

On the Winning Strategies in Generalizations of Nim

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An example: Nim

- Take at least one cookie from any one pile
- The player who takes the last cookie wins

Player 1



Player 2

An example: Nim

- Take at least one cookie from any one pile
- The player who takes the last cookie wins

Player 1



$(2,2,1)$

Player 2

An example: Nim

- Take at least one cookie from any one pile
- The player who takes the last cookie wins

Player 1



$(2,2,0)$

Player 2



An example: Nim

- Take at least one cookie from any one pile
- The player who takes the last cookie wins

Player 1



(1,2,0)

Player 2



An example: Nim

- Take at least one cookie from any one pile
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Player 1



$(1,1,0)$

Player 2



An example: Nim

- Take at least one cookie from any one pile
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Player 2



$(0,1,0)$

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Player 2

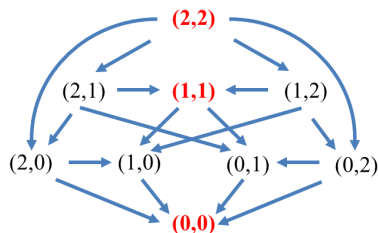


$(0,0,0)$

- Player 2 wins

P-Positions

- The starting position $(2, 3, 1)$ is one where the person to play will always lose assuming optimal play
- We call such positions P-positions (losing positions)
- All other positions are called N-positions (winning positions)
 - Moves from P-positions can only go to N-positions
 - At least one move from every N-position goes to a P-position
 - The zero position $(0, \dots, 0)$ is a P-position
- Winning strategy is to move to a P-position



Winning Strategy for Nim

Theorem (Bouton's Theorem)

In Nim, $P = (a_1, \dots, a_n) \in \mathcal{P}$ if and only if $\bigoplus_{i=1}^n a_i = 0$.

- The operator \oplus is the bitwise XOR operator, (nim-sum) – represent each of the numbers in binary and add them column-wise modulo 2.

Another example: Wythoff's Game

- Take same number of cookies from two piles or any number from one pile

Player 1



Player 2



(4,3)

Another example: Wythoff's Game

- Take same number of cookies from two piles or any number from one pile

Player 1

Player 2



Another example: Wythoff's Game

- Take same number of cookies from two piles or any number from one pile

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Player 2



$(1,1)$



Another example: Wythoff's Game

- Take same number of cookies from two piles or any number from one pile

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Player 2



$(0,0)$

Another example: Wythoff's Game

- Take same number of cookies from two piles or any number from one pile

Player 1



Player 2



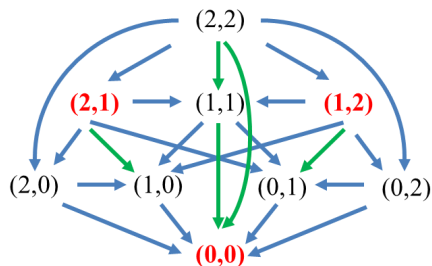
(0,0)

- Player 1 wins

Winning Strategy for Wythoff

Theorem (Wythoff's Theorem)

In Wythoff's game, $P = (a_1, a_2) \in \mathcal{P}$ if and only if $\{a_1, a_2\} = \{\lfloor n\phi \rfloor, \lfloor n\phi^2 \rfloor\}$ for some integer n , where $\phi = \frac{1+\sqrt{5}}{2}$.

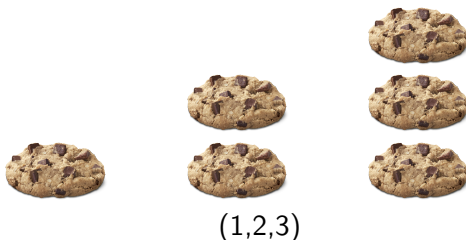


Rectangular Games

- Move consists of taking same number of cookies from specified subsets of piles
 - Based on Cookie Monster game
- Adjoins rules onto the Nim rule (taking at least one cookie from exactly one of the piles)
- We are interested in the properties of the P-positions

Proposed Generalizations

- We specify which subsets of piles are legal to take from:
- Examples of games with three piles
 - 1 One or n game
 - 2 One or Two game
 - 3 Consecutive game
 - 4 Cookie Monster game



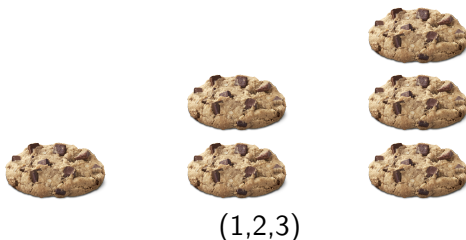
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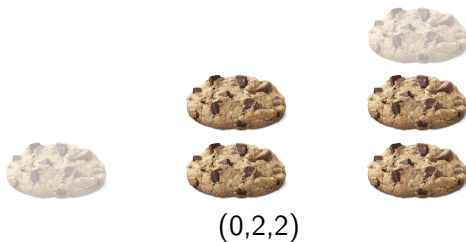
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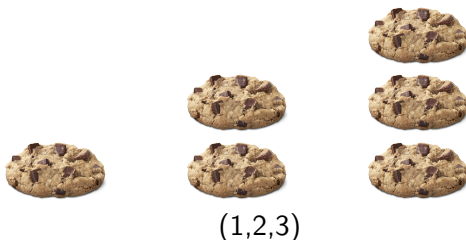
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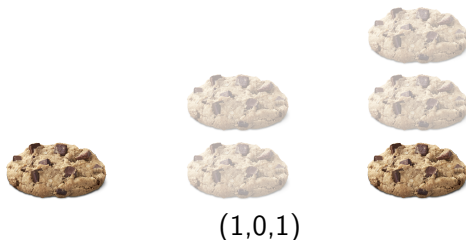
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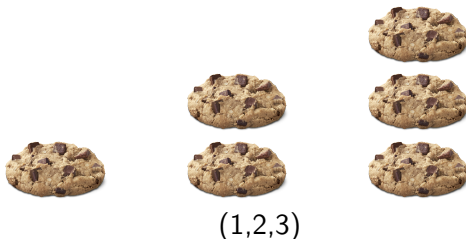
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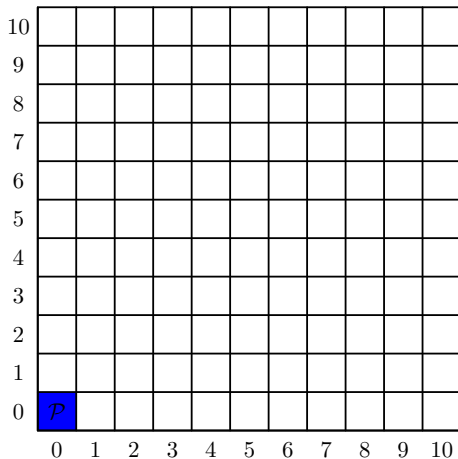
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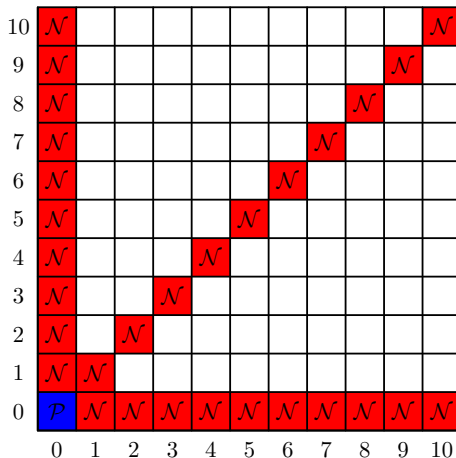
Methods / Programs

- We implement an efficient recursive algorithm for computing P-positions in Java



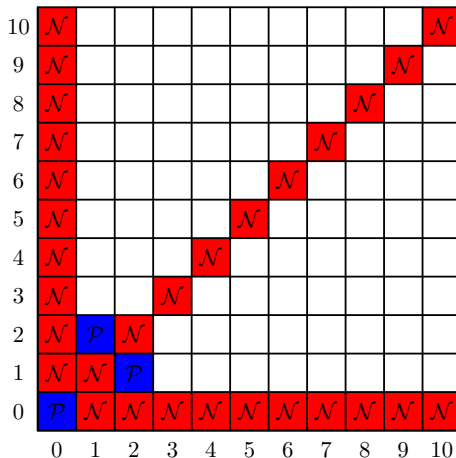
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Methods / Programs

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10	\mathcal{N}	\mathcal{N}	\mathcal{N}						\mathcal{N}	\mathcal{N}	
9	\mathcal{N}	\mathcal{N}	\mathcal{N}					\mathcal{N}	\mathcal{N}	\mathcal{N}	
8	\mathcal{N}	\mathcal{N}	\mathcal{N}				\mathcal{N}	\mathcal{N}	\mathcal{N}		
7	\mathcal{N}	\mathcal{N}	\mathcal{N}				\mathcal{N}	\mathcal{N}	\mathcal{N}		
6	\mathcal{N}	\mathcal{N}	\mathcal{N}			\mathcal{N}	\mathcal{N}	\mathcal{N}			
5	\mathcal{N}	\mathcal{N}	\mathcal{N}		\mathcal{N}	\mathcal{N}	\mathcal{N}				
4	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}					
3	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}						
2	\mathcal{N}	\mathcal{P}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	
1	\mathcal{N}	\mathcal{N}	\mathcal{P}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	
0	\mathcal{P}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	
	0	1	2	3	4	5	6	7	8	9	10

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4	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}					
3	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{P}					
2	\mathcal{N}	\mathcal{P}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}
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0	\mathcal{P}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}
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3	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{P}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}	\mathcal{N}
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Odd Game P-positions

Theorem

If we are only allowed to take from an odd number of piles, the P-positions are the same as the ones in Nim.

Main idea of proof:

- Show that the nim-sum of the position has to change when we use the new moves
- Use the strategy in Nim to get back to a position with zero nim-sum
- Since $(0, \dots, 0) \in \mathcal{P}$, P-positions will be the same

Degrees of Freedom of P-positions

Theorem

For a position with $n - 1$ numbers known, and one number unknown: $P = (a_1, \dots, a_{n-1}, x)$, there is a unique value of x such that $P \in \mathcal{P}$.

- In general, does there exist a function $f(a_1, \dots, a_{n-1}) = x$?
- For Nim, this function is $f_{\text{NIM}}(a_1, \dots, a_{n-1}) = \bigoplus_{i=1}^{n-1} a_i$.

Bounds on P-Positions

- General bound that holds for all rectangular games

Theorem

If $P = (a_1, \dots, a_n) \in \mathcal{P}$ then $2(\sum_{i=1}^n a_i - a_j) \geq a_j$.

- If no move allows us to take from exactly two of the piles

Theorem

If $P = (a_1, \dots, a_n) \in \mathcal{P}$ then $\sum_{i=1}^n a_i - a_j \geq a_j$.

- Both proved by strong induction on $\sum_{i=1}^n a_i - a_j$

Enumeration of P-positions

- This helps to get a sense of the structure and distribution of the P-positions
- We can enumerate based on total sum
- We calculate this sequence for Nim:
 - Three piles, total sum n : $3^{\text{wt}(n)}$ if n is even, 0 otherwise, where $\text{wt}(n)$ is the number of ones in the binary representation of n .
 - 1, 0, 1, 0, 3, 0, 3, 0, 9, 0, 3, 0, 9, 0, 9, 0, 27, 0, ...
- We also calculate for all other games (we show Cookie Monster game with three piles)
 - Three piles, total sum n : 1, 0, 0, 6, 0, 0, 3, 3, 6, 0, 3, 9, 6, 6, 0, 0, 15, 3, ...

Future Research

- Find an explicit formula for P-positions in our games
- Is there a metric that describes how closely the game behaves to Nim and Wythoff?
- Connection between consecutive game and Lie Algebras
- What happens in the misère version?

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