region:

tiles:
tiling:
• Is there a tiling?
• How many?
• About how many?
• Is a tiling easy to find?
• Is it easy to prove a tiling doesn’t exist?
• Is it easy to convince someone that a tiling doesn’t exist?
• What is a “typical” tiling?
• Relations among different tilings
• Special properties, such as symmetry
• Infinite tilings
Is there a tiling?

Tiles should be “mathematically interesting.”

12 pentominos:
Number of tilings of a $6 \times 10$ rectangle: 2339

Found by “brute force” computer search (uninteresting)
Is there a tiling with 31 dominos (or dimers)?
color the chessboard:

Each domino covers one black and one white square, so 31 dominos cover 31 white squares and 31 black squares. There are 32 white squares and 30 black squares in all, so a tiling does not exist.

Example of a coloring argument.
What if we remove one black square and one white square?
What if we remove \textbf{two} black squares and \textbf{two} white squares?
Another coloring argument: can a 10 × 10 board be tiled with 1 × 4 rectangles (in any orientation)?
Every tile covers each color an even number (including 0) of times. But the board has 25 tiles of each color, so a tiling is impossible.
Coloring doesn’t work:

\[ T(1) \quad T(2) \quad T(3) \quad T(4) \]

\( n \) hexagons on each side \((n(n+1)/2 \text{ hexagons in all})\)

Can \( T(n) \) be covered by “tribones”?
Yes for $T(9)$:

Conway: The triangular array $T(n)$ can be tiled by tribones if and only if $n = 12k, 12k + 2, 12k + 9, 12k + 11$ for some $k \geq 0$.

Smallest values: $0, 2, 9, 11, 12, 14, 21, 23, 24, 26, 33, 35, \ldots$.

Cannot be proved by a coloring argument (involves a nonabelian group)
How many tilings?

There are 2339 ways (up to symmetry) to tile a $6 \times 10$ rectangle with the 12 pentominoes. Found by computer search: not so interesting.

First significant result on the enumeration of tilings due to Kasteleyn, Fisher–Temperley (independently, 1961):

The number of tilings of a $2m \times 2n$ rectangle with $2mn$ dominos is

$$4^{mn} \prod_{j=1}^{m} \prod_{k=1}^{n} \left( \cos^2 \frac{j\pi}{2m+1} + \cos^2 \frac{k\pi}{2n+1} \right).$$

$\Pi$ means “product.” $\pi = 180^\circ$. E.g.,

$$\cos \frac{2\pi}{5} = \cos 72^\circ = 0.3090169938 \cdots$$
For instance, \( m = 2, n = 3 \):
\[
4^6 \left( \cos^2 36^\circ + \cos^2 25.71^\circ \right) \left( \cos^2 72^\circ + \cos^2 25.71^\circ \right) \\
\times \left( \cos^2 36^\circ + \cos^2 51.43^\circ \right) \left( \cos^2 72^\circ + \cos^2 51.43^\circ \right) \\
\times \left( \cos^2 36^\circ + \cos^2 77.14^\circ \right) \left( \cos^2 72^\circ + \cos^2 77.14^\circ \right)
\]

\[= 4^6 (1.4662)(.9072)(1.0432)(.4842)(.7040)(.1450) \]

\[= 281 \]
Aztec diamonds:

AZ(1)  AZ(2)  AZ(3)  AZ(7)
Eight domino tilings of AZ(2), the Aztec diamond of order 2:
Elkies-Kuperberg-Larsen-Propp (1992): The number of domino tilings of $AZ(n)$ is $2^{n(n+1)/2}$.

(four proofs originally, now around 12)

<table>
<thead>
<tr>
<th></th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
</tr>
</thead>
<tbody>
<tr>
<td>2</td>
<td>8</td>
<td>64</td>
<td>1024</td>
<td>32768</td>
<td>2097152</td>
<td>268435456</td>
<td></td>
</tr>
</tbody>
</table>

Since $2^{(n+2)(n+1)/2}/2^{(n+1)n/2} = 2^{n+1}$, we would like to associate $2^{n+1}$ Aztec diamonds of order $n + 1$ with each Aztec diamond of order $n$, so that each Aztec diamond of order $n + 1$ occurs exactly once. This is done by domino shuffling.
About how many tilings? $AZ(n)$ is a “skewed” $n \times n$ square. How do the number of domino tilings of $AZ(n)$ and an $n \times n$ square ($n$ even) compare?

If a region with $N$ squares has $T$ tilings, then it has (loosely speaking) $\sqrt[4]{NT}$ degrees of freedom per square.

Number of tilings of $AZ(n)$: $T = 2^{n(n+1)/2}$

Number of squares of $AZ(n)$:

$$N = 2n(n + 1)$$

Number of degrees of freedom per square:

$$\sqrt[4]{NT} = \sqrt[4]{2^{n(n+1)/2}} = 1.189207115 \cdots$$
Number of tilings of $2n \times 2n$ square:

$$4n^2 \prod_{j=1}^{n} \prod_{k=1}^{n} \left( \cos^2 \frac{j\pi}{2n+1} + \cos^2 \frac{k\pi}{2n+1} \right).$$

**Theorem** (Kasteleyn, et al.). The number of domino tilings of a $2n \times 2n$ square is about $C^{4n^2}$, where

$$C = e^{G/\pi}$$

$$= 1.338515152 \cdots.$$

Moreover,

$$G = 1 - \frac{1}{3^2} + \frac{1}{5^2} - \frac{1}{7^2} + \cdots$$

$$= 0.9159655941 \cdots$$

(Catalan’s constant).

Thus the square board is “easier” to tile than the Aztec diamond ($1.3385 \cdots$ degrees of freedom per square vs. $1.189207115 \cdots$).
If a tiling exists then it may be difficult to \textit{find}, but it is easy to \textit{demonstrate}.
What if a tiling doesn’t exist? Is it easy to demonstrate that this is the case?

In general, no (even for $1 \times 3$ rectangular tiles). But yes (!) for domino tilings.

16 white squares and 16 black squares
The six black squares with ● are adjacent to a total of five white squares marked *. No tiling can cover all six black square marked with ●.

**Philip Hall** (1935): If a region cannot be tiled with dominos, then one can *always* find such a demonstration of impossibility.
Tilings rectangles with rectangles: two results

Can a $7 \times 10$ rectangle be tiled with $2 \times 3$ rectangles (in any orientation)?

Clearly no: a $2 \times 3$ rectangle has 6 squares, while a $7 \times 10$ rectangle has 70 squares (not divisible by 6).
Can a $17 \times 28$ rectangle be tiled with $4 \times 7$ rectangles?
Can a $17 \times 28$ rectangle be tiled with $4 \times 7$ rectangles?

**No**: there is no way to cover the first column.
Can a $10 \times 15$ rectangle be tiled with $1 \times 6$ rectangles?
**de Bruijn-Klarner**: an $m \times n$ rectangle can be tiled with $a \times b$ rectangles if and only if:

- The first row and first column can be covered.
- $m$ or $n$ is divisible by $a$, and $m$ or $n$ is divisible by $b$.

Since neither 10 nor 15 are divisible by 6, the $10 \times 15$ rectangle **cannot** be tiled with $1 \times 6$ rectangles.
Let \( x > 0 \), such as \( x = \sqrt{2} \). Can a square be tiled with finitely many rectangle similar to a \( 1 \times x \) rectangle (in any orientation)? In other words, can a square be tiled with finitely many rectangles all of the form \( a \times ax \) (where \( a \) may vary)?
\[ x = \frac{2}{3} \]

\[ 3x - 2 = 0 \]
\begin{align*}
x &= \frac{5 + \sqrt{5}}{10} = 0.7236067977 \cdots \\
5x^2 - 5x + 1 &= 0 \\
\text{Other root: } \frac{5 - \sqrt{5}}{10} &= 0.2763932023 \cdots
\end{align*}
\[ x = 0.5698402910 \cdots \]
\[ x^3 - x^2 + 2x - 1 = 0 \]

Other roots:

\[ 0.215 + 1.307\sqrt{-1} \]
\[ 0.215 - 1.307\sqrt{-1} \]
Freiling-Rinne (1994), LaczkovichSzekeres (1995): A square can be tiled with finitely many rectangles similar to a $1 \times x$ rectangle if and only if:

- $x$ is the root of a polynomial with integer coefficients.
- If $a + b\sqrt{-1}$ is another root of the polynomial of least degree satisfied by $x$, then $a > 0$. 
Examples. \( x = \sqrt{2} \). Then \( x^2 - 2 = 0 \). Other root is \( -\sqrt{2} < 0 \). Thus a square cannot be tiled with finitely many rectangles similar to a \( 1 \times \sqrt{2} \) rectangle.

\[ x = \sqrt{2} + \frac{17}{12}. \] Then 
\[ 144x^2 - 408x + 1 = 0. \]

Other root is 
\[ -\sqrt{2} + \frac{17}{12} = 0.002453 \cdots > 0, \]
so a square can be tiled with finitely many rectangles similar to a \( 1 \times (\sqrt{2} + \frac{17}{12}) \) rectangle.
\[ x = \sqrt[3]{2}. \] Then \( x^3 - 2 = 0 \). Other roots:
\[-\frac{\sqrt[3]{2}}{2} \pm \frac{\sqrt[3]{2}\sqrt{3}}{2}\sqrt{-1}.\]

Since \(-\frac{\sqrt[3]{2}}{2} < 0\), a square cannot be tiled with finitely many rectangles similar to a \( 1 \times \sqrt[3]{2} \) rectangle.

Let \( r/s \) be a rational number and \( x = \frac{r}{s} + \sqrt[3]{2} \). Other roots:
\[
\left( \frac{r}{s} - \frac{\sqrt[3]{2}}{2} \right) \pm \frac{\sqrt[3]{2}\sqrt{3}}{2}\sqrt{-1}.
\]

Therefore a square can be tiled with finitely many rectangles similar to a \( 1 \times (\frac{r}{s} + \sqrt[3]{2}) \) rectangle if and only if
\[
\frac{r}{s} > \frac{\sqrt[3]{2}}{2}.
\]
What is a “typical” tiling?

A random domino tiling of a $12 \times 12$ square:

No obvious structure.
A random tiling of the Aztec diamond of order 50:

“Regular” at the corners, chaotic in the middle.

What is the region of regularity?
**Arctic Circle Theorem** (Jockusch-Propp-Shor, 1995). For very large $n$, and for “most” domino tilings of the Aztec diamond $AZ(n)$, the region of regularity “approaches” the outside of a circle tangent to the four limiting sides of $AZ_n$.

The tangent circle is the **Arctic circle**. Outside this circle the tiling is “frozen.”
Relations among tilings

Two domino tilings of a region in the plane:

A **flip** consists of reversing the orientation of two dominos forming a $2 \times 2$ square.

**Domino flipping theorem** (Thurston, et al.). *If $R$ has no holes* (simply-connected), *then any domino tiling of $R$ can be reached from any other by a sequence of flips.*
Flipping theorem is **false** if holes are allowed.

![Diagram](image-url)
Confronting infinity:
(1) A finite (bounded) region, infinitely many tiles.

Total area: \( \frac{1}{1/2} + \frac{1}{2/3} + \frac{1}{3/4} + \cdots = 1 \)

Can a square of side 1 be tiled with squares of sides 1/2, 1/6, 1/12, \ldots (once each)?
Unsolved (Meir & Moser, 1968), but the tiles will fit into a square of side \(1 + \frac{1}{350}\) (not a tiling, since there is leftover space).
Confronting infinity: (2) Finitely many tiles, but an indeterminately large region.

Which polyominos can tile rectangles?

order 2

order 4
The order of a polyomino is the least number of copies of it needed to tile some rectangle.

No polyomino has order 3.
Known orders: 4, 8, 12, 16, \ldots, 4n, \ldots

1, 2, 10, 18, 50, 138, 246, 270

order 246

order 270

order 468

Unknown: order 6? odd order?
no order

Cannot tile a rectangle (order does not exist).

Deep result from mathematical logic: there does not exist an algorithm (computer program) to decide whether a finite set of polyominos tiles some region containing a given square.
**Consequence.** Let $LS(n)$ be the largest size square that can be covered by disjoint copies of some set $S$ of polynominos with a total of $n$ squares, such that the polynominos in $S$ cannot tile the whole plane (so $LS(n)$ is finite).
If \( f(n) \) is any function that can be computed on a computer (with infinite memory), such as

\[
\begin{align*}
  f(n) &= n^n, \\
  f(n) &= n^{n^n} \\
  f(n) &= n^{n \cdots n} \quad \text{(n n’s), \ldots,}
\end{align*}
\]

then \( \text{LS}(n) > f(n) \) for large \( n \). (Otherwise a computer could simply check all possible tilings such that each polyomino in the tiling overlaps a given square of side length \( k \), for all \( k \leq f(n) \).)
Confronting infinity: (3) Tiling the plane