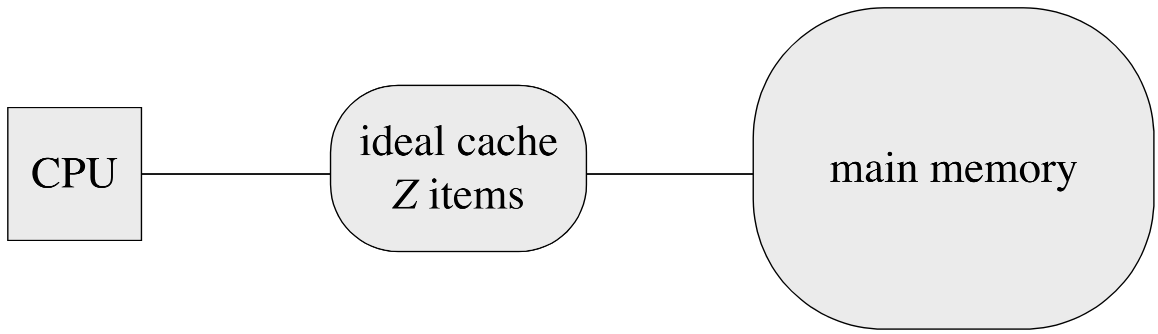

cache hit: CPU needs item in cache (fast)
cache miss: CPU needs item not in cache
— item loaded into cache for future use, replacing some other item

optimal replacement: on cache miss, loaded item replaces item that will not be needed for the *longest time in the future*

[more realistic scheme: **LRU replacement** — replace least recently used item
— provably within small constant factor of optimal, but much harder to analyze]

fully associative — any item in memory can go anywhere in the cache
[real caches have limited associativity, which causes "unlucky" memory-access patterns to go same place in cache
…effectively shrinks cache in these cases]

temporal locality — same item is re-used for several computations that are close to one another in time ⇒ still in-cache ⇒ efficient

[there is also **spatial locality** — items close to one another in main memory are used close in time … exploited by **cache lines**, TBD]

cache complexity — the number of cache misses $Q(n; Z)$ required for a given algorithm running on a problem of size n with cache of size Z
… usually only given as **asymptotic** result for large n, Z, ignoring constant factors

asymptotic notation:

we say a function $f(n)$ is $O(g(n))$ if $g(n)$ is an **asymptotic upper bound** for $f(n)$, ignoring constant factors. Technically, if $f(n) < C\, g(n)$ for some constant C for all sufficiently large n (i.e., for all $n > N$ for some N)

we say a function $f(n)$ is $\Omega(g(n))$ if $g(n)$ is an **asymptotic lower bound** for $f(n)$, ignoring constant factors. Technically, if $f(n) > C\, g(n)$ for some constant C for all sufficiently large n (i.e., for all $n > N$ for some N)

we say a function $f(n)$ is $\Theta(g(n))$ if $g(n)$ is an **asymptotic tight bound** for $f(n)$, ignoring constant factors. Technically, if $f(n)$ is *both* $O(g(n))$ and $\Omega(g(n))$