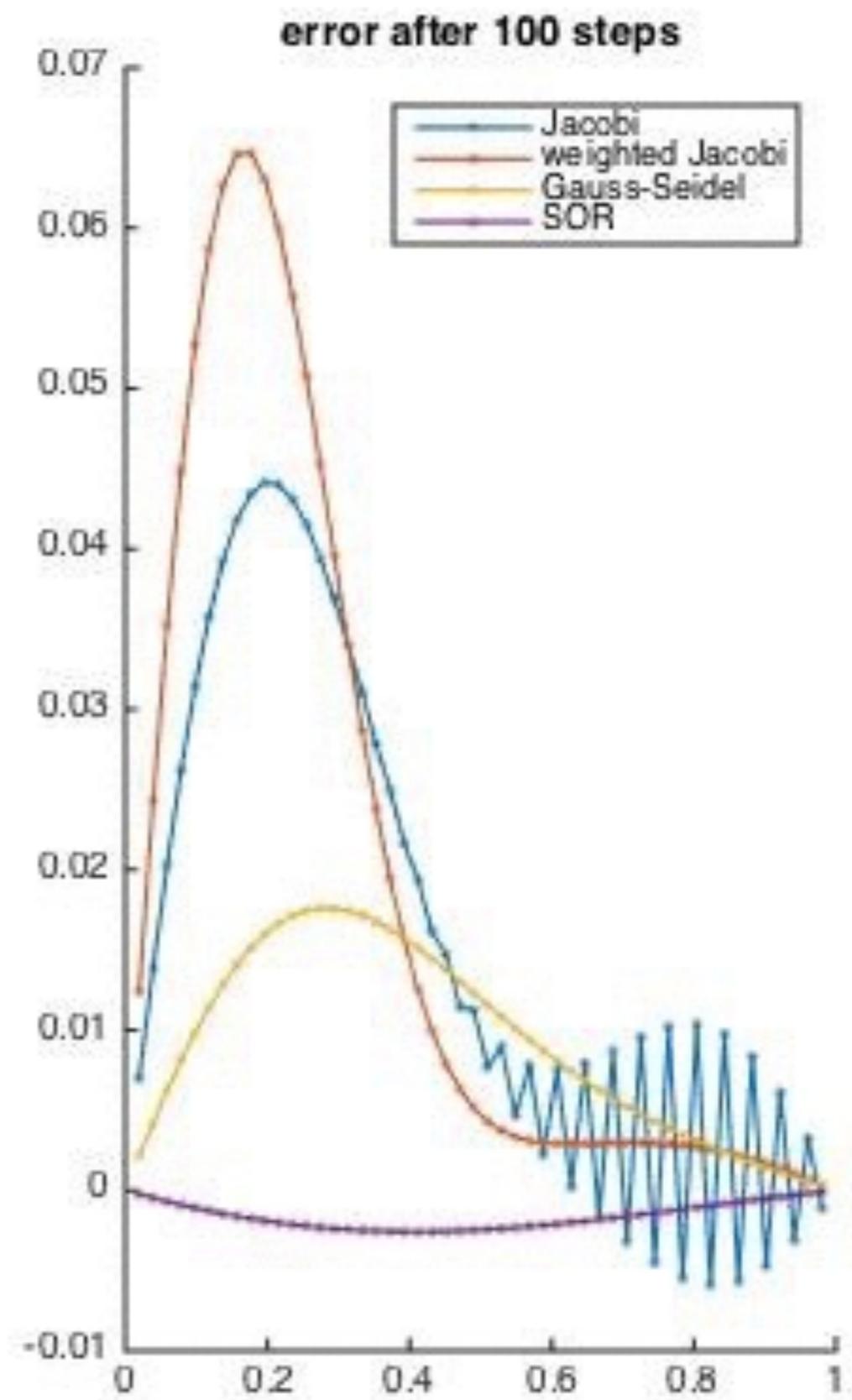
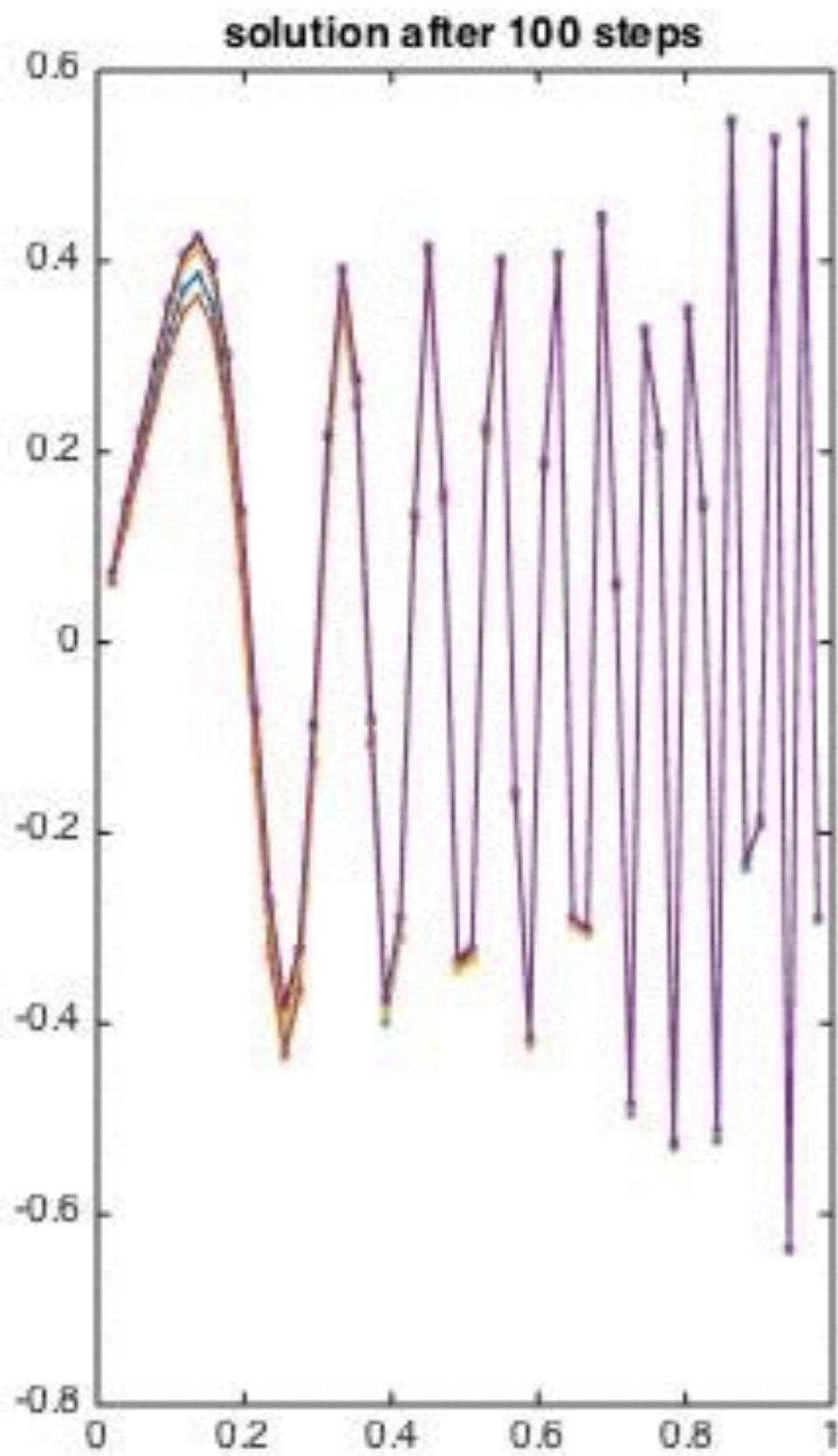


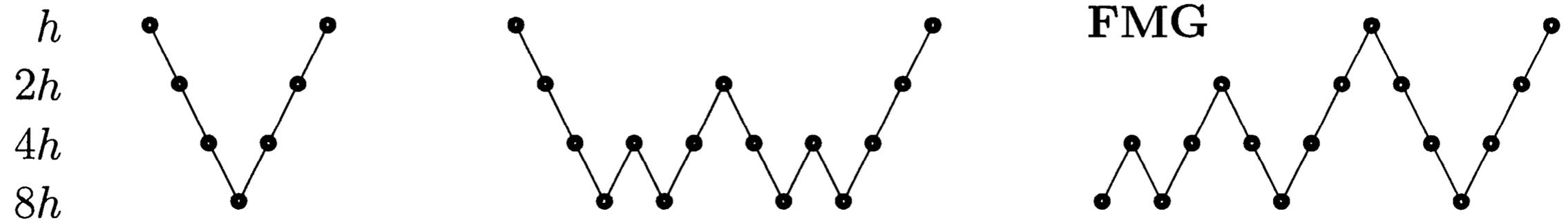
# Lecture 16

18.086



# Multigrid methods (7.3)

- Low freq. on fine mesh => high frequency on coarse mesh
- More than two meshes possible!
- “Notation”:



- Simplest one is the v-cycle with just two meshes ( $h$  and  $2h$  grid spacing)

# v-cycle algorithm

1. **Iterate** on  $A_h u = b_h$  to reach  $u_h$  (say 3 Jacobi or Gauss-Seidel steps).
2. **Restrict** the residual  $r_h = b_h - A_h u_h$  to the coarse grid by  $r_{2h} = R_h^{2h} r_h$ .
3. **Solve**  $A_{2h} E_{2h} = r_{2h}$  (or come close to  $E_{2h}$  by 3 iterations from  $E = 0$ ).
4. **Interpolate**  $E_{2h}$  back to  $E_h = I_{2h}^h E_{2h}$ . Add  $E_h$  to  $u_h$ .
5. **Iterate** 3 more times on  $A_h u = b_h$  starting from the improved  $u_h + E_h$ .

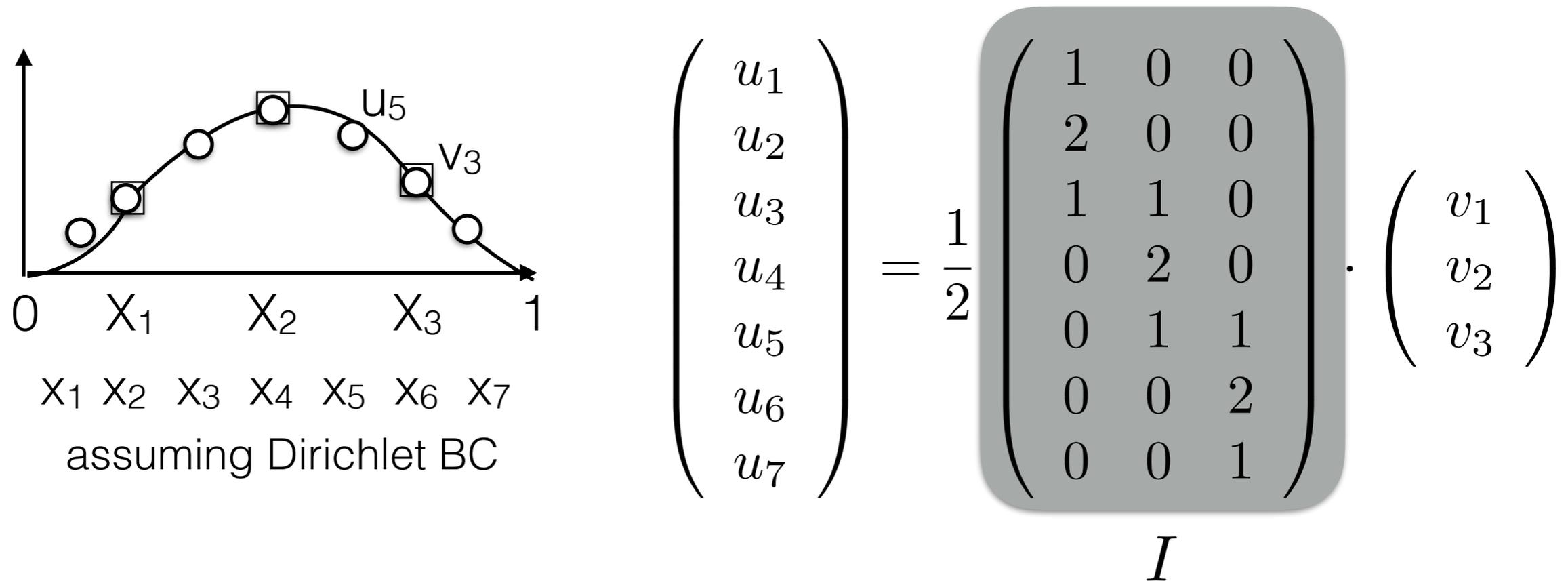
In the following:

$u_i$ : quantities on fine grid

$v_i$ : quantities on coarse grid

# Technicalities of multigrain algorithms

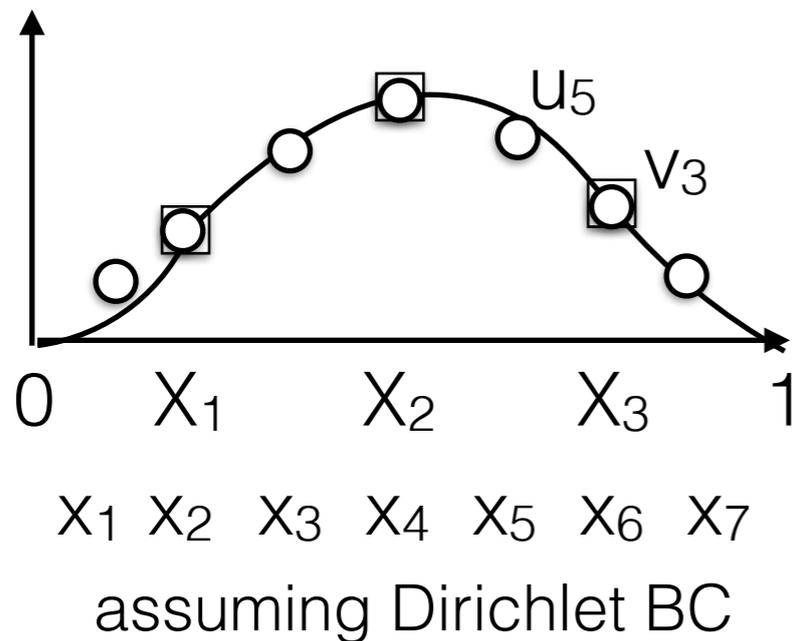
- Interpolation matrix  $I$  from coarse ( $v_i$ )  $\Rightarrow$  fine ( $u_i$ ) grid: linear interpolation



- 2D (see blackboard)

# Technicalities of multigrain algorithms

- Restriction matrix  $R$  from fine  $(u_i) \Rightarrow$  coarse  $(v_i)$  grid:



A possible choice:  $v_1 = u_2$  etc.

- Smarter choice: Weighted average, i.e.  $v_1 = (u_1 + 2u_2 + u_3)/4$ . Then:

$$R = \frac{1}{2} I^T$$

# Technicalities of multigrain algorithms

- Restriction of system matrix  $A$  from fine ( $A_h$ )  $\Rightarrow$  coarse ( $A_{2h}$ ) grid:

$$A_{2h} = RA_hI$$

- Example using  $A=K_5/h^2$

$$A_h = \frac{1}{h^2} \begin{pmatrix} 2 & -1 & 0 & \dots \\ -1 & 2 & -1 & \dots \\ \dots & \dots & \dots & \dots \\ 0 & \dots & -1 & 2 \end{pmatrix} \begin{array}{c} \updownarrow \\ 5 \end{array}$$

5 fine grid points  $\Rightarrow$   
2 coarse grid points

$$A_{2h} = RA_hI = \dots = \frac{1}{(2h)^2} \begin{pmatrix} 2 & -1 \\ -1 & 2 \end{pmatrix}$$

This is just the K-  
matrix on the coarse  
mesh!

# v-cycle algorithm

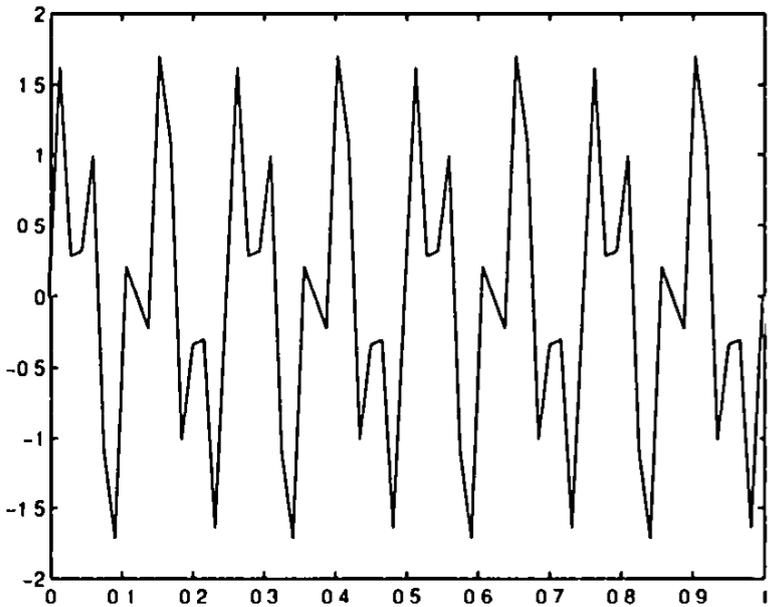
1. **Iterate** on  $A_h u = b_h$  to reach  $u_h$  (say 3 Jacobi or Gauss-Seidel steps).
2. **Restrict** the residual  $r_h = b_h - A_h u_h$  to the coarse grid by  $r_{2h} = R_h^{2h} r_h$ .
3. **Solve**  $A_{2h} E_{2h} = r_{2h}$  (or come close to  $E_{2h}$  by 3 iterations from  $E = 0$ ).
4. **Interpolate**  $E_{2h}$  back to  $E_h = I_{2h}^h E_{2h}$ . Add  $E_h$  to  $u_h$ .
5. **Iterate** 3 more times on  $A_h u = b_h$  starting from the improved  $u_h + E_h$ .

In the following:

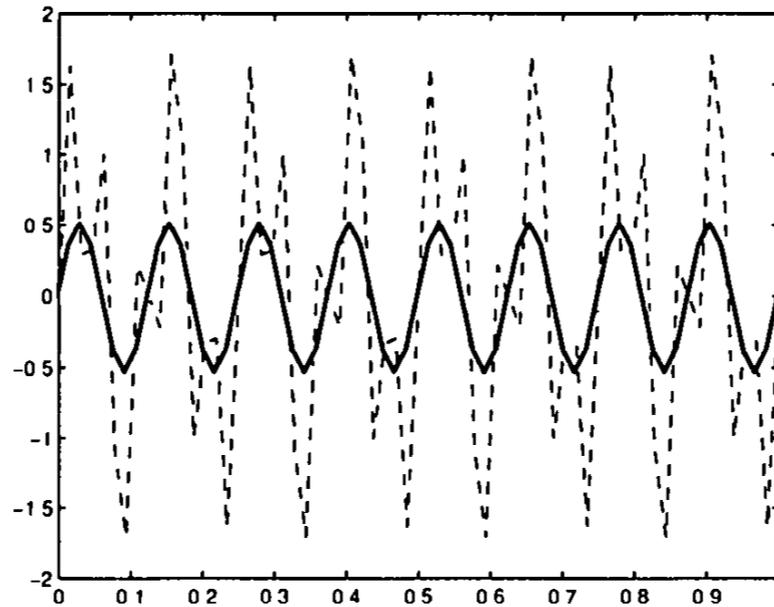
$u_i$ : quantities on fine grid

$v_i$ : quantities on coarse grid

# Error behavior

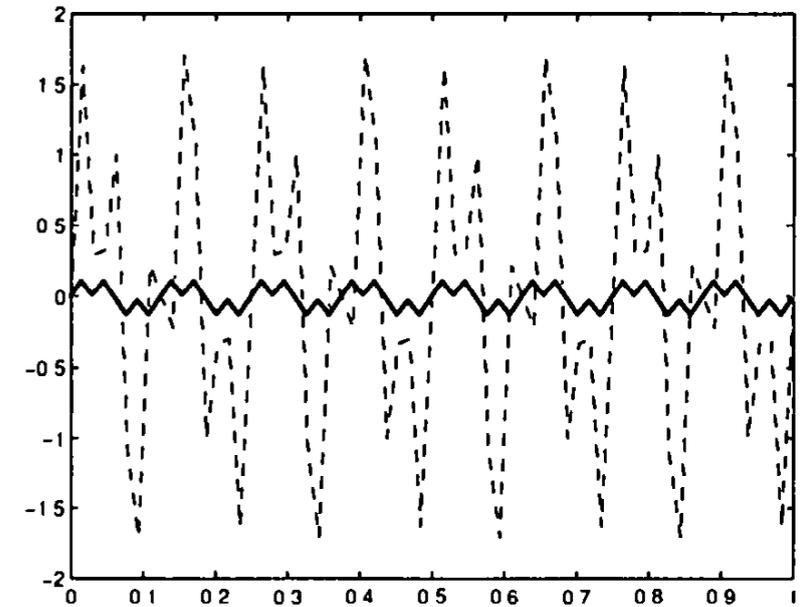


error of initial guess



error  $e_h$

after 3 fine grid iter.



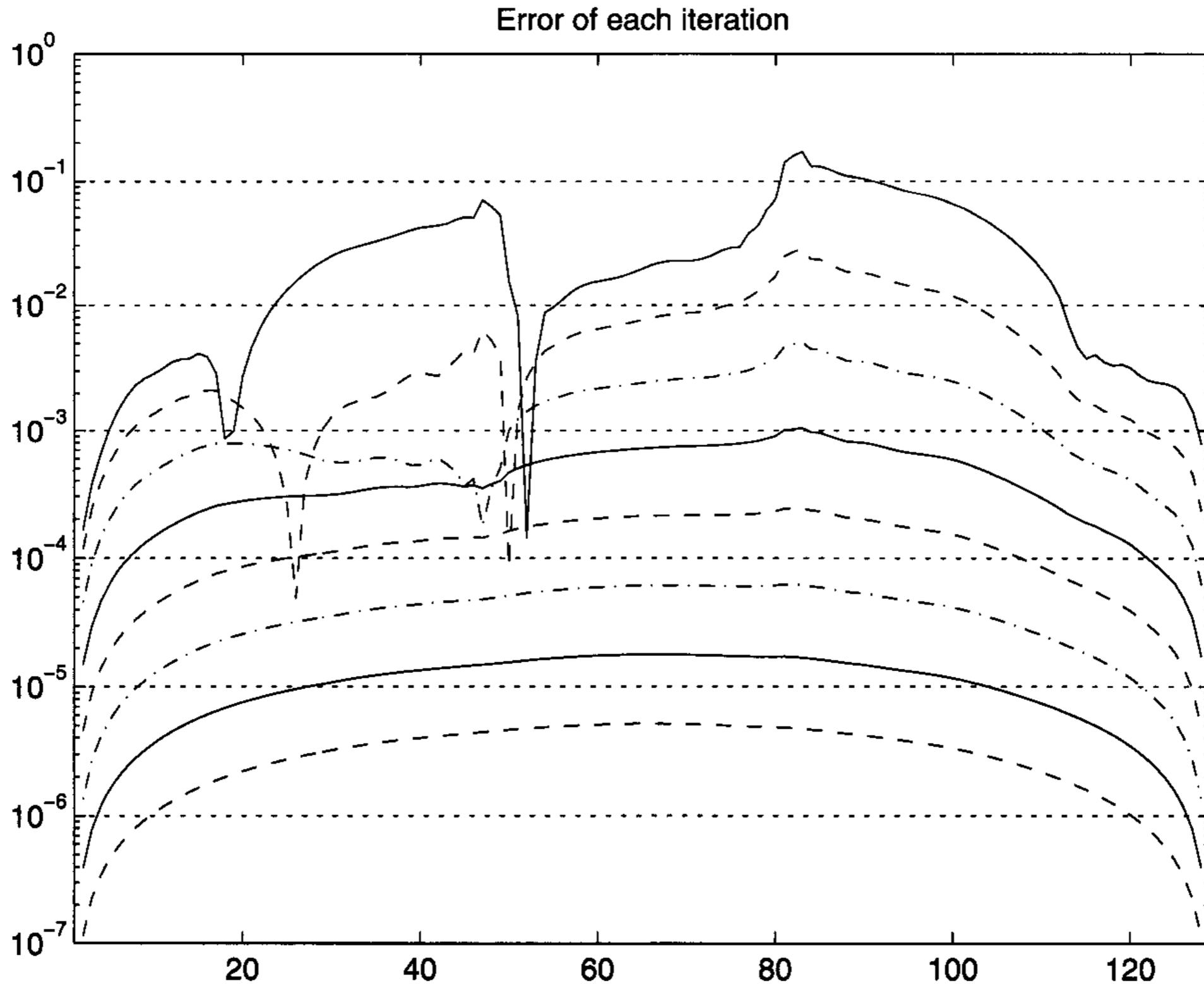
error  $e'_h$

after 3 coarse grid iter.

# Some notes on performance

- Can show:  $\|\text{error after step 5}\| \leq \rho \|\text{error before step 1}\|$   
with typically  $\rho \approx 0.1$  **independent** on  $h$  (i.e.  $N$ )!
- Typically for Jacobi alone:  $\rho \approx 0.99$

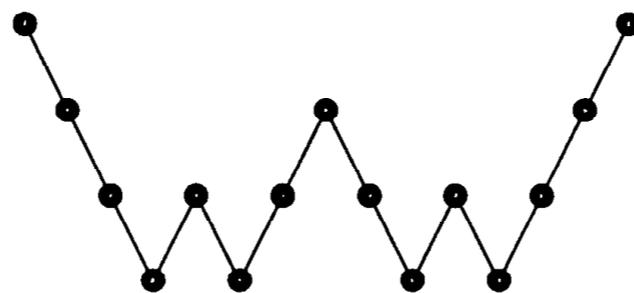
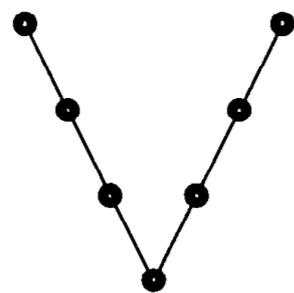
# Some notes on performance



# Some notes on performance

- Total cost per complete v-cycle:  $O(N)$  (N: number of grid points)
- Since error is reduced in each step by a constant independent on  $h$  (and thus  $N$ ), the total cost to get arbitrary low errors still is only  $O(N)$ :
- The finest level grid operations determine the cost!

$h$   
 $2h$   
 $4h$   
 $8h$



**FMG**

