

# 18.211 LECTURE 5 9/15

Last Time: Operators acting on  $\mathbb{C}[x_1, \dots, x_n]$

- div. diff operators  $\partial_i: f \mapsto \frac{1}{x_i - x_{i+1}}(1 - s_i)f$
- Demazure operators  $D_i: f \mapsto \frac{f - \frac{x_{i+1}}{x_i} s_i(f)}{1 - \frac{x_{i+1}}{x_i}}$

Let  $X_i: f \mapsto x_i f$

Then  $\boxed{D_i = \partial_i \circ X_i}$  i.e.  $D_i(f) = \partial_i(x_i f)$

For  $w = s_{i_1} \dots s_{i_\ell} \in S_n$  reduced decomp. of  $w$   
 $\partial_w := \partial_{i_1} \dots \partial_{i_\ell}$  and  $D_w := D_{i_1} \dots D_{i_\ell}$

2 formulas for Schur polys

$$S_\lambda(x_1, \dots, x_n) = \partial_{w_0} (x^{\lambda + \delta}) \stackrel{?}{=} D_{w_0} (x^\lambda) \quad \delta = (n-1, n-2, \dots, 0)$$

Thm:  $\partial_{w_0} \circ X_1^{n-1} X_2^{n-2} \dots X_n^0 = D_{w_0}$

Ex.  $n=3$   $w_0 = s_1 s_2 s_1 \in S_{n, n_2}$   
 $\partial_2 \partial_1 \partial_2 X_1 X_1 X_2 \stackrel{?}{=} \overset{\partial_1}{\partial_1 X_1} \overset{\partial_2}{\partial_2 X_2} \overset{\partial_1}{\partial_1 X_1} (= \partial_2 X_2 \partial_1 X_1 \partial_2 X_2)$

Problem:  $\partial_i$ 's don't <sup>always</sup> commute with  $D_j$ 's

Lemma:  $D_i X_j = X_j D_i$  if  $j \notin \{i, i+1\}$

$\partial_i(X_i X_{i+1}) = (X_i X_{i+1}) \partial_i$

Why:  $\partial_i$  only affects  $x_i$  and  $x_{i+1}$ , doesn't affect any others so for other  $j$  we can commute  
 $\partial_i$  is symmetric wrt  $x_i$  and  $x_{i+1}$

Ex.  $\partial_1 X_1 \partial_2 X_2 \partial_1 X_1$   
 $= \partial_1 \partial_2 (X_1 X_2) \partial_1 X_1$   
 $= \partial_1 \partial_2 \partial_1 X_1 X_2 X_1 = \partial_1 \partial_2 \partial_1 X_1^2 X_2$

Proof (for general  $n$ ):

Write  $w_0 = (s_1 s_2 \dots s_{n-1})(s_1 s_2 \dots s_{n-2}) \dots (s_1 s_2)(s_1)$  reduced decomp

Wiring diagram

$$\begin{aligned} D_{w_0} &= (\partial_1 X_1 \partial_2 X_2 \dots \partial_{n-1} X_{n-1}) &= (\partial_1 \partial_2 \dots \partial_{n-1} (X_1 X_2 \dots X_{n-1})) \\ &(\partial_1 X_1 \partial_2 X_2 \dots \partial_{n-2} X_{n-2}) &= (\partial_1 \partial_2 \dots \partial_{n-2} (X_1 X_2 \dots X_{n-2})) \\ &\vdots &\vdots \\ &(\partial_1 X_1 \partial_2 X_2) &= (\partial_1 \partial_2 (X_1 X_2)) \\ &(\partial_1 X_1) &= (\partial_1 X_1) \end{aligned}$$

Can move all  $X$ 's in each line to the right

Now all combined  $X$  terms symmetric wrt operators on the right, so we can commute them to get exactly what we want.



Ex.  $n=3$

$w_0: D_{w_0} = \partial_{w_0} X_1^2 X_2$

$S_i: \partial_1 X_1, S_2: \partial_2 X_2, i=1$

$S_1 S_2: \partial_1 X_1 \partial_2 X_2 = \partial_1 \partial_2 X_1 X_2$

BUT  $S_2 S_1: \partial_2 X_2 \partial_1 X_1 \neq \partial_2 \partial_1 X_2 X_1$

$S_2 S_1 = 312$

Claim: Can write the  $\partial_i$  first then  $X_i$  iff  $w$  is 312 avoiding

Schubert Poly  $G_u = \partial_{u^{-1}w_0}(X^S)$

Demazure char  $ch_{\lambda,w} = D_w(X^\lambda)$

Question: When is  $G_u = ch_{\lambda,w}$ ?

$\lambda = (\lambda_1 \geq \dots \geq \lambda_n), w \in S_n$

not necess. = n

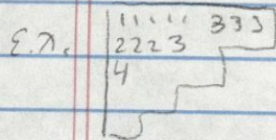
Thm 1 If  $w$  is a 312-avoiding perm in  $S_n$ , then  $ch_{\lambda,w} = G_u$  for some perm  $u \in S_n$

Thm 2 If  $u$  is a 2143-avoiding perm in  $S_m$ , then  $G_u = ch_{\lambda,w}$  for some  $\lambda, w$

↑ vexillary perms

Combinatorial Def of  $S_\lambda$ :

$S_\lambda(x_1, \dots, x_n) = \sum_{T \text{ SSYT of shape } \lambda} x^{\text{weight}(T)}$



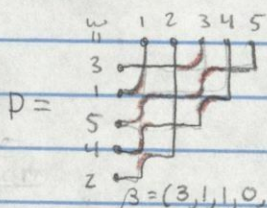
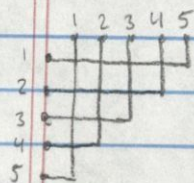
weight = (#1's, #2's, ..., #n's)

Combinatorial formula for  $G_w$

RC-graphs aka pipe dreams

[Forman-Stanley, Billey-Jokushi-Stanley]

Def: Pipe Dream



No double crossings

Any two wires intersect at most once

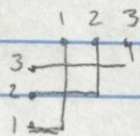
Def:  $\text{weight}(P) = (\beta_1, \beta_2, \dots, \beta_n)$

$\beta_i = \#$  crossings in  $i^{\text{th}}$  row

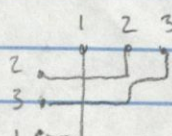
Thm:  $G_w = \sum_{P \text{ pipe dream}} x^{\text{weight}(P)}$

Ex.  $n=3$

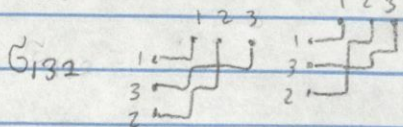
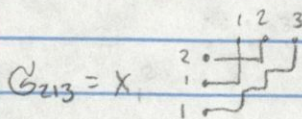
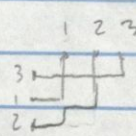
$G_{w_0} = X_1^2 X_2$



$G_{231} = X_1 X_2$



$G_{312} = X_1^2$



two options  $\Rightarrow G_{132} = X_1 + X_2$



Schubert Polynomials  $G_w, w \in S_n$ 

- Defined div. diff. operators  $\partial_i$ :

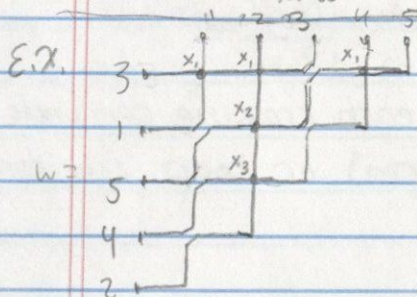
$$(1) G_{w_0} = x_1^{n-1} x_2^{n-2} \dots x_{n-1}$$

$$(2) G_w = \partial_i(G_{ws_i}) \text{ if } l(ws_i) = l(w) + 1$$

- A combinatorial rule for  $G_w$  via pipe dreams

Thrm: [Billey-Jokush-Stanley] cf [Fomin-Stanley]

$$G_w = \sum_{\substack{\text{pipe dream} \\ \text{for } w}} x^{\text{weight}(P)}$$



it should be reduced

(any 2 wires intersect at most once)

$$w = 31542 \quad x^{\text{weight}(P)} = x_1^3 x_2 x_3$$

Nil Coxeter algebra  $N_n$  (over  $\mathbb{C}$ )

- generators  $u_1, u_2, \dots, u_{n-1}$

- relations (1)  $u_i^2 = 0$

$$(2) u_i u_j = u_j u_i, \quad |i-j| \geq 2$$

$$(3) u_i u_{i+1} u_i = u_{i+1} u_i u_{i+1}$$

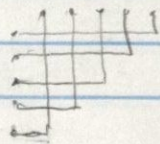
← Note: Same relations as  $\partial_i$ 's

Today: Want to prove Thrm.

Show formula satisfies 2 relation

(1) is "obvious". Just draw pipe dream diagram for  $w_0$ .

(There is exactly 1 way to do this)



→ linear basis of  $N_n$   $u_w, w \in S_n$

$$w = s_{i_1} s_{i_2} \dots s_{i_k} \text{ (reduced decomp)}$$

$$u_w := u_{i_1} u_{i_2} \dots u_{i_k}$$

NOTE: If  $s_{i_1} \dots s_{i_k}$  is not reduced then  $u_{i_1} \dots u_{i_k} = 0$ .

$\forall v, w \in S_n$

$$u_v \cdot u_w = \begin{cases} u_{vw} & \text{if } l(vw) = l(v) + l(w) \\ 0 & \text{otherwise} \end{cases}$$

Commutative variables  $x_1, x_2, \dots, x_{n-1}$  that also commute with  $u_i$ 's

$$h_i(x) := 1 + x u_i$$

$$A_i(x) := h_{n-1}(x) h_{n-2}(x) \dots h_i(x)$$

$$G := A_1(x_1) A_2(x_2) \dots A_n(x_n)$$



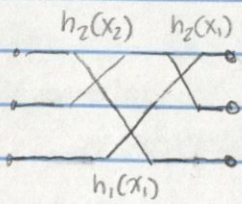
Thrm: [F.S]

← The Schubert poly

$$\tilde{G} = \sum_{w \in S^n} \tilde{G}_w(x_1, x_2, \dots, x_n) U_w$$

E.g.  $n=3 \quad \tilde{G} = A_1(x_1) A_2(x_2)$

$$\begin{aligned} &= h_2(x_1) h_1(x_1) h_2(x_2) \\ &= (1+x_1 U_1)(1+x_1 U_1)(1+x_2 U_2) \\ &= \underbrace{1}_{\tilde{G}_{id}} + \underbrace{x_1 U_1}_{\tilde{G}_{s_1}} + \underbrace{(x_1+x_2) U_2}_{\tilde{G}_{s_2}} + \underbrace{x_1 x_2 U_1 U_2}_{\tilde{G}_{s_1 s_2}} + \underbrace{x_1^2 U_2 U_1}_{\tilde{G}_{s_2 s_1}} + \underbrace{x_1^2 x_2 U_2 U_1 U_2}_{\tilde{G}_{s_2 s_1 s_2}} \end{aligned}$$



For  $w_0$ . And the for each crossing can undo it (corresponds to the 1 term) or keep (corresponds to  $x_i U_j$  term)

Let  $\tilde{G} = \sum_{w \in S_n} \tilde{G}_w(x_1, \dots, x_{n-1}) U_w$   
 WTS  $\tilde{G}_w = G_w$

Need to check (1)  $\tilde{G}_{w_0} = x_1^{n-1} x_2^{n-2} \dots x_{n-1}$

(2)  $\tilde{G}_w = \partial_i(\tilde{G}_{ws_i})$  if  $l(ws_i) = l(w) + 1$

We need to prove the identity

(\*)  $\partial_i(\tilde{G}) = \tilde{G} \cdot U_i$

$$\partial_i(\tilde{G}) = \sum_{w \in S_n} \partial_i(\tilde{G}_w) U_w$$

$$\tilde{G} U_i = \sum_{\substack{w \in S_n \\ l(ws_i) = l(w) + 1}} \tilde{G}_w U_w S_i$$

Proof of (\*)  $\partial_i(A_1(x_1) A_2(x_2) \dots A_i(x_i) A_{i+1}(x_{i+1}) \dots A_{n-1}(x_n))$

NTS =  $A_1(x_1) A_2(x_2) \dots A_i(x_i) A_{i+1}(x_{i+1}) \dots A_{n-1}(x_n) U_i$

BTS  $\partial_i(A_i(x_i) A_{i+1}(x_{i+1})) = A_i(x_i) A_{i+1}(x_{i+1}) U_i$  (\*\*)

Lemma 1: (1)  $h_i(x) h_i(y) = h_i(x+y) \quad h_i(0) = 1$

(2)  $h_i(x) h_j(y) = h_j(y) h_i(x)$  if  $|i-j| \geq 2$

(3)  $h_i(x) h_{i+1}(x+y) h_i(y) = h_{i+1}(y) h_i(x+y) h_{i+1}(x)$

← Yang-Baxter relations

Lemma 2:  $A_i(x)$  &  $A_i(y)$  commute with each other

$A_i(x) A_i(y) = A_i(y) A_i(x)$  (expression symmetric in  $x$  &  $y$ )

Lemma 3:  $A_i(x) A_{i+1}(y) U_i = A_i(y) A_{i+1}(x) U_i$  (" ")

Lemma 4:  $A_i(x) A_{i+1}(y) - A_i(y) A_{i+1}(x) = (x-y) A_i(x) \cdot A_{i+1}(y) \cdot U_i$



Lemma 4  $\Rightarrow (**)$

Ex. of Lemma 2  $n=3, i=2$

$$h_2(x)h_2(y) = h_2(x+y) = h_2(y) + h_2(x)$$

$$\frac{h_2(x)h_1(x)}{A_1(x)} \quad \frac{h_2(y)h_1(y)}{A_1(y)}$$

$$= h_2(x)h_1(x)h_2(y)h_1(y-x)h_1(x-y)h_1(y)$$

$$= h_2(x)h_2(y-x)h_1(y)h_2(x)h_1(x)$$

$$= h_2(y)h_1(y)h_2(x)h_1(x)$$



Schub Polys  $\rightarrow$  div. diff operators  
 $\rightarrow$  pipe dreams

$\lambda = (\lambda_1, \dots, \lambda_k) \subset k \times (n-k)$  rectangle  
 partition

$\rightarrow w(\lambda) \in S_n$  Grassmanian perm.  $w(\lambda) = \lambda_k + 1, \lambda_k + 2, \dots, \lambda_1 + k$

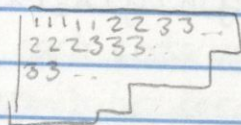
Then  $S_\lambda(x_1, \dots, x_k) = G_{w(\lambda)}$

(all other indices in incr. order)

Comb. Formula for  $S_\lambda(x_1, \dots, x_k)$

Thm:  $S_\lambda(x_1, \dots, x_k) = \sum_{T \in \text{SSYT}(\lambda)} x^{\text{wt}(T)} \leftarrow x_1^{\#1s} x_2^{\#2s} \dots$

Recall: SSYT( $\lambda$ ) = Semi-standard Young tableaux  
 of shape  $\lambda$



Pipe dream formula for  $G_w$   $\rightarrow$  SSYT formula for  $S_\lambda$

(Want to see SSYT formula as special case of pipe dreams)

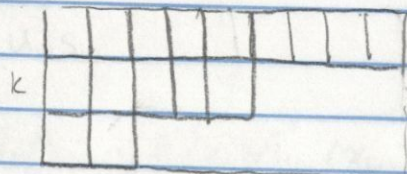
Claim: Pipe dreams  $P$  for  $w(\lambda)$  are in bij. with SSYT( $\lambda$ )

Rule: Order the wires of  $P$  from top to bottom by their left end

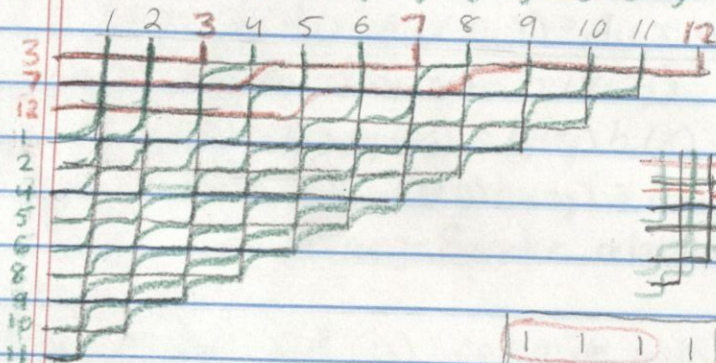
For  $i = 1, \dots, k$ ,

# crossings of  $i^{\text{th}}$  wire of  $P$  in  $j^{\text{th}}$  row = # entries  $k+1-j$  in  $(k+1-i)^{\text{th}}$  row of  $T$

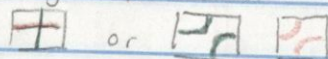
Ex:  $n=12, k=3$   
 $\lambda = (9, 5, 2)$



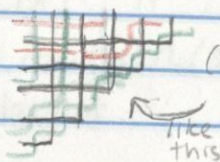
$w(\lambda) = 3, 7, 12, 1, 2, 4, 5, 6, 8, 9, 10, 11, 12$



each grid point either



(Should be drawn inside boxes, oops)



like this

1	1	1	1	2	2	3	3	3
2	2	2	3	3				
3	3							

Always get

have choices

(Calculate rows from bottom up)



- Thm:  $S_\lambda(x_k, x_{k-1}, \dots, x_1) = \sum_{\tau \in \text{SSYT}(\lambda)} x^{\text{wt}(\tau)}$
- Lemma:  $S_\lambda = (x_1, \dots, x_k) = G_{w(\lambda)}$  is symmetric in  $x_1, \dots, x_k$
- Lemma:  $G_w$  is symmetric in  $x_i$  &  $x_{i+1}$  if  $l(ws_i) = l(w) + 1$

### Double Schubert Polynomials

$$G_w(x, y) := G_w(x_1, \dots, x_{n-1}, y_1, \dots, y_{n-1}) \quad w \in S_n$$

Almost do the same as for usual Schub

$$(1)' \quad G_{w_0}(x, y) = \prod_{\substack{(i,j) \\ i,j \geq 1 \\ i+j \leq n}} (x_i - y_j)$$

$$(2) \quad G_w(x, y) = \partial_i^* G_{ws_i}(x, y) \quad \text{if } l(ws_i) = l(w) + 1$$

$$\partial_i^* : f \mapsto \frac{1}{x_i - x_{i+1}} (f(x_1, \dots, \overset{\leftarrow}{x_i}, x_{i+1}, \dots) - f(x_1, \dots, x_{i+1}, \overset{\leftarrow}{x_i}, \dots))$$

Prop: Positivity:  $G_w(x_1, \dots, x_{n-1}, -y_1, -y_2, \dots, -y_{n-1})$  has positive int. coeffs.

### Pipe dream formula for $G_w(x, y)$

$$\text{Thm: } G_w(x, y) = \sum_{\substack{P \text{ pipe dream} \\ \text{for } w}} \prod_{(i,j) \text{ st. } P \text{ has a} \\ \text{crossing in the } i^{\text{th}} \text{ row} \\ \text{ \& } j^{\text{th}} \text{ column}} (x_i - y_j)$$

$$\text{Cor (Symmetry): } G_w(x_1, \dots, x_{n-1}, y_1, \dots, y_{n-1}) \\ = G_w(-y_1, \dots, -y_{n-1}, -x_1, \dots, -x_{n-1})$$

### Algebra (over $\mathbb{C}$ )

generators:  $u_1, \dots, u_{n-1}, x_1, \dots, x_{n-1}, y_1, \dots, y_{n-1}$

relations:  $x_i, y_j$  commute with each other and  $u_k$ 's

$$(1)' \quad u_i^2 = 0$$

$$(2) \quad u_i u_j = u_j u_i \quad |i-j| \geq 2$$

$$(3) \quad u_i u_{i+1} u_i = u_{i+1} u_i u_{i+1}$$

$h_i(a) = 1 + a u_i$  satisfy YB relations  
Yang-Baxter

$$G^{x,y} := \prod_{\substack{(i,j) \\ i,j \geq 1 \\ i+j \leq n}} h_{i+j-1}(x_i - y_j)$$

order  $(i,j)$ 's as  $(1, n-1)(1, n-2) \dots (1, 1)$   
 $(2, n-2)(2, n-3) \dots (2, 1)$   
 $(3, n-3) \dots (3, 1)$   
 $\vdots$   
 $(n-1, 1)$



# LECTURE 9 9/27

Cauchy identity for  $S_\lambda$

$$\sum_{\lambda=(\lambda_1, \dots, \lambda_n)} S_\lambda(x_1, \dots, x_n) S_\lambda(y_1, \dots, y_n) = \prod_{i,j \in [n]} \frac{1}{1-x_i y_j}$$

Let  $x^\alpha y^\beta = x_1^{\alpha_1} \dots x_n^{\alpha_n} y_1^{\beta_1} \dots y_n^{\beta_n}$

LHS:  $[x^\alpha y^\beta] = \# \{ (P, Q) \mid P, Q \text{ SSYT of shape } \lambda \text{ s.t. } \begin{matrix} \alpha = \text{wt}(P) \\ \beta = \text{wt}(Q) \end{matrix} \}$

RHS:  $\prod_{i,j} \frac{1}{1-x_i y_j} = \prod_{i,j} \sum_{a_{ij} \geq 0} (x_i y_j)^{a_{ij}} = \sum_{\substack{A=(a_{ij}) \\ \text{nonneg-int} \\ n \times n \text{ matrix}}} \prod_i x_i^{\sum_j a_{ij}} \prod_j y_j^{\sum_i a_{ij}}$

Robinson-Schensted-Knuth correspondence (RSK)

Robinson-Schensted correspondence

special case where  $\alpha = \beta = (i_1, i_2, \dots, i_r)$

$P, Q$  are SYT's of the same shape

$A$  - permutation matrix

$$S_n \xleftrightarrow{\text{RS}} \{ (P, Q) \mid \text{SYT's of same shape } \lambda \vdash n \}$$

$w = w_1 \dots w_n$

Ex Example of RSK  $n=2$

$$A = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \xleftrightarrow{\text{RSK}} (P, Q)$$

$$P = \begin{array}{|c|c|} \hline 1 \dots 1 & 2 \dots 2 \\ \hline 2 \dots 2 & \\ \hline \end{array} \begin{matrix} \lambda_1 \\ \lambda_2 \end{matrix}$$

$$Q = \begin{array}{|c|c|} \hline 1 \dots 1 & 2 \dots 2 \\ \hline 2 \dots 2 & \\ \hline \end{array} \begin{matrix} \lambda_1 \\ \lambda_2 \end{matrix}$$

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \longleftrightarrow \begin{pmatrix} \lambda_2 \leq e \\ f \leq \lambda_1 \end{pmatrix}$$

col. sum  $a+c=e$

row sums  $a+b=f$

$$b+d = \lambda_1 + \lambda_2 - e$$

$$c+d = \lambda_2 + \lambda_1 - f$$

These 4 e.g. are not lin. independent  $\Rightarrow$  Not enough to solve for RSK  
 Answer:  $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \longleftrightarrow \begin{pmatrix} \min(b,c) & a+c \\ a+b & a+\min(b,c) \end{pmatrix}$  correspondence

In general can view RSK as piecewise linear transform of matrices



# Classical Construction of RSK

Ex.

$$A = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

→ biword  $w$  with  $a_{ij}$  entries  $\binom{i}{j}$  arranged lexicographically

$$w = \binom{2}{1} \binom{2}{1} \binom{2}{1} \binom{2}{4} \binom{3}{3}$$

$\uparrow$                        $\uparrow$                        $\uparrow$   
 1 in pos.              2 is pos              1 in pos  
 (2,1)                      (2,1)                      (2,4)

will insert entries into  $P$  &  $Q$  using

## Schnested insertion algorithm

Ex.  $P = \begin{bmatrix} 1 & 1 & 1 & 3 & 3 & 8 \\ 2 & 2 & 2 & 4 & 5 \\ 3 & 4 & 6 & 6 & 7 \\ 5 \end{bmatrix} \leftarrow j=2$

1st entry in 1st row which  $> j$

1st entry in 2nd row which  $> 3$

1st entry in 3rd row which  $> 4$

$$\begin{bmatrix} 1 & 1 & 1 & 2 & 3 & 3 & 8 \\ 2 & 2 & 2 & 3 & 5 \\ 3 & 4 & 4 & 6 & 7 \\ 5 & 6 \end{bmatrix}$$

Back to Ex. from top

	$\binom{1}{2}$	$\binom{2}{1}$	$\binom{2}{1}$	$\binom{2}{4}$	$\binom{3}{3}$	$\binom{i}{j}$
$P$	$\emptyset$	insert 2	insert 2	insert 2	insert 2	insert 2
	$\begin{bmatrix} 1 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 & 2 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 & 2 & 4 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 \end{bmatrix}$	
$Q$	$\emptyset$	insert 1	insert 1	insert 1	insert 1	insert 1
	$\begin{bmatrix} 1 \end{bmatrix}$	$\begin{bmatrix} 1 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 & 1 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 & 1 & 2 \\ 2 \end{bmatrix}$	$\begin{bmatrix} 1 & 2 & 2 \\ 2 & 3 \end{bmatrix}$	

- Insert  $j$  to  $P$  using Schnested insertion
- Insert  $i$  to  $Q$  in whatever box got added to  $P$

Thm: (Knuth) This construction  $A \mapsto (P, Q)$  is bij. with needed properties.

Proof: Bijection: Can undo algorithm.

Idea

Find max entry in  $Q$  tableaux. That tells where we added in  $P$  tableau  
 un-add via schnested in reverse



Thrm: (Knuth) If  $A \xrightarrow{RSK} (P, Q)$ , then  $A^T \xrightarrow{RSK} (Q, P)$

Not so clear why we should have this symmetry from this version of RSK

Thrm: (Knuth) If  $\lambda$  is the shape of  $P \& Q$  tableaux, then

$\lambda_1$  = length of a "max weakly increasing subsequence" in  $w$

$\lambda'_1$  = length of a "max strictly decreasing subseq." in  $w$

Def: **Weakly increasing subseq.**

$w = \binom{i_{a_1}}{j_{a_1}} \binom{i_{a_2}}{j_{a_2}} \dots \binom{i_{a_x}}{j_{a_x}}$  st.

$i_{a_1} \leq i_{a_2} \leq \dots \leq i_{a_x}$

$j_{a_1} \leq j_{a_2} \leq \dots \leq j_{a_x}$

**Strictly decreasing subseq**

$i_{a_1} < \dots < i_{a_x}$

$j_{a_1} > \dots > j_{a_x}$

Ex. In example before,  $\lambda_1 = 3$ ,  $\lambda'_1 = 2$

A different correspondence that lets us see Thrm properties better

SSYT  $\longleftrightarrow$  Gelfand-Tsetlin Patterns

Ex.  
 $n=5$

1	1	1	2	2	<del>5</del>	<del>5</del>
2	2	2	3	3	3	
3	3	<del>4</del>	<del>4</del>	<del>5</del>	<del>5</del>	
<del>5</del>						

$\rightsquigarrow$

8	6	6	2	0
6	6	4	0	
6	6	2		
6	3			
4				

- Write # elts in each row of T
- Cross out all of highest number
- Repeat for next row of GT-pattern