

## Lecture notes on matroid union

From any matroid  $M = (E, \mathcal{I}(M))$ , one can construct a dual matroid  $M^* = (E, \mathcal{I}(M^*))$ .

**Theorem 1** *Let  $\mathcal{I}(M^*) = \{X \subseteq E(M) : E(M) \setminus X \text{ contains a base of } M\}$ . Then  $M^* = (E, \mathcal{I}(M^*))$  is a matroid with rank function*

$$r_{M^*}(X) = |X| + r_M(E \setminus X) - r_M(E).$$

There are several ways to show this. One is to first show that indeed the size of the largest subset of  $X$  in  $\mathcal{I}(M^*)$  has cardinality  $|X| + r_M(E \setminus X) - r_M(E)$  and then show that  $r_{M^*}$  satisfies the three conditions that a rank function of a matroid needs to satisfy (the third one, submodularity, follows from the submodularity of the rank function for  $M$ ).

One can use Theorem 1 and matroid intersection to get a good characterization of when a graph  $G = (V, E)$  has two edge-disjoint spanning trees. Indeed, letting  $M$  be the graphic matroid of the graph  $G$ , we get that  $G$  has two edge-disjoint spanning trees if and only if

$$\max_{S \in \mathcal{I}(M) \cap \mathcal{I}(M^*)} |S| = |V| - 1.$$

For the graphic matroid, we know that  $r_M(F) = n - \kappa(F)$  where  $n = |V|$  and  $\kappa(F)$  denotes the number of connected components of  $(V, F)$ . But by the matroid intersection theorem, we can write:

$$\begin{aligned} \max_{S \in \mathcal{I}(M) \cap \mathcal{I}(M^*)} |S| &= \min_{E_1 \subseteq E} r_M(E_1) + r_{M^*}(E \setminus E_1) \\ &= (n - \kappa(E_1)) + (|E \setminus E_1| + \kappa(E) - \kappa(E_1)) \\ &= n + 1 + |E \setminus E_1| - 2\kappa(E_1), \end{aligned}$$

where we replaced  $\kappa(E)$  by 1 since otherwise  $G$  would even have one spanning tree. Rearranging terms, we get that  $G$  has two edge-disjoint spanning trees if and only if for all  $E_1 \subseteq E$ , we have that  $|E \setminus E_1| \geq 2(\kappa(E_1) - 1)$ . If this inequality is violated for some  $E_1$ , we can add to  $E_1$  any edge that does not decrease  $\kappa(E_1)$ . In other words, if the connected components of  $E_1$  are  $V_1, V_2, \dots, V_p$  then we can assume that  $E_1 = E \setminus \delta(V_1, V_2, \dots, V_p)$  where  $\delta(V_1, \dots, V_p) = \{(u, v) \in E : u \in V_i, v \in V_j \text{ and } i \neq j\}$ . Thus we have shown:

**Theorem 2**  *$G$  has two edge-disjoint spanning trees if and only if for all partitions  $V_1, V_2, \dots, V_p$  of  $V$ , we have*

$$|\delta(V_1, \dots, V_p)| \geq 2(p - 1).$$

Going back to the spanning tree game, it is now clear that if the graph does not have two edge-disjoint spanning trees then player 1 has a winning strategy. He/she just needs to delete edges from  $\delta(V_1, \dots, V_p)$  for the partition given by theorem 2.

Theorem 2 can be generalized.

**Theorem 3**  $G$  has  $k$  edge-disjoint spanning trees if and only if for all partitions  $V_1, V_2, \dots, V_p$  of  $V$ , we have

$$|\delta(V_1, \dots, V_p)| \geq k(p-1).$$

From two matroids  $M_1 = (E, \mathcal{I}(M_1))$  and  $M_2 = (E, \mathcal{I}(M_2))$ , we can also define its union by  $M_1 \cup M_2 = (E, \mathcal{I})$  where  $\mathcal{I} = \{S_1 \cup S_2 : S_1 \in \mathcal{I}(M_1), S_2 \in \mathcal{I}(M_2)\}$ .

We can show that:

**Theorem 4**  $M_1 \cup M_2$  is a matroid with rank function

$$r_{M_1 \cup M_2}(X) = \min_{F \subseteq E} \{|E \setminus F| + r_{M_1}(F) + r_{M_2}(F)\}.$$

**Proof:** To show that it is a matroid, assume that  $X, Y \in \mathcal{I}$  with  $|X| < |Y|$ . Let  $X = X_1 \cup X_2$  and  $Y = Y_1 \cup Y_2$  where  $X_1, Y_1 \in \mathcal{I}(M_1)$  and  $X_2, Y_2 \in \mathcal{I}(M_2)$ . We can furthermore assume that the  $X_i$ 's are disjoint and so are the  $Y_i$ 's. Finally we assume that among all choices for  $X_1, X_2, Y_1$  and  $Y_2$ , we choose the one maximizing  $|X_1 \cap Y_1| + |X_2 \cap Y_2|$ . Since  $|Y| > |X|$ , we can assume that  $|Y_1| > |X_1|$ . Thus, there exists  $e \in (Y_1 \setminus X_1)$  such that  $X_1 \cup \{e\}$  is independent for  $M_1$ . The maximality implies that  $e \notin X_2$  (otherwise consider  $X_1 \cup \{e\}$  and  $X_2 \setminus \{e\}$ ). But this implies that  $X \cup \{e\} \in \mathcal{I}$  as desired.

We now show the expression for the rank function. The fact that it is  $\leq$  is obvious as an independent set  $X \in \mathcal{I}$  has size  $|X \setminus F| + |X \cap F| \leq |E \setminus F| + r_{M_1}(F) + r_{M_2}(F)$  and this is true for any  $F$ .

For the converse, notice that  $X \in \mathcal{I}$  is such that  $X = X_1 \cup X_2$  with  $X_1 \in \mathcal{I}(M_1)$  and  $X_2 \in \mathcal{I}(M_2)$ . We can furthermore assume that  $X_1$  and  $X_2$  are disjoint and that  $r_{M_2}(X_2) = r_{M_2}(E)$  (otherwise add elements to  $X_2$  and possibly remove them from  $X_1$ ). Thus we can assume that  $|X| = |X_1| + r_{M_2}(E)$  and that  $X_1 \in \mathcal{I}(M_1) \cap \mathcal{I}(M_2^*)$ . The proof is completed by using the matroid intersection theorem and Theorem 1:

$$\begin{aligned} r_{M_1 \cup M_2}(E) &= \max_{X_1 \in \mathcal{I}(M_1) \cap \mathcal{I}(M_2^*)} (|X_1| + r_{M_2}(E)) \\ &= \min_{E_1 \subseteq E} (r_{M_1}(E_1) + r_{M_2^*}(E \setminus E_1) + r_{M_2}(E)) \\ &= \min_{E_1 \subseteq E} (r_{M_1}(E_1) + |E \setminus E_1| + r_{M_2}(E_1) - r_{M_2}(E) + r_{M_2}(E)) \\ &= \min_{E_1 \subseteq E} (|E \setminus E_1| + r_{M_1}(E_1) + r_{M_2}(E_1)), \end{aligned}$$

as desired. △