The Pixel Array method for nonlinear systems, and applications to numerical PDEs

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in the Numerical Methods for PDE seminar
Introduction

What to expect
A few examples
Linearizing is the key to success

Details on the Pixel Array method

Relation to Numerical methods for PDE

Open questions

Conclusion
The subject of this talk is a new method for solving nonlinear systems.

- It’s very different: you get *all solutions* inside a bounding box.
What to expect from the talk

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- It’s very different: you get all solutions inside a bounding box.
- It arose from my work on applied category theory.
- There will be no category theory (CT) in this talk.
- The idea originated there, but all you need to know is matrix mult.
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- It’s very different: you get all solutions inside a bounding box.
  - It arose from my work on applied category theory.
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  - The idea originated there, but all you need to know is matrix mult.

- The talk is more about solving systems of equations than about PDE.
  - I’ll discuss how the method applies to PDE near the end.

- I know next to nothing about PDEs and numerical methods.
  - I’m a foreigner here; please forgive my ’dialect’ and ’cultural ignorance’.
  - I’d appreciate any feedback you may have.
  - I hope our communities—applied math and CT—can become friends.
An image, for your imagination

Separately plot the solutions to equations: $f(x, w) = 0$ and $g(w, y) = 0$.

- Plot each in its own bounding box, say in the range $[-1.5, 1.5]$.
- Consider the plots as matrices $M, N$ whose entries are on/off pixels.
- That is, $M$ and $N$ are boolean matrices corresponding to $f$ and $g$. 
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Multiplying these two matrices \( MN \) yields...

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David I. Spivak (MIT)

The Pixel Array method

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- Consider the plots as matrices \( M, N \) whose entries are on/off pixels.
- That is, \( M \) and \( N \) are boolean matrices corresponding to \( f \) and \( g \).
- Multiplying these two matrices \( MN \) yields the simultaneous solution.
- For example, plot equations \( x^2 = w \) and \( w = 1 - y^2 \), and multiply.
A more complex example

The following eq’s are not differentiable, nor even defined everywhere.

\[
\cos \left( \ln(z^2 + 10^{-3}x) \right) - x + 10^{-5}z^{-1} = 0 \quad \text{(Equation 1)}
\]
\[
\cosh(w + 10^{-3}y) + y + 10^{-4}w = 2 \quad \text{(Equation 2)}
\]
\[
\tan(x + y)(x - 2)^{-1}(x + 3)^{-1}y^{-2} = 1 \quad \text{(Equation 3)}
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Q: For what values of \(w\) and \(z\) does a simultaneous solution exist? \(^1\)

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Another way to linearize

A popular refrain in mathematics:

- The only problems we know how to solve are linear ones.
- We approach any other problem by trying to linearize it.
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A popular refrain in mathematics:

- The only problems we know how to solve are linear ones.
- We approach any other problem by trying to linearize it.

The Pixel Array slogan: *Plotting is a way of linearizing.*

- Once we plot, we have matrices (more generally, arrays / tensors).
- We can use *linear algebra* to put plots together and solve big systems.
Selling points

The Pixel Array method has the following features:

- it returns all solutions in a given bounding box;
- it’s much faster than quasi-Newton methods for finding “all solutions”;
- it introduces no false negatives;
- it works for non-differentiable or even discontinuous functions;
- it’s not iterative and requires no initial guess, in contrast with quasi-Newton methods;
- it’s elementary—relying only on generalized matrix arithmetic—hence has low barrier to entry; and
- it provides insights, by showing the whole solution set at once.
So where are the limitations hiding?

Takes original plots as input.

- The Pixel Array (PA) method requires a plot for each equation.
- The plot of $f(x_1, \ldots, x_n) = 0$ is $n$-dimensional; can be costly.
  - But there are methods: sampling, zero-crossing, “interval arithmetic”.
  - Nice aside: you can even use raw data as your plots.
- PA method is fairly robust, as we saw in butterfly picture.
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Introduces false positives.

- False positives can occur arbitrarily far from real solutions.
- Accuracy often degrades when a local derivative approaches 0 or \( \infty \).
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Speed comes from unexposed variables.

- I will discuss exposed/unexposed variables in depth.
- For now, recall butterfly picture: we exposed $w$ and $z$, but not $x$ or $y$.
- Comparison with quasi-Newton is not straightforward.
Outline

1. Introduction

2. Details on the Pixel Array method
   - Overview
   - Why it works
   - Wiring diagrams
   - Clustering
   - False positives, true negatives
   - Visualizing higher-order arrays

3. Relation to Numerical methods for PDE

4. Open questions
Equations and wiring diagrams

Consider an arbitrary system of equations having the following form:

\[ f_1(t, u, v) = 0 \]
\[ f_2(v, w, x) = 0 \]
\[ f_3(u, w, x, y) = 0 \]
\[ f_4(x, z) = 0 \]

Bold variables are those we want to expose; others are latent or unexposed.
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Bold variables are those we want to expose; others are latent or unexposed.

Said another way, we want \( \{(t, v, z) \mid \exists u, w, x, y : f_1 = f_2 = f_3 = f_4 = 0\} \).
Using array multiplication to solve systems

To solve the system for $t$, $v$, $z$, we plot each equation as an array.

\[
\begin{align*}
  f_1(t, u, v) &= 0 \\
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\end{align*}

PA says: multiply the arrays according to “variable sharing” diagram:
There are now two important directions to go from here:
- Explain general array multiplication as dictated by a wiring diagram.
- Demystify relationship between multiplying arrays and solving systems.
Let’s start by demystifying why the Pixel Array method works.
- As we’ll see, it just comes down to logic: AND, OR, TRUE, FALSE.
To keep things simple, let’s restrict our attention to matrix multiplication.
Multiplying Boolean matrices

The matrix multiplication formula works well in any semiring.

$$(MN)_{i,k} = \sum_j M_{i,j} \ast N_{j,k}.$$}

- Roughly, a semiring is a set with $0, 1, +, \ast$ that act reasonably.
- It’s like a ring, but you don’t need negatives (e.g. $\mathbb{N}$).
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- It’s like a ring, but you don’t need negatives (e.g. \(\mathbb{N}\)).

Today we’ll focus on the Boolean semiring, \(\mathbb{B}\).

- It has two elements \(\mathbb{B} = \{0, 1\}\).
- 0 means FALSE, and 1 means TRUE.
- Multiplication is given by boolean AND (denoted \(\wedge\)).
- Addition is given by boolean OR (denoted \(\vee\)).
- The only slightly unexpected thing is that \(1 + 1 = 1\).
The logic of matrix multiplication

In general, the PA method is to multiply higher-order arrays (tensors).

- But for now, let’s suppose $A, B$ are boolean matrices.
- Say that $A$ and $B$ are plots of $f(x, y) = 0$ and $g(y, z) = 0$, resp.
- So $A_{ij} = 1$ means $f(x, y) = 0$ in corresponding pixel (else $A_{ij} = 0$).
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- The $(i, k)$-entry of their product $AB$ is given by the formula:

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(AB)_{i,k} = \sum_j A_{i,j} \ast B_{j,k}
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\[
= \bigvee_j A_{i,j} \land B_{j,k}
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= \exists_j (A_{i,j} \land B_{j,k})
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$$ = \exists_j (A_{i,j} \land B_{j,k}) $$

- “There exists some $y$ such that $f(x, y) = 0$ and $g(y, z) = 0.$”
Visual example again

M: \text{pixelMatrix}( x^2 = w )

N: \text{pixelMatrix}( w = 1-y^2 )

ordinary matrix product, M*N
Demo

Here is a little demo.²

Terminal:

> cd /Users/davidspivak/Dropbox\  \(\text{(MIT)}/\text{Code/}\)
> notebook

²Thanks to David Sanders and Andreas Noack for lots of help with Julia.
Multiplying larger-order arrays

When two arrays share a common dimension, they can be multiplied.

- For example, say that $A$ is $r_1 \times \cdots \times r_m$ and that $B$ is $s_1 \times \cdots \times s_n$.
- If $r_1 = s_1$, the product is an $r_2 \times \cdots \times r_m \times s_2 \times \cdots \times s_n$ array.
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We will be drawing these situations using wiring diagrams.
A single formula exists to multiply arrays according to any wiring diagram.

- Basically: iterate over all the links, multiply array entries, and sum.
- But this is very naive: $O(n^{\#\text{links}})$.
  - Modern techniques are much faster, especially given parallel processors.
  - Plots of equations are sparse matrices.
  - Boolean matrices are very special, bit arithmetic.
  - So we can do much better than our current “naive” implementation.
Skip: a naive general array multiplication algorithm

Precondition: Wiring diagram $\Phi: P_1, \ldots, P_n \rightarrow P'$ and arrays $A_i \in \text{Arr}(P_i)$.

1 function $\text{Arr}(\Phi)(A_1, \ldots, A_n)$
2 $A' := 0$
3 for $i \in \text{multi_index}(\Phi)$ do
4     $a_i := 1$
5     for $j \in \{0, \ldots, n\}$ do
6         $i_j := \text{multi_index}_\Phi^j(i)$
7         $a_i := a_i \ast A_j(i_j)$
8     $i' := \text{multi_index}_\Phi'(i)$
9     $A'(i') := A'(i') + a_i$
10 return $A'$

$\triangleright A' \in \text{Arr}(P')$
$\triangleright$ One index per link
$\triangleright$ Indices for $j^{\text{th}}$ pack
$\triangleright$ Indices for outer pack
Example wiring diagrams for named operations

The same general array multiplication formula returns famous matrix products for the following wiring diagrams:

- **Multiplication:** $MN$
- **Khatri-Rao:** $M \odot N$
- **Trace:** $\text{Tr}(M)$
- **Hadamard:** $M \odot N$
- **Kronecker:** $M \otimes N$
- **Marginalize:** $\sum_i M_{i,j}$
Clustering as associative law

The associative law for matrix multiplication can be seen as clustering.
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**Diagram:**

- **Top Diagram:**
  - $A$, $B$, $C$
  - $AB \rightarrow BC$

- **Bottom Diagram:**
  - $f_1$, $f_2$, $f_3$, $f_4$
  - $f_1 \Rightarrow f_2 \Rightarrow f_3 \Rightarrow f_4$

- **Association Diagram:**
  - $A$, $B$, $C$
  - $A \Rightarrow BC$
  - $f_1$, $f_2$, $f_3$, $f_4$
  - $f_1 \Rightarrow f_23 \Rightarrow f_4$
**Clustering for speed**

Order of array multiplication doesn’t affect solution, but does affect speed.

Naive cost for multiplying arrays: $O(n^{\#\text{Links}})$

- Even if we can use sparsity, etc. solving big systems is expensive.
- Instead, we should use the associative law: cluster.

Naive cost goes from $O(n^7)$ to $O(n^4)$. 
Cluster trees

Different cluster strategies lead to different speeds.

\[
\text{cost} = (2 \times 16 \times 2) + (2 \times 2 \times 2) = 72
\]

\[
\text{cost} = (16 \times 2 \times 2) + (2 \times 16 \times 2) = 128
\]

Each strategy can be drawn as a “cluster tree”:

\[
A \rightarrow 64 \\
B \rightarrow 8 \quad \text{(Total: 72)}
\]

\[
A \rightarrow 64 \\
B \rightarrow 64 \quad \text{(Total: 128)}
\]

Tree tells you how much the computation will cost, using the strategy.

- Given serial processors, the cost is the sum of node values.
- Given parallel processors, the cost is the length of longest path.
Some clustering algorithms

We need to hierarchically cluster our wiring diagrams to minimize cost.
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Most hypergraph clustering algorithms don’t apply “out of the box”:

- Our wiring diagrams are “pointed” hypergraphs (outside circle).
- Our clustering should be hierarchical: all the way down.

Some existing hypergraph clustering algorithms:

- Greedy algorithm: look for pairwise cluster of minimize cost.
- Try 1000 random samples and pick the best. Or exhaustively search.
- Try to use min cut algorithms or spectral graph theory.
Where false positives come from

Discretization and matrix multiplication together introduce false positives.

- Discretization corresponds to an equivalence relation, denoted $\sim$.
  - All numbers in the same bin are equivalent.
- This causes discrepancy between what we want and what we get:
  - Want: $\{(x, z) \mid \exists y : f(x, y) = 0 \text{ and } g(y, z) = 0\}$.
  - Get: $\{(x, z) \mid \exists y_1 \sim y_2 : f(x, y_1) = 0 \text{ and } f(y_2, z) = 0\}$
- This provably approaches the correct solution as mesh is refined.
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True negatives: “Get” $\geq$ “Want”. Further refine the “on” pixels.
Before moving on to PDEs, here’s something I found interesting.

- Start with a sphere $x^2 + y^2 + z^2 = 1$, and plot it as a pixel array.
- The plot is 3D, say $100 \times 100 \times 100$.
- Now reshape it to a $100 \times 10000$ array.
- Plot it in a square box. What do you get?
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1 Introduction

2 Details on the Pixel Array method

3 Relation to Numerical methods for PDE
   ■ Interconnected dynamical systems
   ■ Pixel Array analysis

4 Open questions

5 Conclusion
In this section I’ll discuss how the Pixel Array method may apply to PDE.

- The whole PA idea came from category theory.
- What do pixel arrays, PDEs, and category theory have in common?
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- The whole PA idea came from category theory.
- What do pixel arrays, PDEs, and category theory have in common?
- Compositionality: analyze a system by assembling simple components.

Let’s go a bit deeper and look at the heat equation.
**The heat equation**

Spacial discretization of the 1-d heat equation $u_t = u_{xx}$ over $\Omega = [0, 5]$.

![Wiring Diagram](image)

Discretize $x$ with distance $h = 1$, we use Taylor’s thm to obtain four ODEs.

$$\dot{u}_i = u_{i-1} - 2u_i + u_{i+1} \quad (i = 1, \ldots, 4)$$
The heat equation

Spatial discretization of the 1-d heat equation $u_t = u_{xx}$ over $\Omega = [0, 5]$.

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Wiring diagram (1) is more general than the heat equation. It represents:

- Four interconnected dynamical systems, each defined by an ODE.
- Each ODE is a function of the neighbors’ time-varying states.
- Some ODEs, $u_1$ and $u_4$, involve parameters from the “outside world”.
Open dynamical systems

Let’s look at the boxes in isolation.

A box of the above shape represents an ODE $\dot{u} = f(u, a_1, \ldots, a_n)$.

- The left-hand ports represent input parameters $a_1(t), \ldots, a_n(t)$.
- The right-hand port represents the state $u(t)$, being exported.
- Called an open dynamical system (ODS), because it interacts with the outside world.
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- A closed dynamical system would look like this: \[ \ddot{u} \]

You could also have multiple outputs, \( u_1, \ldots, u_m \) representing functions, possibly projections, of the state: \( u_i = \pi_i(u) \).
Setup for the PA method: interconnected ‘open dynamical systems’.

- In each box \((i,j)\), there is an ODS \(\dot{u}_{ij} = f(u_{ij}, u_{\text{neighbors}})\).
  - These open dynamical systems are sharing variables with neighbors.
  - Together these form a more complicated ODS on the outer box.
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- Examples of interconnected ODS’s:
  - Finite differences, in any dimension, discretized in space.
**Interconnected open dynamical systems (ODS’s)**

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- **Examples of interconnected ODS’s:**
  - Finite differences, in any dimension, discretized in space.
  - Differential equation on a network (compartment model).
  - Systems of systems (smart grid, National Airspace System).
The Pixel Array method lets you assemble local results.

- **First choose some sort of data to plot for each ODS.**
  - Most obvious: plot the steady states $u = 0$.
  - But you could also try plotting solutions over time $[0, \tau]$.

- The plot is an array, indexed by the input and output wires.
  - For steady state data: index by constant input/output functions.
  - For $[0, \tau]$ solutions: index by classifying $[0, \tau]$ input/output functions.
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- The plot is an array, indexed by the input and output wires.
  - For steady state data: index by constant input/output functions.
  - For $[0, \tau]$ solutions: index by classifying $[0, \tau]$ input/output functions.
- Then put the plots together using the Pixel Array method.
- The result is a plot of the chosen sort for the whole system (PDE).
In the box we have an open dynamical system $\dot{u} = f(u, a_1, \ldots, a_n)$.

- A plot for this ODS will be an order-$(n + 1)$ array.
- For example, you can plot steady-state solutions.
  - Set $\dot{u} = 0$ and get an equation $f(u, a_i) = 0$.
  - This is just the bifurcation diagram of the dynamical system.
  - Plot it as a pixel array: PA method combines bifurcation diagrams.
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  - This is just the bifurcation diagram of the dynamical system.
  - Plot it as a pixel array: PA method combines bifurcation diagrams.
- More exotic: plot bounded-time solutions, e.g. on $[0, \tau]$.
  - Bin the Fourier coefficients for each $a_i$ and $u$: a bin is a range.
  - A pixel corresponds to a class of periodic functions for each $a_i$ and $u$.
  - Turn the pixel on if there is a solution for that sequence of classes.
How to obtain these plots?

How do we obtain these plots? “Good luck!”
How to obtain these plots?

How do we obtain these plots? Develop numerical techniques!

The Pixel Array method says: “plotting is a way of linearizing”.

- Nonlinear solvers traditionally linearize the functions using Jacobians.
- The PA method goes around this, “linearizing” in a different way.
- For PA the challenge is plotting, rather than approximating Jacobians.
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Some methods for plotting $f(x) = 0$:

- Sample points at corner of each cell; turn pixel on if $\text{sign}(f)$ changes.
- Sample one point in each cell and put $2^{-|f(x)|}$ (not Boolean).
- Use “Interval arithmetic” (e.g. $[1, 2] \ast [2, 3] = [2, 6]$).
- If your function is Lipschitz, you can obtain a perfect plot.
- Use your favorite $n$-dimensional contour plotter (with 0-contour).
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- Sample points at corner of each cell; turn pixel on if \( \text{sign}(f) \) changes.
- Sample one point in each cell and put \( 2^{-|f(x)|} \) (not Boolean).
- Use “Interval arithmetic” (e.g. \([1, 2] \star [2, 3] = [2, 6]\)).
- If your function is Lipschitz, you can obtain a perfect plot.
- Use your favorite \( n \)-dimensional contour plotter (with 0-contour).

Once you have the plots, the PA method lets you combine them.
Outline

1. Introduction
2. Details on the Pixel Array method
3. Relation to Numerical methods for PDE
4. Open questions
   - Speed?
   - Accuracy?
   - Clustering?
5. Conclusion
**Apples and oranges**

The inputs and outputs of the PA method are different than Newton’s.

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Our speed test assumes you want to produce “all solutions” in range.

How might we use Newton to find all solutions??
How to compare speed?

Our comparison with Newton assumes you want “all solutions” in range.

- We iterated over all boxes in the grid, each as an initial guess.

\[
\begin{align*}
\cos \left( \ln(z^2 + 10^{-3}x) \right) - x + 10^{-5}z^{-1} &= 0 \\
\cosh(w + 10^{-3}y) + y + 10^{-4}w &= 2 \\
\tan(x + y)(x - 2)^{-1}(x + 3)^{-1}y^{-2} &= 1
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\]
\[
\cosh(w + 10^{-3}y) + y + 10^{-4}w = 2 \quad \text{(Equation 2)}
\]
\[
\tan(x + y)(x - 2)^{-1}(x + 3)^{-1}y^{-2} = 1 \quad \text{(Equation 3)}
\]

- Is this unfair? (Ruben compares to QR vs. inverse iteration.)

With that assumption, for the above case, PA was over 7200x faster.
- PA: 1.5 seconds; Newton: we stopped Julia’s NLsolve after 3 hours.
Determining accuracy

The PA method can introduce false positives, but no false negatives.
- The plotting process can easily introduce false negatives.
- And the Pixel Array method will of course propagate them.
- But assuming plots have true negatives, so do results.

\[ \text{error} := \max_{p \in \text{calculated solution}} \left( \min_{q \in \text{real solution}} \text{dist}(p, q) \right) \]

Fact: for a given resolution, the error is “unbounded”:

For example, take \( f(x, y) = y \), \( g(y, z) = y - \text{step size} \). \( \frac{2}{2} \).

PA returns the “all 1’s” pixel matrix; correct solution is “all 0’s”.

Again, refinement solves this, but at a cost.

\[ 3^\text{Probably either the } \ell_1^- \text{ or the } \ell_\infty^- \text{ norm is best suited for defining distance here.} \]
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We thus define the error to be:\(^3\)

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The source of error

When multiplying arrays, one should think of lining up matching indices.

\[
\begin{array}{ccc}
  y & \times & y \\
  x & \times & z \\
\end{array}
\]

= \[
\begin{array}{ccc}
  Get & \times & z \\
  z & \times & z \\
\end{array}
\]

Note that both of these have small $y$-coverage.

- This is related to having a small value of $\frac{\partial y}{\partial x}$.
- Interestingly, a similar condition comes up in Newton’s method.
The source of error

When multiplying arrays, one should think of lining up matching indices.

\[
\begin{array}{ccc}
\begin{array}{c}
\text{Get}
\end{array} & \times & \begin{array}{c}
\text{Want}
\end{array}
\end{array}
\]

Note that both of these have small \( y \)-coverage.

- This is related to having a small value of \( \frac{\partial y}{\partial x} \).
- Interestingly, a similar condition comes up in Newton’s method.

But errors (false positives) also propagate as we combine arrays.

- A bound for the error in solving the whole system would be useful.
- I’ve got UROPs on the case, but we are looking for a good idea.
Suppose given a cost function for multiplying two arrays.

- For example, my naive cost function was $r \cdot \#\text{variables}$.  
- But if matrices are sufficiently sparse, this goes way down.

Then for any cluster tree, you get the parallel and serial costs.

- Finding a good one can save you a lot of time in multiplying arrays.
- But finding a good one is probably NP.
Clustering cost function

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It remains to find a good algorithm for minimizing this cost function.
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- Such systems come up in PDE when discretizing in space.
- For example, we can plot a pixel array for steady states of each ODE.
- Any networked system of ODS’s can be analyzed using pixel arrays.
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- We have an array for each equation.
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Thanks for inviting me to speak!