

1

(a) $C_{jk} = 2^{-p}$, where $p = \#$ of persons between person j and k .

This is because each communication loses a factor of $1/2$ and this is multiplicative.

(b) $C = \begin{bmatrix} 1 & 1/2 & 1 & 1/2 \\ 1/2 & 1 & 1/2 & 1 \\ 1 & 1/2 & 1 & 1/2 \\ 1/2 & 1 & 1/2 & 1 \end{bmatrix}$

(c) $(F^{-1}C)e_1 = (DF^{-1})e_1 = D(F^{-1}e_1)$

Now, the 1st column of F^{-1} is $\begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$

$\Rightarrow (F^{-1}C)e_1 = D \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix} = \begin{pmatrix} d_1 \\ \vdots \\ d_n \end{pmatrix} = \text{diagonal entries}$

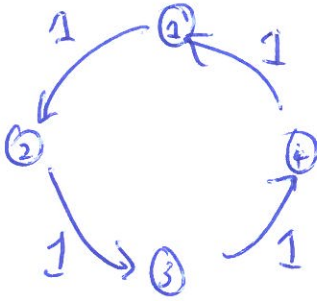
(d) Rank 1: $\begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$ Rank 4: $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$

Rank 2: $\begin{bmatrix} 1 & 1/2 & 1 & 1/2 \\ 1/2 & 1 & 1/2 & 1 \\ 1 & 1/2 & 1 & 1/2 \\ 1/2 & 1 & 1/2 & 1 \end{bmatrix}$ Rank 3: $\begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix}$

(e) Take $\begin{bmatrix} | & | & | & | \\ | & | & | & | \\ | & | & | & | \\ | & | & | & | \end{bmatrix}$ (for example.)

• ~~Everyone talk~~

Corresponding to:



2

(a) $\sin(4\pi t) = \cos(4\pi t + \pi/2)$

$f(t) = \cos(2\pi(2)\pi t + \pi/2) + \cos(\overset{2\pi(3)t}{\cancel{6\pi t}} + \pi/2)$

$\therefore A_2 = 1, a_2 = \pi/2, A_3 = 1, a_3 = \pi/2$

(b) ~~∗~~ Nyquist's sampling rate: need at least 2 pts per wavelength to exactly recover a signal.

(c) Points are $k/4$ $0 \leq k \leq 4$. AND cosine 2π -periodic:

so $\cos(6\pi k/4 + \pi/2) = \cos(6\pi k/4 + \pi/2 + 2\pi Lk)$ for any integer k .
 $= \cos(2(3+4L)\pi k/4 + \pi/2)$

$M = 3 + 4L$ for any integer L .

e.g. $M = -9, -5, -1, 3, 7, \dots$

(d) $\sin(4\pi k/4) = \sin(k\pi) = 0$ $0 \leq k \leq 4$.

$f(k/4) = \cos(6\pi k/4 + \pi/2) = \cos(-2\pi k/4 + \pi/2)$
 $= \cos(2\pi k/4 - \pi/2)$

Recover signal is $\cos(2\pi t - \pi/2)$, i.e. $A_0 = 0, a_0 = 0$
 $A_1 = 1, a_1 = -\pi/2$.

~~(d) Nyquist's sampling rate: Need at least two points samples per wavelength to exactly recover a signal.~~

e) Easy question : Since $f(x)$ is periodic, sample f at k/N , $1 \leq k \leq N$ instead.

Procedure :

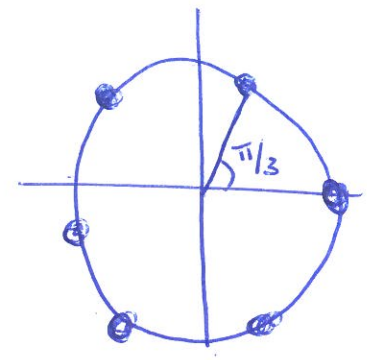
- ① Sample f at ~~k/N~~ k/N , $1 \leq k \leq N$.
- ② ~~Apply~~ Apply FFT to the vector $f(1), f(1/N), f(2/N), \dots, f((N-1)/N)$.
- ~~③~~ N must be large enough. Here, $N \geq 7$,
(anything more than 7 is fine)

3

Give the DFT for $N=6$:

$$X_k = \sum_{n=0}^5 c_n e^{-2\pi i n k / 6} \quad , \quad 0 \leq k \leq 5$$

$w = e^{-2\pi i / 6} = \frac{1}{2} + \frac{\sqrt{3}}{2} i$ so equivalent to



In matrix form

$$\begin{bmatrix} X_0 \\ X_1 \\ X_2 \\ X_3 \\ X_4 \\ X_5 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 & w^4 & w^5 \\ 1 & w^2 & w^4 & w^6 & w^8 & w^{10} \\ 1 & w^3 & w^6 & w^9 & w^{12} & w^{15} \\ 1 & w^4 & w^8 & w^{12} & w^{16} & w^{20} \\ 1 & w^5 & w^{10} & w^{15} & w^{20} & w^{25} \end{bmatrix} \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \\ c_4 \\ c_5 \end{bmatrix}$$

DECOMPOSE POLYNOMIAL :

$$P_5(z) = \sum_{n=0}^5 c_n z^n = c_0 + c_1 z + c_2 z^2 + c_3 z^3 + c_4 z^4 + c_5 z^5$$

$$= \underbrace{c_0 + c_2 z^2 + c_4 z^4}_{\text{even}(z^2)} + z \underbrace{(c_1 + c_3 z^2 + c_5 z^4)}_{\text{odd}(z^2)}$$

Equivalently, one could do:

$$P_5(z) = \underbrace{c_0 + c_3 z^3}_{P_1(z^3)} + z \underbrace{(c_1 + c_4 z^3)}_{P_2(z^3)} + z^2 \underbrace{(c_2 + c_5 z^3)}_{P_3(z^3)}$$

MAIN IDEA : The DFT of size $N=6$ can be broken up into 2 DFTs of size $N=3$.

EQUIVALENTLY : THE DFT of size $N=6$ can be broken up into 3 DFTs of size $N=2$.

MAIN IDEA needs the fact that

$$\left(e^{-2\pi i n k / 6}\right)^2 = e^{-2\pi i n k / 3}$$

ROOTS OF UNITY
OF $N=3$
(REPEATED TWICE)

EQUIVALENTLY:

$$\left(e^{-2\pi i n k / 6}\right)^3 = e^{-2\pi i n k / 2}$$

ROOTS OF UNITY
OF $N=2$
(REPEATED 3x)

\Rightarrow $p_{\text{even}}(z_k^2)$, $p_{\text{odd}}(z_k^2)$, ~~$p_{\text{even}}(z_k)$~~ are DFTs

Equivalently: $p_1(z_k^3)$, $p_2(z_k^3)$, and $p_3(z_k^3)$ are DFTs

Full credit also if they go for the matrix factorization:

$$\begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & \omega & \omega^2 & \omega^3 & \omega^4 & \omega^5 \\ 1 & \omega^2 & \omega^4 & \omega^3 & \omega^2 & \omega^6 \\ 1 & \omega^3 & \omega^6 & \omega^9 & \omega^{12} & \omega^5 \\ 1 & \omega^4 & \omega^8 & \omega^{12} & \omega^6 & \omega^{20} \\ 1 & \omega^5 & \omega^{10} & \omega^{15} & \omega^{20} & \omega^{25} \end{bmatrix} = \underbrace{\begin{bmatrix} I_3 & I_3 \\ I_3 & -I_3 \end{bmatrix}}_{\text{combine}} \underbrace{\begin{bmatrix} 1 & & & & & \\ & 1 & & & & \\ & & \omega & & & \\ & & & \omega^2 & & \\ & & & & \omega & \\ & & & & & \omega^2 \end{bmatrix}}_Z \underbrace{\begin{bmatrix} 1 & & & & & \\ & \omega^2 & & & & \\ & & \omega^4 & & & \\ & & & \omega^2 & & \\ & & & & \omega^4 & \\ & & & & & \omega^2 \end{bmatrix}}_{\text{DFT of size } N=3}$$

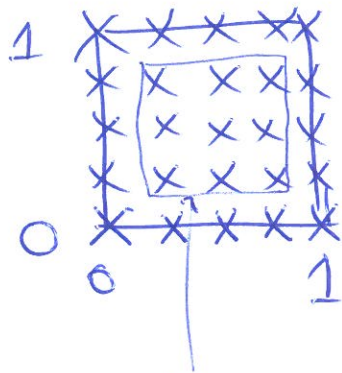
EQUIVALENTLY:

$$= \begin{bmatrix} I_3 & I_3 & I_3 \\ I_3 & \omega^2 I_3 & \omega^4 I_3 \\ I_3 & \omega^4 I_3 & \omega^8 I_3 \end{bmatrix} \begin{bmatrix} 1 & & & & & \\ & 1 & & & & \\ & & \omega & & & \\ & & & \omega^2 & & \\ & & & & \omega & \\ & & & & & \omega^2 \end{bmatrix} \begin{bmatrix} 1 & & & & & \\ & 1 & & & & \\ & & \omega & & & \\ & & & \omega^2 & & \\ & & & & \omega & \\ & & & & & \omega^2 \end{bmatrix}$$

(b) $N=7$ is a prime number \Rightarrow no factorization.

4

(a)



5x5 grid.

Edge values are known to be 0.

X is 3x3 and contains the values of solution at these grid points.

(b)

MX represents yu . So take y on the grid
i.e. ~~1/4~~ $1/4, 1/2, 3/4$
(excluding endpoints)

~~$MX = u$~~ If $X =$ values of u on grid, then
 $MX =$ values of yu on grid.

(c)

$$K_n X + X(SDS^{-1}) + MX = G$$

$$K_n [XS] + [XS]D + M[XS] = GS$$

(multiply
S on the
right)

Setting $Y = XS$ we have

$$K_n Y + YD + MY = GS$$

$$\Rightarrow (K_n + M)Y + YD = GS$$

(d)

$$YD = \begin{bmatrix} d_1 y_{11} & d_2 y_{12} & \dots & d_n y_{1n} \\ d_1 y_{21} & d_2 y_{22} & \dots & d_n y_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ d_1 y_{n1} & d_2 y_{n2} & \dots & d_n y_{nn} \end{bmatrix} = \left[d_1 y_{i1} \mid \dots \mid d_n y_{in} \right]$$

~~$k_n + M$~~
 ∴ To solve for k^{th} column of Y we have

$$(k_n + M) y_k + d_k y_k = \underbrace{b_k}_{\substack{k^{\text{th}} \text{ column} \\ \text{of } GS}}$$

$$\underbrace{(k_n + M + d_k I_n)}_{\text{tridiagonal}} y_k = b_k$$

(e)

- ① Compute GS $O(n^2 \log n)$
- ② Solve $\underbrace{(k_n + M + d_k I_n)}_{\text{tridiagonal}} y_k = b_k \quad 1 \leq k \leq n$ $O(n^2)$
- ③ Compute $X = Y S^{-1}$ $O(n^2 \log n)$

$$O(n^2 \log n)$$