

# Lovász local lemma

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Paper Presentation

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- Lovász local lemma (Symmetric Case)
- Lovász local lemma (Asymmetric Case)
- Example of application of Lovász local lemma.
- Algorithmic Lovász local lemma.
- Applications and proving a theorem using the Lovász local lemma.

# Lemma: Symmetric Case

- Let  $A_1, A_2, \dots, A_n$  be events in an arbitrary probability space. Suppose each event  $A_i$  is independent of all the other events, except at most  $d$  of them, and that  $\Pr[A_i] \leq p$  for all  $A_i$ , and if certain conditions on  $p$  and  $d$  are met, then:

$$\Pr \left[ \bigwedge_{i=1}^n \bar{A}_i \right] > 0$$

## Conditions on $p$ and $d$

- **Lemma I** (Lovász and Erdős 1973; published 1975)

$$4pd \leq 1$$

- **Lemma II** (Lovász 1977; published by [Joel Spencer](#))

$$ep(d+1) \leq 1$$

## **Lemma III** (Shearer 1985)

$$\begin{cases} p < \frac{(d-1)^{d-1}}{d^d} & d > 1 \\ p = \frac{1}{2} & d = 1 \end{cases}$$

## Lemma: Asymmetric (general) Case (events can have different probability bounds)

- Let  $A_1, A_2, \dots, A_n$  be events in an arbitrary probability space. A **directed digraph**  $D=(V,E)$  on the set of vertices  $V = \{1, 2, \dots, n\}$  is called a **dependency digraph** for the events  $A_1, A_2, \dots, A_n$  if for each  $1 \leq i \leq n$ , the event  $A_i$  is mutually independent of all the events  $\{A_j : (i,j) \notin E\}$ . Suppose  $D=(V,E)$  is a dependency digraph for the above events, and suppose there are real numbers  $x_1, x_2, \dots, x_n$  such that  $0 \leq x_i < 1$  and

$$\Pr \left[ \bigwedge_{i=1}^n \bar{A}_i \right] \leq x_i \prod_{(i,j) \in E} (1 - x_j)$$

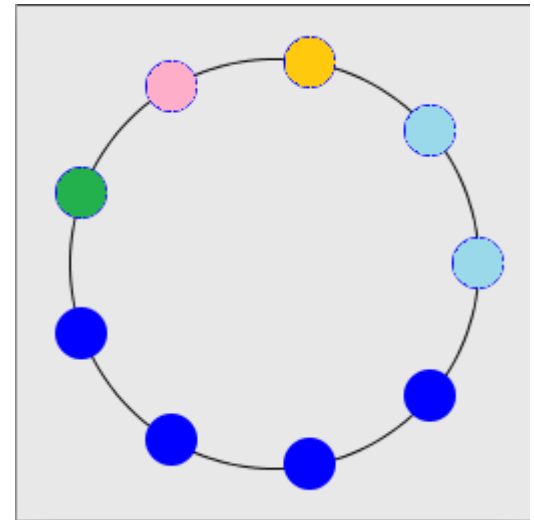
for all  $1 \leq i \leq n$ , then

$$\Pr \left[ \bigwedge_{i=1}^n \bar{A}_i \right] \geq \prod_{i=1}^n (1 - x_i)$$

## Example

- Suppose  $11n$  points are placed on a circle and colored with  $n$  different colors.
- Each color is applied to exactly 11 points.
- In any such coloring, there must be subset of  $n$  points each colored a different color and not containing any adjacent points.
- Let's pick a point of each color randomly ( $p=1/11$ ).
- Event we want to avoid corresponds to picking a pair  $(a,b)$  of adjacent points ( $p=1/121$ ).
- 21 pairs include same color as  $a$ , 21 pairs include same color as  $b$ , if disjoint,  $(a,b)$  is dependent on 42 pairs.  $d=42$ .

$$ep(d + 1) \approx 0.966 < 1.$$



# Algorithmic Lovász local lemma

## Algorithm

Given a set of bad events  $A$  that we wish to avoid, and that are determined by a collection of mutually independent random variables  $P$ :

1. For every random variable  $P$ , make an assignment.
2. While there exists a satisfied event  $A$ :
  1. Pick a satisfied event  $A$ .
  2. For all  $P$  that determine  $A$ , give a new assignment to  $P$ .
3. Return the final assignment

# Applications

- Probabilistic method proofs.
- Random Graphs.



# Theorem:

Let  $H=(V,E)$  be a hypergraph in which every edge has at least  $k$  elements, and suppose that each edge of  $H$  intersects at most  $d$  other edges. If  $e(d+1) \leq 2^{k-1}$  then  $H$  is two-colorable.

## Proof

- Color each vertex  $v$  of  $H$  randomly and independently, either blue or red with equal probability. For each edge  $f$  of  $H$ , let  $A_f$  be the event that  $f$  is monochromatic.

$$\Pr [A_f] \geq 2/2^{|f|} \leq 1/2^{k-1}$$

- Therefore if  $ep(d+1) \leq 1$  (Lovász local lemma) then  $e(d+1) \leq 2^{k-1}$

and  $H$  is two-colorable.

# References

- Alon Noga, Spencer Joel H. “5. The Local Lemma”, *The Probabilistic Method*. Hoboken, NJ: John Wiley & Sons, Inc. 2008. 3<sup>rd</sup> ed.
- <http://en.wikipedia.org>:Lovász local lemma
- <http://en.wikipedia.org>:Algorithmic Lovász local lemma