Group messaging in the XRD private communication system

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Acknowledgements

Thank you to Albert Kwon for mentoring me

Thank you to Prof. Devadas for PRIMES-CS

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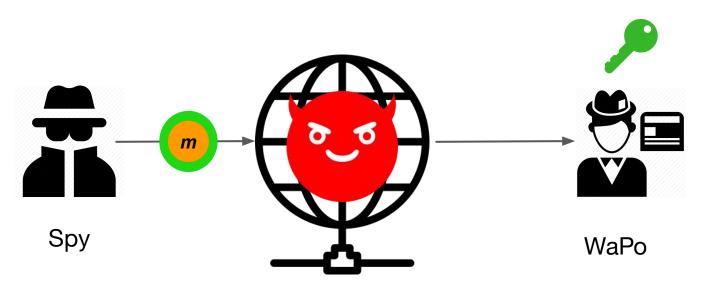
Thank you to my parents for their support

Motivation

Spy hides message *content* through encryption.

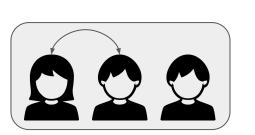
However, spy still leaks metadata:

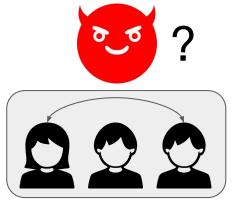
- Identities
- Timing
- Size



Privacy guarantee

 Provide metadata private messaging against powerful adversaries

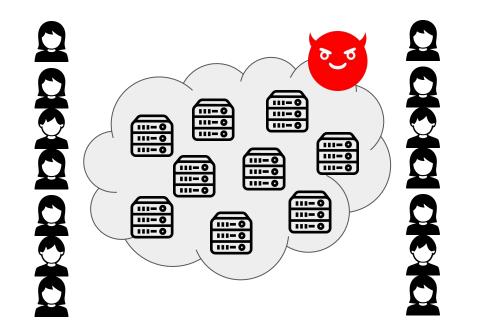




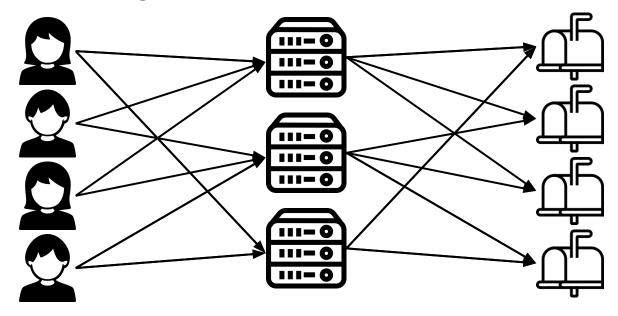


Deployment and threat model

• Global network adversary



XRD: basic design



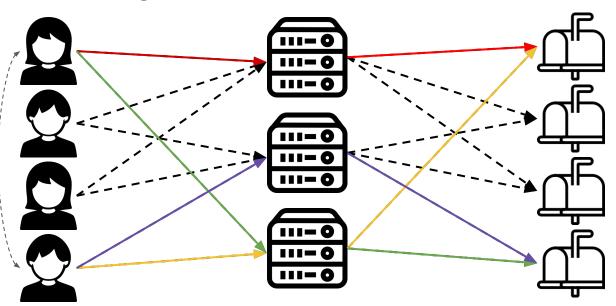
1. Send messages to *l* servers

2. Forward messages to mailboxes

 $\ell = 2$

XRD: basic design

If 1 and 4 are talking to each other with $\ell = 2$

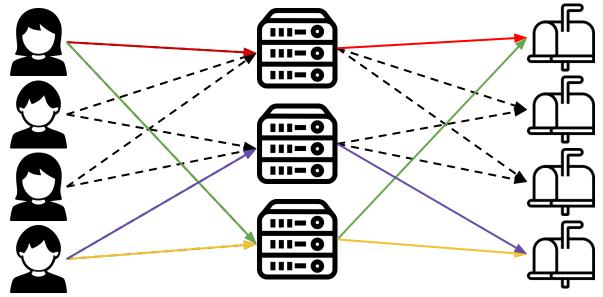


1. Send messages to *l* servers

2. Forward messages to mailboxes

XRD: basic design

If 1 and 4 are *not* talking to each other with $\ell = 2$

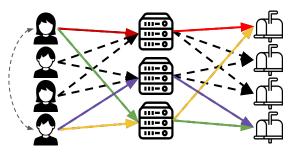


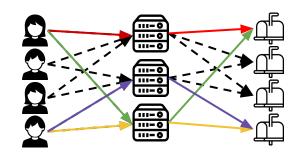
1. Send messages to *l* servers

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Security argument

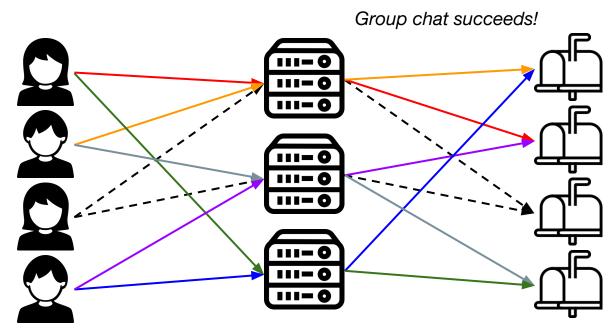
- Every mailbox gets exactly *l* messages
 - Mailboxes are identical
- Every pair of users intersects
 - Hides which users are talking with each other
- The origin of the message is hidden by the server
 - Hides swap-or-not





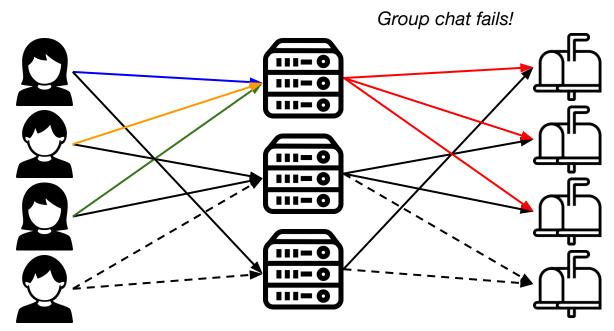
Group messaging

• Multiple 1-1 conversations



Group messaging

• Multiple 1-1 conversations



Goal

- Enable group conversations of any size
- Minimize the max number of messages that a server gets
 This determines overall latency of XRD

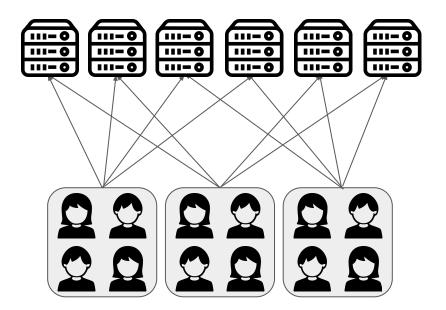
Approaches

- Randomized algorithm
- Parallel algorithm
- Linear programming
- Mixed integer programming
- Greedy algorithms (this talk)

Greedy algorithm #1

- **Goal:** Every pair of users intersects at least x times
- Divide into (ℓ /x+1) groups
- $l = \sqrt{2nx + x^2/4} x/2$ minimizes load per server (for n servers)

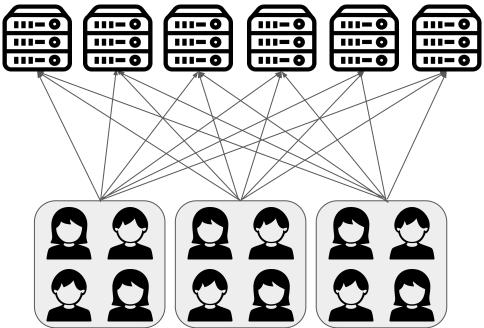
6 servers, x = 23 groups, $\ell = 4$



Greedy algorithm #2

- Start with every user sending message to every server
- Remove messages that least impact intersections between users

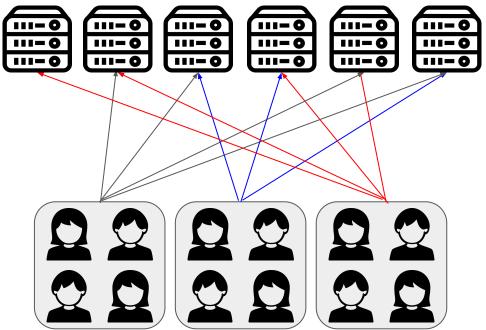
6 servers, x = 2, 3 groups



Greedy algorithm #2

- Start with every user sending message to every server
- Remove messages that least impact intersections between users

6 servers, x = 2, 3 groups



Results

Max number of messages per server vs. group size (100 servers, 1M users)



Size of group

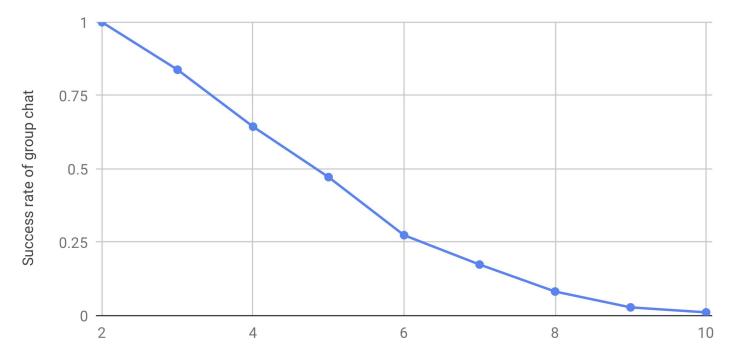
Conclusion

- Two greedy algorithms to enable group messaging in XRD
- For 100 servers and 1M users:
 - Groups of up to 5 users with 96% latency increase
 - Groups of up to 20 users with 250% latency increase

Backup

Success rate of current algorithm

Success rate of group chat vs. group size (100 servers, 1M users)



Size of group

Scalability properties of original system

For *m* users and *n* chains,

- We can make sure all users intersect with $l = \sqrt{2n}$
- Each chain handles m^{*}ℓ/n = (√2)^{*}m/(√n) messages
 If you increase n, the load per chain goes down (scalable)