# Extremal Bounds on Peripherality Measures

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### Conventions

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Finite simple graphs, usually connected. The name of the graph is always G and the number of vertices is always denoted n.

## Centrality and Peripherality

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### Definition

A peripherality measure is the opposite of a centrality measure; peripheral vertices are the least important in a graph.

# Applications

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• Atmospheric networks

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- Neural networks

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The peripherality of a graph is the sum of the peripheralities of its vertices.

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$$n_{G}(v_{1}, v_{2}) < n_{G}(v_{2}, v_{1}) n_{G}(v_{1}, v_{3}) < n_{G}(v_{3}, v_{1}) n_{G}(v_{1}, v_{4}) < n_{G}(v_{4}, v_{1}) n_{G}(v_{3}, v_{2}) < n_{G}(v_{2}, v_{3}) n_{G}(v_{4}, v_{2}) < n_{G}(v_{2}, v_{4}) n_{G}(v_{3}, v_{4}) = n_{G}(v_{4}, v_{3}) peri(v_{1}) = 0$$

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$$\begin{array}{l} n_G(v_1, v_2) < n_G(v_2, v_1) \\ n_G(v_1, v_3) < n_G(v_3, v_1) \\ n_G(v_1, v_4) < n_G(v_4, v_1) \\ n_G(v_3, v_2) < n_G(v_2, v_3) \\ n_G(v_4, v_2) < n_G(v_2, v_4) \\ n_G(v_3, v_4) = n_G(v_4, v_3) \\ \text{peri}(v_2) = 3 \end{array}$$

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#### Theorem

The peripherality of a graph is the number of unordered pairs (v, x) of vertices such that  $n_G(v, x) \neq n_G(x, v)$ .

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#### Corollary

The peripherality of an *n*-vertex graph is at most  $\binom{n}{2}$ .

Geneson and Tsai found constructions of the equality case for each  $n \ge 9$ . We determined the maximum for each n < 9.

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#### Definition

The edge peripherality of an edge, denoted eperi( $\{u, v\}$ ), is the number of vertices x such that  $n_G(x, u) > n_G(u, x)$  and  $n_G(x, v) > n_G(v, x)$ .

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#### Old bound (Geneson and Tsai)

The maximum edge peripherality of an *n*-vertex graph lies in the interval  $[\frac{2}{125}n^3, \frac{1}{2}n^3]$ .

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#### New bound

The maximum edge peripherality of an *n*-vertex graph lies in the interval  $\left[\frac{\sqrt{3}}{24}n^3(1-o(1)), \frac{1}{6}n^3\right]$ .

The edge sum peripherality of an edge, denoted  $espr({u, v})$ , is defined as

$$\sum_{v\in V-\{u,v\}}(n_G(x,u)+n_G(x,v)).$$

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The maximum edge sum peripherality of an *n*-vertex graph of diameter 2 is  $\frac{4}{27}n^4 - O(n^3)$ .

#### New bound

The maximum edge sum peripherality of an *n*-vertex bipartite graph of diameter at most 3 is  $\frac{1}{8}n^4 - O(n^2)$ .

The Trinajstić index of an unordered pair (u, v) of vertices is  $NT(u, v) = (n_G(u, v) - n_G(v, u))^2$ .

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#### Definition

The Trinajstić index of a graph is the sum of NT(u, v) over all  $\binom{n}{2}$  pairs.

### Conjecture (Furtula)

For sufficiently large *n*, the Trinajstić index of an *n*-vertex graph is maximized by the generalization of this graph:



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#### Verdict

The conjecture is false. This family of graphs achieves  $NT(G) \le 0.25n^4(1 + o(1))$ . The maximum of NT(G) is actually  $0.5n^4(1 - o(1))$ .

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As the number of "arms" and the length of each "arm" both go to infinity,  $NT(G) = 0.5n^4(1 - o(1))$ .

### Conjecture (Furtula)

For sufficiently large n, the Trinajstić index of an n-vertex tree is minimized by the generalization of this graph:



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This conjecture is still open.

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If the Trinajstić Index of a graph is 0, then every vertex has the same degree.

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#### Verdict

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In fact, these can be used to generate arbitrarily large counterexamples.

I thank Dr. Jesse Geneson for suggesting this research topic, telling me about many possible directions for research, helping me format my results into a paper, giving me feedback on drafts of the paper, helping me submit it to arXiv and a journal, and giving me feedback on my presentation rehearsal. I thank Dr. Tanya Khovanova for giving me feedback on drafts of the paper and on my presentation, as well as PRIMES organizers for making this amazing research opportunity possible.

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