

Sample MidTerm Examination Questions

- Let  $\Sigma = \{a, b, c\}$  and let  $A = \{a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i = j \text{ or } i = k\}$ . Describe (in English) a pushdown automaton that recognizes  $A$ .
  - Let  $R$  be the regular expression  $\Sigma^* 1100 \Sigma^*$  where  $\Sigma = \{0, 1\}$ . Let  $D = L(R)$  and let  $E = \overline{D}$ , the complement of  $D$ . Give the state diagram of a DFA with at most 5 states that recognizes  $E$ .
- Let  $\Sigma = \{(, )\}$  and let  $P$  be the language consisting of all strings of properly nested parentheses. For example,  $P$  contains “ $()()$ ”, “ $((()))$ ”, “ $((()((()((()))))))$ ” and “ $\epsilon$ ”, but not “ $()()$ ” and “ $((()$ ”.
  - Give a CFG that generates  $P$ .
  - Show that  $P$  is not a regular language.
- Let  $A = \{a^i b^j c^i \mid i \leq j \leq 2i\}$ . Prove that  $A$  is not a context-free language.
  - Let  $B = \{a^i b^j \mid i \leq j \leq 2i\}$ . Give an *unambiguous* context-free grammar generating  $B$ .
- Let  $D = \{\langle M \rangle \mid M \text{ is a TM that accepts the input string } 101\}$ .
  - Show that  $D$  is undecidable.  
(Do not use Rice’s theorem. If you don’t know Rice’s theorem, ignore this comment.)
  - Show that the complement of  $D$  is not Turing-recognizable.
- A *2-way pushdown automaton* (2WAY-PDA) is a nondeterministic pushdown automaton that has a single stack and that can move its input head in both directions on the input tape. In addition we assume that a 2WAY-PDA is capable of detecting when its input head is at either end of its input tape. A 2WAY-PDA accepts its input by entering an accept state.
  - Show that a 2WAY-PDA can recognize the language  $\{a^m b^m c^m \mid m \geq 0\}$ .
  - Let  $E_{2\text{WAY-PDA}} = \{\langle P \rangle \mid P \text{ is a 2WAY-PDA which recognizes the empty language}\}$ . Show that  $E_{2\text{WAY-PDA}}$  is not decidable.
- Consider the infinite two-dimensional grid,  $G = \{(m, n) \mid m \text{ and } n \text{ are integers}\}$ . Every point in  $G$  has 4 neighbors, North, South, East, and West, obtained by varying  $m$  or  $n$  by  $\pm 1$ . Starting at the origin  $(0, 0)$ , a string of commands **N**, **S**, **E**, **W**, generates a path in  $G$ . For example, the string **NESW**, generates a path clockwise around a unit square touching the origin. Say that a path is *closed* if it starts at the origin and ends at the origin.

Let  $C$  be the collection of all strings over  $\Sigma = \{N, S, E, W\}$  that generate a closed path.

  - Give a clear mathematical description of  $C$  as a language.
  - Describe in English two CFLs,  $A$  and  $B$ , such that  $C = A \cap B$ .  
Give a CFG that generates  $A$ .
  - Prove that  $C$  is not context-free.
- Let  $\Sigma = \{0, 1\}$ . Consider the problem of testing whether a PDA accepts some string of the form  $\{w \mid w \in 0^* 1^*\}$ . Is this problem decidable? Prove your answer.